Real-Time Systems

Lecture 03: Duration Calculus I

2013-04-23

Dr. Bernd Westphal

Albert-Ludwigs-Universität Freiburg, Germany

Contents & Goals

Last Lecture:

- Model of timed behaviour: state variables and their interpretation
- First order predicate-logic for requirements and system properties

This Lecture:

- **Educational Objectives:** Capabilities for following tasks/questions.
  - Read (and at best also write) Duration Calculus formulae.
- **Content:**
  - Classes of requirements (safety, liveness, etc.)
  - Duration Calculus:
    Assertions, Terms, Formulae, Abbreviations, Examples

Recall: Correctness

Recall: Correctness

- Let 'Req' be a **requirement**,
- 'Des' be a **design**, and
- 'Impl' be an **implementation**.

Recall: each is a set of evolutions, i.e. a subset of $\big(\text{Time} \to \times_{i=1}^{n} \mathcal{D}(obs_i)\big)$, described in any form.

We say

- 'Des' is a **correct design** (wrt. 'Req') if and only if

$$\text{Des} \subseteq \text{Req}.$$

- 'Impl' is a **correct implementation** (wrt. 'Des' (or 'Req')) if and only if

$$\text{Impl} \subseteq \text{Des} \quad (\text{or } \text{Impl} \subseteq \text{Req})$$

If 'Req' and 'Des' are described by formulae of first-oder predicate logic, proving the design correct amounts to proving that 'Des $\implies$ Req' is valid.

Classes of Timed Properties

Safety Properties

- A **safety property** states that
  something bad must never happen [Lamport].
- Example: train inside level crossing with gates open.
- More general, assume observable $C : \text{Time} \to \{0, 1\}$ where $C(t) = 1$ represents a critical system state at time t.
  Then

$$\forall\, t \in \text{Time} \bullet \neg C(t)$$

  is a safety property.
- In general, a safety property is characterised as a property that can be **falsified** in bounded time.
- But safety is not everything...

Liveness Properties

- The simplest form of a **liveness property** states that
  something good eventually does happen.
- Example: gates open for road traffic.
- More general, assume observable $G : \text{Time} \to \{0,1\}$ where $G(t) = 1$ represents a good system state at time t.

  Then
  $$\exists\, t \in \text{Time} \bullet G(t)$$
  is a liveness property.
- Note: not falsified in finite time.
- With real-time, liveness is too weak...

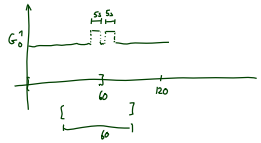
Bounded Response Properties

- A **bounded response property** states that
  the desired reaction on an input occurs in time interval $[b,e]$.
- Example: from request to secure level crossing to gates closed.
- More general, re-consider good thing $G : \text{Time} \to \{0,1\}$ and request $R : \text{Time} \to \{0,1\}$.

  Then
  $$\forall\, t_1 \in \text{Time} \bullet (R(t_1) \implies \exists\, t_2 \in [t_1 + \cancel{10}\, b, t_1 + \cancel{120}\, e] \bullet G(t_2))$$
  is a bounded liveness property.
- This property can again be falsified in finite time.
- With gas burners, this is still not everything...

Duration Properties

- A **duration property** states that
  for observation interval $[b,e]$ characterised by a condition $A(b,e)$
  the **accumulated time** in which the system is in a certain critical state has an upper bound $u(b,e)$.
- Example: leakage in gas burner.



Duration Properties

- A **duration property** states that
  for observation interval $[b,e]$ characterised by a condition $A(b,e)$
  the **accumulated time** in which the system is in a certain critical state has an upper bound $u(b,e)$.
- Example: leakage in gas burner.
- More general, re-consider critical thing $C : \text{Time} \to \{0,1\}$.

  Then
  $$\forall\, b, e \in \text{Time} \bullet \left(A(b,e) \implies \int_b^e C(t)\, dt \le u(b,e) \right)$$
  is a duration property.
- This property can again be falsified in finite time.

  gas burner: $A(b,e) := e - b \geq 60$
  $u(b,e) := \frac{e-b}{20} - \frac{1}{20}(e-b)$

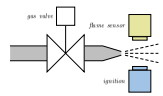
Duration Calculus

Duration Calculus: Preview

- Duration Calculus is an **interval logic**.
- Formulae are evaluated in an (**implicitly given**) interval.

Back to our gas burner:

- $G, F, I, H : \text{Time} \to \{0,1\}$
- Define $L : \text{Time} \to \{0,1\}$ as $G \wedge \neg F$.



Strangest operators:

- **everywhere** — Example: $\lceil G \rceil$
  (Holds in a given interval $[b,e]$ iff the gas valve is open almost everywhere.)
- **chop** — Example: $(\lceil \neg I \rceil \,;\, \lceil I \rceil \,;\, \lceil \neg I \rceil) \implies \ell \geq 1$
  (Ignition phases last at least one time unit.)
- **integral** — Example: $\ell \geq 60 \implies \int L \leq \frac{\ell}{20}$
  (At most 5% leakage time within intervals of at least 60 time units.)

Duration Calculus: Overview

We will introduce three (or five) syntactical "levels":

(i) **Symbols:**

$p, q,$ $f, g,$ $true, false, =, <, >, \leq, \geq$ x, y, z $X, Y, Z,$ d

(ii) **State Assertions:**

$P ::= 0 \mid 1 \mid X = d \mid \neg P_1 \mid P_1 \wedge P_2$

(iii) **Terms:**

$\theta ::= x \mid \ell \mid f() \mid f(\theta_1, \ldots, \theta_n)$

(iv) **Formulae:**

$F ::= p(\theta_1, \ldots, \theta_n) \mid \neg F_1 \mid F_1 \wedge F_2 \mid \forall x \bullet F_1 \mid F_1 ; F_2$

(v) **Abbreviations:**

$\lceil\rceil, \quad \lceil P \rceil, \quad \lceil P \rceil^t, \quad \lceil P \rceil^{\leq t}, \quad \Diamond F, \quad \Box F$

Symbols: Syntax

- f, g: **function symbols**, each with arity $n \in \mathbb{N}_0$.
  Called **constant** if $n = 0$.
  Assume: constants $0, 1, \cdots \in \mathbb{N}_0$; binary '$+$' and '$\cdot$';
- p, q: **predicate symbols**, also with arity.
  Assume: constants $true, false$; binary $=, <, >, \leq, \geq$.
- $x, y, z \in$ GVar: **global variables**.
- $X, Y, Z \in$ Obs: **state variables** or **observables**, each of a data type $\mathcal{D}$ (or $\mathcal{D}(X), \mathcal{D}(Y), \mathcal{D}(Z)$ to be precise).
  Called **boolean observable** if data type is $\{0, 1\}$.
- d: **elements** taken from data types $\mathcal{D}$ of observables.

Symbols: Semantics

- **Semantical domains** are
  - the **truth values** $\mathbb{B} = \{\mathrm{tt}, \mathrm{ff}\}$,
  - the **real numbers** $\mathbb{R}$,
  - **time** Time,
    (mostly Time $= \mathbb{R}_0^+$ (continuous), exception Time $= \mathbb{N}_0$ (discrete time))
  - and **data types** $\mathcal{D}$.
- The semantics of an n-ary **function symbol** f is a (mathematical) function from $\mathbb{R}^n$ to $\mathbb{R}$, denoted $\hat{f}$, i.e.
  $$\hat{f} : \mathbb{R}^n \to \mathbb{R}.$$
- The semantics of an n-ary **predicate symbol** p is a function from $\mathbb{R}^n$ to $\mathbb{B}$, denoted $\hat{p}$, i.e.
  $$\hat{p} : \mathbb{R}^n \to \mathbb{B}.$$
- For constants (arity $n = 0$) we have $\hat{f} \in \mathbb{R}$ and $\hat{p} \in \mathbb{B}$.

Symbols: Examples

- The **semantics** of the function and predicate symbols **assumed above** is fixed throughout the lecture:
  - $\hat{true} = \mathrm{tt}$, $\hat{false} = \mathrm{ff}$
  - $\hat{0} \in \mathbb{R}$ is the (real) number **zero**, etc.
  - $\hat{+} : \mathbb{R}^2 \to \mathbb{R}$ is the **addition** of real numbers, etc.
  - $\hat{=} : \mathbb{R}^2 \to \mathbb{B}$ is the **equality** relation on real numbers,
  - $\hat{<} : \mathbb{R}^2 \to \mathbb{B}$ is the **less-than** relation on real numbers, etc.
- "Since the semantics is the expected one, we shall often simply use the symbols $0, 1, +, \cdot, =, <$ when we mean their semantics $\hat{0}, \hat{1}, \hat{+}, \hat{\cdot}, \hat{=}, \hat{<}$."

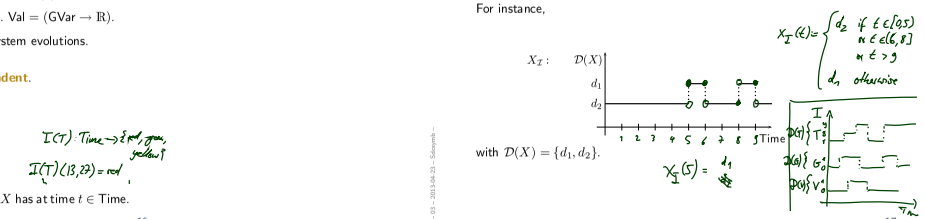
Symbols: Semantics

- The semantics of a **global variable** is not fixed (throughout the lecture) but given by a **valuation**, i.e. a mapping
  $$\mathcal{V} : \mathrm{GVar} \to \mathbb{R}$$
  assigning each global variable $x \in \mathrm{GVar}$ a real number $\mathcal{V}(x) \in \mathbb{R}$.
  We use Val to denote the set of all valuations, i.e. $\mathrm{Val} = (\mathrm{GVar} \to \mathbb{R})$.
  Global variables are though **fixed over time** in system evolutions.
- The semantics of a **state variable** is **time-dependent**.
  It is given by an interpretation $\mathcal{I}$, i.e. a mapping
  $$\mathcal{I} : \mathrm{Obs} \to (\mathrm{Time} \to \mathcal{D})$$
  assigning each state variable $X \in \mathrm{Obs}$ a function
  $$\mathcal{I}(X) : \mathrm{Time} \to \mathcal{D}(X)$$
  such that $\mathcal{I}(X)(t) \in \mathcal{D}(X)$ denotes the value that X has at time $t \in$ Time.

Symbols: Representing State Variables

- For convenience, we shall abbreviate $\mathcal{I}(X)$ to $X_{\mathcal{I}}$.
- An **interpretation** (of a state variable) can be displayed in form of a **timing diagram**.

  For instance,



  with $\mathcal{D}(X) = \{d_1, d_2\}$.

Duration Calculus: Overview

We will introduce three (or five) syntactical "levels":

(i) **Symbols:**

$$f, g, \quad true, false, =, <, >, \leq, \geq, \quad x, y, z, \quad X, Y, Z, \quad d$$

(ii) **State Assertions:**

$$P ::= 0 \mid 1 \mid X = d \mid \neg P_1 \mid P_1 \wedge P_2$$

(iii) **Terms:**

$$\theta ::= x \mid \ell \mid \int P \mid f(\theta_1, \ldots, \theta_n)$$

(iv) **Formulae:**

$$F ::= p(\theta_1, \ldots, \theta_n) \mid \neg F_1 \mid F_1 \wedge F_2 \mid \forall x \bullet F_1 \mid F_1 ; F_2$$

(v) **Abbreviations:**

$$\lceil \rceil, \quad \lceil P \rceil, \quad \lceil P \rceil^t, \quad \lceil P \rceil^{\leq t}, \quad \Diamond F, \quad \Box F$$

State Assertions: Syntax

- The set of **state assertions** is defined by the following grammar:

$$P ::= 0 \mid 1 \mid X = d \mid \neg P_1 \mid P_1 \wedge P_2$$

with $d \in \mathcal{D}(X)$.

We shall use P, Q, R to denote state assertions.

- **Abbreviations:**
  - We shall write X instead of $X = 1$ if $\mathcal{D}(X) = \mathbb{B}$.
  - Define $\vee$, $\Longrightarrow$, $\Longleftrightarrow$ as usual.

State Assertions: Semantics

- Given an evaluation $\mathcal{I}$.
- The **semantics** of **state assertion** P is a function

$$\mathcal{I}[\![P]\!] : \text{Time} \to \{0, 1\}$$

i.e. $\mathcal{I}[\![P]\!](t)$ denotes the truth value of P at time $t \in \text{Time}$.

- The value is defined **inductively** on the structure of P:

$$\mathcal{I}[\![0]\!](t) = \hat{0} \in \mathbb{R}, \quad \hat{0} = 0$$

$$\mathcal{I}[\![1]\!](t) = \hat{1} = 1 \in \mathbb{R}$$

$$\mathcal{I}[\![X = d]\!](t) = \begin{cases} 1, & \text{if } X_{\mathcal{I}}(t) = \hat{d} \\ 0, & \text{otherwise} \end{cases}$$

$$\mathcal{I}[\![\neg P_1]\!](t) = 1 - \mathcal{I}[\![P_1]\!](t)$$

$$\mathcal{I}[\![P_1 \wedge P_2]\!](t) = \begin{cases} 1, & \text{if } \mathcal{I}[\![P_1]\!](t) = \mathcal{I}[\![P_2]\!](t) = 1 \\ 0, & \text{otherwise} \end{cases}$$

State Assertions: Notes

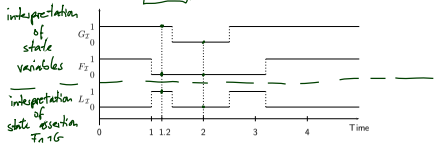
- $\mathcal{I}[\![X]\!](t) = \mathcal{I}[\![X = 1]\!](t) = \mathcal{I}(X)(t) = X_{\mathcal{I}}(t)$, if X boolean, i.e. $\mathcal{D}(X) = \mathbb{B}$
- $\mathcal{I}[\![P]\!]$ is also called **interpretation** of P.

  We shall write $P_{\mathcal{I}}$ for it.

- Here we prefer 0 and 1 as boolean values (instead of tt and ff) — for reasons that will become clear immediately.

State Assertions: Example

- Boolean observables G and F.
- State assertion $L := G \wedge \neg F$,



- $L_{\mathcal{I}}(1.2) = 1$, because
- $L_{\mathcal{I}}(2) = 0$, because

References

References

[Olderog and Dierks, 2008] Olderog, E.-R. and Dierks, H. (2008). *Real-Time Systems - Formal Specification and Automatic Verification*. Cambridge University Press.