Real-Time Systems

Lecture 17: Automatic Verification of DC Properties for TA

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Contents & Goals

Last Lecture:

- Undecidability Results for TBA

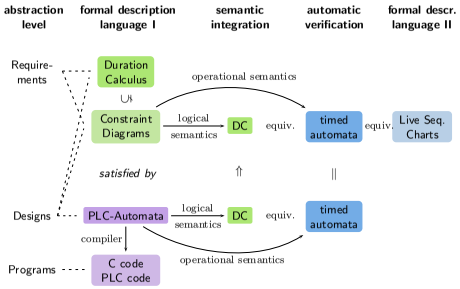
This Lecture:

- **Educational Objectives:** Capabilities for following tasks/questions.
  - How can we relate TA and DC formulae? What's a bit tricky about that?
  - Can we use Uppaal to check whether a TA satisfies a DC formula?
- **Content:**
  - An evolution-of-observables semantics of TA
  - A satisfaction relation between TA and DC
  - Model-checking DC properties with Uppaal

You Are Here

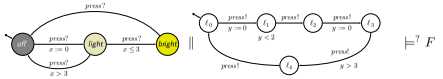
Content

Tying It All Together



Observer-based Automatic Verification of DC Properties for TA

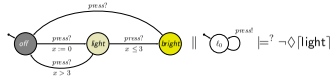
Model-Checking DC Properties with Uppaal



$\models^? F$

- **First Question**: what is the "$\models$" here?
- **Second Question**: what kinds of DC formulae can we check with Uppaal?
  - **Clear**: Not every DC formula.
    (Otherwise contradicting undecidability results.)
  - **Quite clear**: $F = \Box\lceil\text{off}\rceil$ or $F = \neg\Diamond\lceil\text{light}\rceil$
    (Use Uppaal's fragment of TCTL, something like $\forall\Box$off,
    but not exactly (see later).)
  - **Maybe**: $F = \ell > 5 \implies \Diamond\lceil\text{off}\rceil^5$
  - **Not so clear**: $F = \neg\Diamond(\lceil\text{bright}\rceil\,;\lceil\text{light}\rceil)$

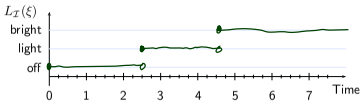
Example: Let's Start With Single Runs



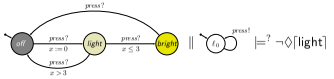
$\models^? \neg\Diamond\lceil\text{light}\rceil$

$$\xi = \langle \begin{smallmatrix}\text{off}\\0\end{smallmatrix}\rangle, 0 \xrightarrow{2.5} \langle \begin{smallmatrix}\text{off}\\2.5\end{smallmatrix}\rangle, 2.5 \xrightarrow{\tau} \langle \begin{smallmatrix}\text{light}\\0\end{smallmatrix}\rangle, 2.5 \xrightarrow{2.0} \langle \begin{smallmatrix}\text{light}\\2.0\end{smallmatrix}\rangle, 4.5 \xrightarrow{\tau} \langle \begin{smallmatrix}\text{bright}\\2.0\end{smallmatrix}\rangle, 4.5\ldots$$

Construct interpretation $L_{\mathcal{I}}(\xi)$: Time $\to$ {off, light, bright}:

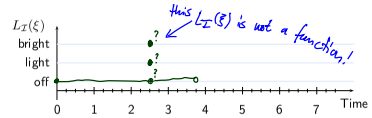


Example 2: Another Single Run



$\models^? \neg\Diamond\lceil\text{light}\rceil$

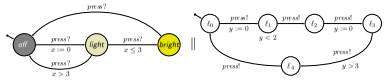
$$\xi = \langle \begin{smallmatrix}\text{off}\\0\end{smallmatrix}\rangle, 0 \xrightarrow{2.5} \langle \begin{smallmatrix}\text{off}\\2.5\end{smallmatrix}\rangle, 2.5 \xrightarrow{\tau} \langle \begin{smallmatrix}\text{light}\\0\end{smallmatrix}\rangle, 2.5 \xrightarrow{\tau} \langle \begin{smallmatrix}\text{bright}\\0\end{smallmatrix}\rangle, 2.5 \xrightarrow{\tau} \langle \begin{smallmatrix}\text{off}\\0\end{smallmatrix}\rangle, 2.5 \xrightarrow{1.0} \ldots$$



We know this problem from the exercises...

Observing Timed Automata

DC Properties of Timed Automata



$\models^? F$

Wanted: A satisfaction relation between networks of timed automata and DC formulae, a notion of $\mathcal{N}$ **satisfies** F, denoted by $\mathcal{N} \models F$.

Plan:

- Consider network $\mathcal{N}$ consisting of TA
  $$\mathcal{A}_{e,i} = (L_i, C_i, B_i, U_i, X_i, V_i, I_i, E_i, \ell_{ini,i})$$
- Define observables Obs($\mathcal{N}$) of $\mathcal{N}$.
- Define evolution $\mathcal{I}_\xi$ of Obs($\mathcal{N}$) induced by computation path $\xi \in CompPaths(\mathcal{N})$ of $\mathcal{N}$,
  $CompPaths(\mathcal{N}) = \{\xi \mid \xi \text{ is a computation path of } \mathcal{N}\}$
- Say $\mathcal{N} \models F$ if and only if $\forall\, \xi \in CompPaths(\mathcal{N}) : \mathcal{I}_\xi \models_0 F$.

Observables of TA Network

Let $\mathcal{N}$ be a network of n extended timed automata

$$\mathcal{A}_{e,i} = (L_i, C_i, B_i, U_i, X_i, V_i, I_i, E_i, \ell_{ini,i})$$

(locations, clocks, variables)

For simplicity: assume that the L_i and X_i are pairwise disjoint and that each V_i is pairwise disjoint to every L_i and X_i (otherwise rename).

- **Definition**: The observables Obs($\mathcal{N}$) of $\mathcal{N}$ are
  $$\{\ell_1, \ldots, \ell_n\} \cup \bigcup_{1 \le i \le n} V_i$$
  (current location of ate, (would be less confusing if we used $\{\mathbb{Q}, \ldots, \mathbb{Q}_n\}$)

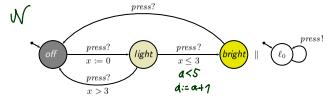
  with
  - $\mathcal{D}(\ell_i) = L_i$,
  - $\mathcal{D}(v)$ as given, $v \in V_i$.

Observables of TA Network: Example

$\mathcal{A}_{n,i} = (L_i, C_i, B_i, U_i, X_i, V_i, I_i, E_i, \ell_{ini,i})$.

The observables Obs($\mathcal{N}$) of $\mathcal{N}$ are $\{\ell_1, \ldots, \ell_n\} \cup \bigcup_{1 \le i \le n} V_i$ with

- $\mathcal{D}(\ell_i) = L_i$,
- $\mathcal{D}(v)$ as given, $v \in V_i$.



Obs($\mathcal{N}$) = $\{\ell_1, \ell_2, a\}$

$\mathcal{D}(\ell_1)$ = {off, light, bright} $\mathcal{D}(\ell_2) = \{\ell_0\}$

$\mathcal{D}(a) = \{0, .., 5\}$

References

References

[Olderog and Dierks, 2008] Olderog, E.-R. and Dierks, H. (2008). *Real-Time Systems - Formal Specification and Automatic Verification*. Cambridge University Press.