

Formal Methods for C

Seminar – Summer Semester 2014

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Tools & Modules

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Hello, Again

```
1 #include <stdio.h>
2
3 int g( int x ) { return x/2; }
4
5 int f() { return g(1); }
6
7 int main() {
8     printf( "HelloWorld\n" );
9     return f();
10 }
```

- % gcc helloWorld.c
- % ls
- a.out helloWorld.c
- % ./a.out
- Hello World.
- % echo \$?

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Zoom In: Preprocessing, Compiling, Linking

```
1 #include <stdio.h>
2
3 int g( int x ) { return x/2; }
4
5 int f() { return g(1); }
6
7 int main() {
8     printf( "HelloWorld\n" );
9     return f();
10 }
```

- % gcc -E helloWorld.c > helloWorld.i
- % gcc -c helloWorld.i
- % ld -o helloWorld.o [...] helloWorld.o [...]
- % ./helloWorld
- Hello World.
- %

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Modules

```
1 #include <stdio.h>
2
3 int g( int x ) {
4     return x/2;
5 }
6
7 int f() {
8     return g(1);
9 }
10
11 int main() {
12     printf( "HelloWorld\n" );
13     return f();
14 }
```

- Split into:
- .h (header) declarations
- .c: definitions, use headers
- to "import" declarations

```
g.h
1 #ifndef GH
2 #define GH
3
4 extern int
5 g( int x );
6 #endif
```

```
f.c
1 #include "g.h"
2
3 int f() {
4     return g(1);
5 }
```

```
helloWorld.c
1 #include <stdio.h>
2 #include "f.h"
3
4 int main() {
5     printf( "HelloWorld\n" );
6     return f();
7 }
```

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`gdb(1)`, `ddd(1)`, `nm(1)`, `make(1)`

- **Command Line Debugger:**

`gdb a.out [core]`

- **GUI Debugger:**

`ddd a.out [core]`

(works best with debugging information compiled in (`gcc -g`))

- **Inspect Object Files:**

`nm a.o`

- **Build Utility:**

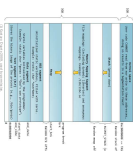
`make`

See battery controller exercise for an example.

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Core Dumps

- **Recall:** Anatomy of a Linux Program in Memory
- **Core dump:** (usually) this memory written to a file.



```
1 int main() {
2     int *p;
3     *p = 27;
4     return 0;
}
```

```
1 % gcc -g core.c
2 % ./core
3 % ./a.out
4 % limit coredumpsize 1g
5 % ./a.out
6 % ./a.out
7 % ls -lh core
8 -rw-r--r-- 1 user user 23K Feb 29 11:11 core
9 % cat core (cat) 7.41--debian
10 [...]
11 [...]
12 [...]
13 [...]
14 % ./a.out
15 % ./a.out
16 % ./a.out
17 % ./a.out
18 % ./a.out
```

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Formal Methods for C

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Correctness and Requirements

Correctness

- **Correctness is defined with respect to a specification.**
- A program (function, ...) is **correct** (wrt. specification ϕ) **if and only if** it satisfies ϕ .
- Definition of "satisfies": **in a minute.**

Examples

- ϕ_1 : the return value is 10 divided by parameter (if parameter not 0)
- ϕ_2 : the value of variable x is "always" strictly greater than 3
- ϕ_3 : the value of i increases in each loop iteration
- ...

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Common Patterns

- **State Invariants:**
"at this program point, the value of p must not be NULL"
"at all program points, the value of p must not be NULL"
(cf: **sequence points** (Annex C))
- **Data Invariants:**
"the value of n must be the length of s "
- **(Function) Pre/Post Conditions:**
Pre-Condition: the parameter must not be 0
Post-Condition: the return value is 10 divided by the parameter
- **Loop Invariants:**
"the value of i is between 0 and array length minus 1"

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Diagnostics (7.2)

```
1 #include <assert.h>
2 void assert( /* scalar */ expression );
```

Poor Man's Requirements Specification
aka. How to Formalize Requirements in C?

- "The assert macro puts diagnostic tests into programs: [...]
- When it is executed, if expression (which shall have a scalar type) is false (that is, compares equal to 0), the assert macro
- writes information about the particular call that failed [...] on the standard error stream in an implementation-defined format.
 - It then calls the `abort` function."

Pitfall:

- If macro `DEBUG` is defined when including `<assert.h>`, expression is not evaluated (thus should be side-effect free).

abort (7.20.4.1)

```
1 #include <stdlib.h>
2
3 void abort();
```

- "The abort function causes abnormal program termination to occur, unless [...]
- [...] An implementation-defined form of the status unsuccessful termination is returned to the host environment by means of the function call `raise(SIGABRT)`."

(→ Core Dumps)

Common Patterns with assert

- **State Invariants:**
"at this program point, the value of *p* must not be NULL"
"at all program points, the value of *p* must not be NULL"
(cf. **sequence points** (Annex C))
- **Data Invariants:**
the value of *n* must be the length of *s*"
- **(Function) Pre/Post Conditions:**
Pre-Condition: the parameter must not be 0
Post-Condition: the return value is 10 divided by the parameter
- **Loop Invariants:**
"the value of *i* is between 0 and array length minus 1"

State Invariants with <assert.h>

```
1 void f() {
2     int* p = (int*)malloc(sizeof(int));
3     if (!p)
4         return;
5     assert(p); // assume p is valid from here
6     // ...
7     void g() {
8         Node* p = find( 'a' );
9         assert(p); // we inserted 'a' before
10        // ...
11    }
12 }
13
14
15
16 }
```

Data Invariants with <assert.h>

```
1 typedef struct {
2     char* s;
3     int n;
4 } str;
5
6 str* construct( char* s ) {
7     str* x = (str*)malloc( sizeof(str) );
8     // ...
9     assert( (x->s == NULL && x->n == -1)
10            || (x->n == strlen( x->s ) ) );
11 }
```

Pre/Post Conditions with <asserth>

```
1 int f ( int x ) {
2   assert( x != 0 ); // pre-condition
3
4   int r = 10/x;
5
6   assert( r == 10/x ); // post-condition
7
8   return r;
9 }
```

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Loop Invariants with <asserth>

```
1 void f( int a[] , int n ) {
2   int i = 0;
3
4   // holds before the loop
5   assert( 0 <= i && i <= n );
6   while ( i <= 1 || a[i-1] == 0 );
7
8   // holds before each iteration
9   assert( i <= 1 || a[i-1] == 0 );
10
11   a[i++] = 0;
12
13   // holds after exiting the loop
14   assert( 0 <= i && i <= n );
15
16   return;
17 }
18
19
20 }
```

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Old Variables, Ghost Variables

```
1 void xorSwap( unsigned int* a, unsigned int* b ) {
2   #ifndef NDEBUG
3     #warning int *old_a = a, *old_b = b;
4     #endif
5     assert( a && b ); assert( a != b ); // pre-condition
6
7     *a = *a + *b;
8     *b = *a - *b;
9     *a = *a - *b;
10
11     assert( *a == *old_b && *b == *old_a ); // post-condition
12     assert( a == old_a && b == old_b ); // ditto
13 }
```

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Outlook

- Some verification tools simply verify for each `assert` statement:
When executed, expression is not false.
- Some verification tools support sophisticated requirements specification languages like ACSL with explicit support for
 - pre/post conditions
 - ghost variables, old values
 - data invariants
 - loop invariants
 - ...

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Dependable Verification (Jackson)

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Dependability

- **"The program has been verified,"** tells us **not very much**.
- One wants to know (and should state):
 - **Which specifications** have been considered?
 - Under **which assumptions** was the verification conducted?
 - Platform assumptions: finite words (size?), mathematical integers, ...
 - Environment assumptions, input values, ...
- Assumptions are often implicit: **"in the tool!"**
- And **what does verification mean** after all?
 - In some contexts: **testing**.
 - In some contexts: **review**.
 - In some contexts: **model-checking** procedure.
 - ("We verified the program" - "What did the tool say?" - "Verification failed.")
 - In some contexts: **model-checking tool claims correctness**.

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Distinguish

Most **generic errors** boil down to:

- specified but **unwanted behaviour**,
e.g. under/overflows
- **initialisation issues**
e.g. automatic block scope objects
- **unspecified behaviour** (J1)
e.g. order of evaluation in some cases
- **undefined behaviour** (J2)
- **implementation defined behaviour** (J3)
the compiler

Common Errors

Conformance (4)

- "A program that is
 - correct in all other aspects,
 - operating on correct data,
 - containing **unspecified behavior**
- shall be a correct program and act in accordance with 5.1.1.2.3. (Program Execution)
- A conforming program is one that is acceptable to a conforming implementation. (~~is-acceptable~~)
 - Strictly conforming programs are intended to be maximally portable among conforming implementations.
 - An implementation [of C, a compiler] shall be accompanied by a document that defines all implementation-defined and local-specific characteristics and all extensions.

Distinguish

Most **generic errors** boil down to:

- specified but **unwanted behaviour**,
e.g. under/overflows
- **initialisation issues**
e.g. automatic block scope objects
- **unspecified behaviour** (J1)
e.g. order of evaluation in some cases
- **undefined behaviour** (J2)
- **implementation defined behaviour** (J3)
the compiler

Over- and Underflows, Casting

- Not specific to C...

Over- and Underflows

```
1 void f( short a, int b ) {  
2     a = b; // Spng ok, but...  
3 }  
4 short a; // provisioning, implicit cast  
5 if (++a < 0) { /* no */ }  
6 if (++i > MAX_INT) {  
7     /* no */ }  
8  
9  
10  
11  
12 int e = 0;  
13  
14 void set-error() { e++; }  
15 void clear-error() { e = 0; }  
16  
17 void g() { if (e) { /* ... */ } }
```

Initialisation (6.7.8)

Initialisation (6.7.8)

- If an object that has automatic storage duration is not initialized explicitly, its value is indeterminate.

```
1 void f() {
2     int a;
3
4     printf( "%i\n", a ); // surprise...
5 }
```

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Unspecified Behaviour (1.1)

- Each implementation (of a compiler) documents how the choice is made.
- For example**
- whether two string literals result in distinct arrays (6.4.5)
 - the order in which the function designator, arguments, and subexpressions within the arguments are evaluated in a function call (6.5.2.2)
 - the layout of storage for function parameters (6.9.1)
 - the result of rounding when the value is out of range (7.12.9.5, ...)
 - the order and contiguity of storage allocated by successive calls to malloc (7.20.3)
- etc. pp.

```
1 char a[] = "hello", b[] = "hello"; // a == b?
2
3 i = 0; f(++i, ++i, ++i); // f(1,2,3)?
4
5 int g() { int a, b; } // &a > &b ?
6
7 int* p = malloc( sizeof( int ) );
8 int* q = malloc( sizeof( int ) ); // q > p?
```

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Undefined Behaviour (3.4.3)

"Behaviour, upon use of a non-portable or erroneous program construct or of erroneous data, for which this International Standard imposes no requirements."

Possible undefined behaviour ranges from

- ignoring the situation completely with **unpredictable results**,
- to behaving during **translation or program execution** in a documented manner characteristic of the environment (with or without the issuance of a diagnostic message),
- to terminating a **translation or execution** (with the issuance of a diagnostic message)''

"An example of undefined behaviour is the behaviour on **integer overflow**."

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Unspecified Behaviour (1.1)

Each implementation (of a compiler) documents how the choice is made.

For example

- whether two string literals result in distinct arrays (6.4.5)
 - the order in which the function designator, arguments, and subexpressions within the arguments are evaluated in a function call (6.5.2.2)
 - the layout of storage for function parameters (6.9.1)
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- etc. pp.

```
1 char a[] = "hello", b[] = "hello"; // a == b?
2
3 i = 0; f(++i, ++i, ++i); // f(1,2,3)?
4
5 int g() { int a, b; } // &a > &b ?
6
7 int* p = malloc( sizeof( int ) );
8 int* q = malloc( sizeof( int ) ); // q > p?
```

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Undefined Behaviour (1.2)

More examples:

- an identifier [...] contains an invalid multibyte character (5.2.1.2)
 - an object is referred to outside of its lifetime (6.2.4)
 - the value of a pointer to an object whose lifetime has ended is used (6.2.4)
 - conversion to or from an integer type produces a value outside the range that can be represented (6.3.1.6)
 - conversion between two pointer types produces a result that is incorrectly aligned (6.3.2.3)
 - the program attempts to modify a string literal (6.4.5)
 - an exceptional condition occurs during the evaluation of an expression (6.5)
 - the value of the second operand of the / or % operator is zero (6.5.5)
 - pointers that do not point into, or just beyond, the same array object are subtracted (6.5.6)
 - An array subscript is out of range [...] (6.5.6)
 - the program removes the definition of a macro whose name begins with an underscore and either an uppercase letter or another underscore (7.1.3)
- etc. pp.

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Null-Pointer

```
1 int main () {
2     int *p;
3     *p = 27;
4     return 0;
5 }
```

- "An integer constant expression with the value 0, or such an expression cast to type `void*` is called a **null pointer constant** [...]"
- "The macro **NULL** is defined in `<stddef.h>` (and other headers) as a null pointer constant; see 7.17."
- "Among the invalid values for dereferencing a pointer by the unary `*` operator are a null pointer [...] (6.5.3.2)"

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Segmentation Violation

```
1 int main () {
2     int *p = (int*)0x12345678;
3     *p = 27;
4
5     *(int*)((void*)p) + 1) = 13;
6     return 0;
7 }
```

- Modern operating systems provide **memory protection**.
- Accessing memory which the process is not allowed to access is observed by the operating system.
- Typically an instance of "accessing an object outside its lifetime".
- **But:** other way round does not hold, accessing an object outside its lifetime does not imply a segmentation violation.
- Some platforms (e.g. SPARC) unaligned memory access, i.e. outside word boundaries, not supported by hardware ("bus error").
- Operating system notifies process: default handler: terminate: dump core

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Implementation-Defined Behaviour (I.3)

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Implementation-Defined Behaviour (I.3)

"A conforming implementation is required to document its choice of behavior in each of the areas listed in this subclause. The following are implementation-defined:"

- J.3.2 Environment, e.g.
- The set of signals, their semantics, and their default handling (7.14)
- J.3.3 Identifiers, e.g.
- The number of significant initial characters in an identifier (5.2.4.1, 6.4.2)
- J.3.4 Characters, e.g.
- The number of bits in a byte (3.6)
- J.3.5 Integers, e.g.
- Any extended integer types that exist in the implementation (6.2.5)
- J.3.6 Floating Point, e.g.
- The accuracy of the floating-point operations [...] (5.2.4.2.2)
- J.3.7 Arrays and Pointers, e.g.
- The result of converting a pointer to an integer or vice versa (6.3.2.3)
- etc. pp.

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Locale and Common Extensions (I.4, I.5)

- J.4 Locale-specific behaviour
- J.5 Common extensions
- "The following extensions are widely used in many systems, but are not portable to all implementations."

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References

[ISO 1999] ISO (1999). Programming languages – C. Technical Report ISO/IEC 9899:1999, ISO, Second edition, 1999-12-01.

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