Softwaretechnik / Software-Engineering

Lecture 11: Structural Software Modelling II

2019-06-24

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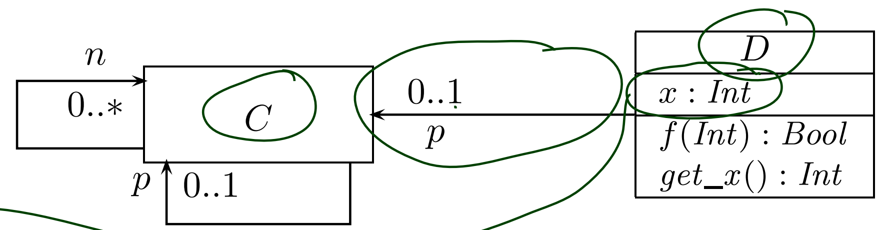
Topic Area Architecture & Design: Content

VL 10	• **Introduction and Vocabulary**
	• **Software Modelling**
⋮	– model; views / viewpoints; 4+1 view
	• **Modelling structure**
VL 11	– (simplified) Class & Object diagrams
⋮	– (simplified) Object Constraint Logic (OCL)
VL 12	• **Principles of Design**
	– modularity, separation of concerns
⋮	– information hiding and data encapsulation
	– abstract data types, object orientation
	• **Design Patterns**
VL 13	• **Modelling behaviour**
	– Communicating Finite Automata (CFA)
⋮	– Uppaal query language
	– CFA vs. Software
VL 14	– Unified Modelling Language (UML)
	– basic state-machines
⋮	– an outlook on hierarchical state-machines
	• **Model-driven/-based Software Engineering**

From Abstract to Concrete Syntax



$$\mathscr{S} = (\mathscr{T}, \mathscr{C}, V, atr, F, mth)$$

- $\mathscr{T} = \{Int, Bool\}$
- $\mathscr{C} = \{C, D\}$
- $V = \{x : Int, p : C_{0,1}, n : C_*\}$
- $atr = \{C \mapsto \{n, p\}, D \mapsto \{x, p\}\}$
- $F = \{f : Int \to Bool, \ldots$
- $mth = \{C \mapsto \emptyset, \ldots$

Basic Object System Structure Example

Wanted: a structure for signature

$$\mathscr{S}_0 = (\{Int, Bool\}, \{C, D\}, \{x : Int, p : C_{0,1}, n : C_*\}, \{C \mapsto \{p, n\}, D \mapsto \{p, x\}\},$$
$$\{f : Int \to Bool, get\_x : Int\}, \{C \mapsto \emptyset, D \mapsto \{f, get\_x\}\})$$

A structure $\mathscr{D}$ maps

- $\tau \in \mathscr{T}$ to **some** $\mathscr{D}(\tau)$, $C \in \mathscr{C}$ to **some** identities $\mathscr{D}(C)$ (infinite, pairwise disjoint),
- C_* and $C_{0,1}$ for $C \in \mathscr{C}$ to $\mathscr{D}(C_{0,1}) = \mathscr{D}(C_*) = 2^{\mathscr{D}(C)}$.

$$\begin{aligned}
\mathscr{D}(Int) &= \mathbb{Z} \\
\mathscr{D}(C) &= \mathbb{N}^+ \times \{C\} = \{1_C, 2_C, 3_C, \dots\} \\
\mathscr{D}(D) &= \mathbb{N}^+ \times \{D\} = \{1_D, 2_D, 3_D, \dots\} \\
\mathscr{D}(C_{0,1}) = \mathscr{D}(C_*) &= 2^{\mathscr{D}(C)} \\
\mathscr{D}(D_{0,1}) = \mathscr{D}(D_*) &= 2^{\mathscr{D}(D)}
\end{aligned}$$

$\mathscr{D}'$: $\{3, 17, 25\}$

$\{\bullet, \blacktriangle, \dots\}$

$\{a, aa, aaa, \dots\}$

System State Examples

$$\mathscr{S}_0 = (\{Int, Bool\}, \{C, D\}, \{x : Int, p : C_{0,1}, n : C_*\}, \{C \mapsto \{p, n\}, D \mapsto \{p, x\}\},$$
$$\{f : Int \to Bool, get\_x : Int\}, \{C \mapsto \emptyset, D \mapsto \{f, get\_x\}\})$$
$$\mathscr{D}(Int) = \mathbb{Z}, \quad \mathscr{D}(C) = \{1_C, 2_C, 3_C, ...\}, \quad \mathscr{D}(D) = \{1_D, 2_D, 3_D, ...\}$$

A system state is a partial function $\sigma : \mathscr{D}(\mathscr{C}) \nrightarrow (V \nrightarrow (\mathscr{D}(\mathscr{T}) \cup \mathscr{D}(\mathscr{C}_*)))$ such that

- $\mathrm{dom}(\sigma(u)) = atr(C)$,
- $\sigma(u)(v) \in \mathscr{D}(\tau)$ if $v : \tau, \tau \in \mathscr{T}$,
- $\sigma(u)(v) \in \mathscr{D}(C_*)$ if $v : D_*$ or $v : D_{0,1}$ with $D \in \mathscr{C}$.

$$\sigma_1 = \{ \underset{\substack{\uparrow \\ \mathscr{D}(\mathscr{C})}}{2_C} \mapsto \{ \overbrace{p \mapsto \{2_C\}}^{\text{link}}, n \mapsto \emptyset \}, 1_D \mapsto \{ p \mapsto \{2_C\}, x \mapsto 27 \} \}$$

$$\sigma_2 = \emptyset$$

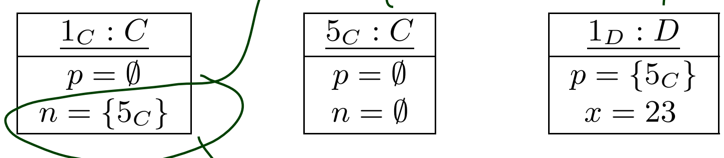
$$\sigma_3 = \{ 5_C \mapsto \{ p \mapsto \{13_C\}, n \mapsto \emptyset \} \} \quad \checkmark$$

Object Diagrams

$$\mathscr{S}_0 = (\{Int, Bool\}, \{C, D\}, \{x : Int, p : C_{0,1}, n : C_*\}, \{C \mapsto \{p, n\}, D \mapsto \{p, x\}\},$$
$$\{f : Int \to Bool, get\_x : Int\}, \{C \mapsto \emptyset, D \mapsto \{f, get\_x\}\}), \qquad \mathscr{D}(Int) = \mathbb{Z}$$

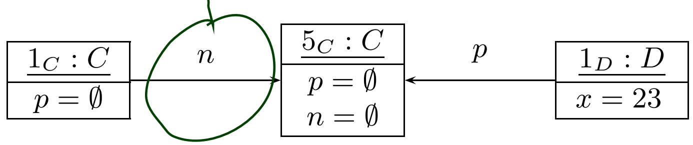
$$\sigma = \{1_C \mapsto \{p \mapsto \emptyset, n \mapsto \{5_C\}\}, 5_C \mapsto \{p \mapsto \emptyset, n \mapsto \emptyset\}, 1_D \mapsto \{p \mapsto \{5_C\}, x \mapsto 23\}\}$$

- We may **represent** σ graphically as follows:

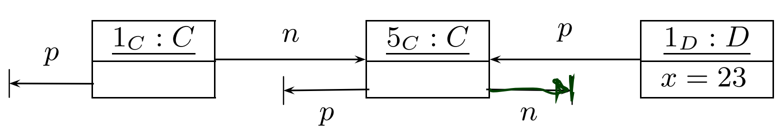


This is an **object diagram**.

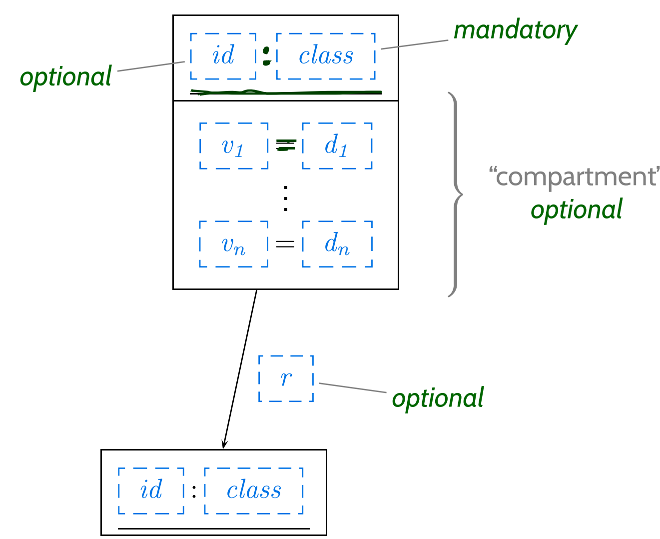
- Alternative notation:



- Alternative **non-standard** notation:



Concrete Syntax:



Content

- **Object Diagrams Cont'd**
  - dangling references
  - partial vs. complete
  - object diagrams at work
- **Proto-OCL**
  - syntax, semantics
  - Proto-OCL vs. OCL
  - Putting It All Together: Proto-OCL vs. Software

Object Diagrams Cont'd

Special Case: Dangling Reference

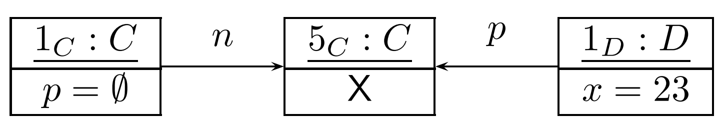
Definition.
Let $\sigma \in \Sigma_{\mathscr{S}}^{\mathscr{D}}$ be a system state and $u \in \mathrm{dom}(\sigma)$ an alive object of class C in σ.

We say $r \in atr(C)$ **is a dangling reference in** u if and only if
$r : C_{0,1}$ or $r : C_*$ and u refers to a **non-alive** object via r, i.e.

$$(\sigma(u))(r) \not\subseteq \mathrm{dom}(\sigma).$$

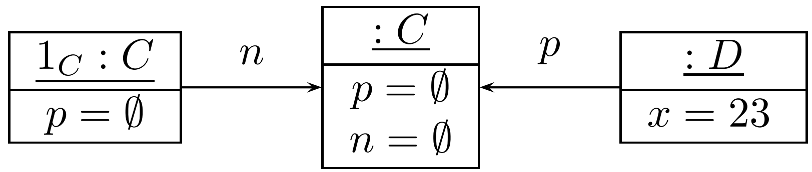
Example:

- $\sigma = \{1_C \mapsto \{p \mapsto \emptyset, n \mapsto \{5_C\}\}, 1_D \mapsto \{p \mapsto \{5_C\}, x \mapsto 23\}\}$
- Object diagram representation:



Special Case: Anonymous Objects

If the object diagram



is considered as **complete**, then it denotes the set of all system states

$$\{1_C \mapsto \{p \mapsto \emptyset, n \mapsto \{c\}\}, \boxed{c} \mapsto \{p \mapsto \emptyset, n \mapsto \emptyset\}, \boxed{d} \mapsto \{p \mapsto \{c\}, x \mapsto 23\}\}$$

where $\quad c \in \mathscr{D}(C), \quad d \in \mathscr{D}(D), \quad \underline{c \neq 1_C}.$

Intuition: different boxes represent different objects.

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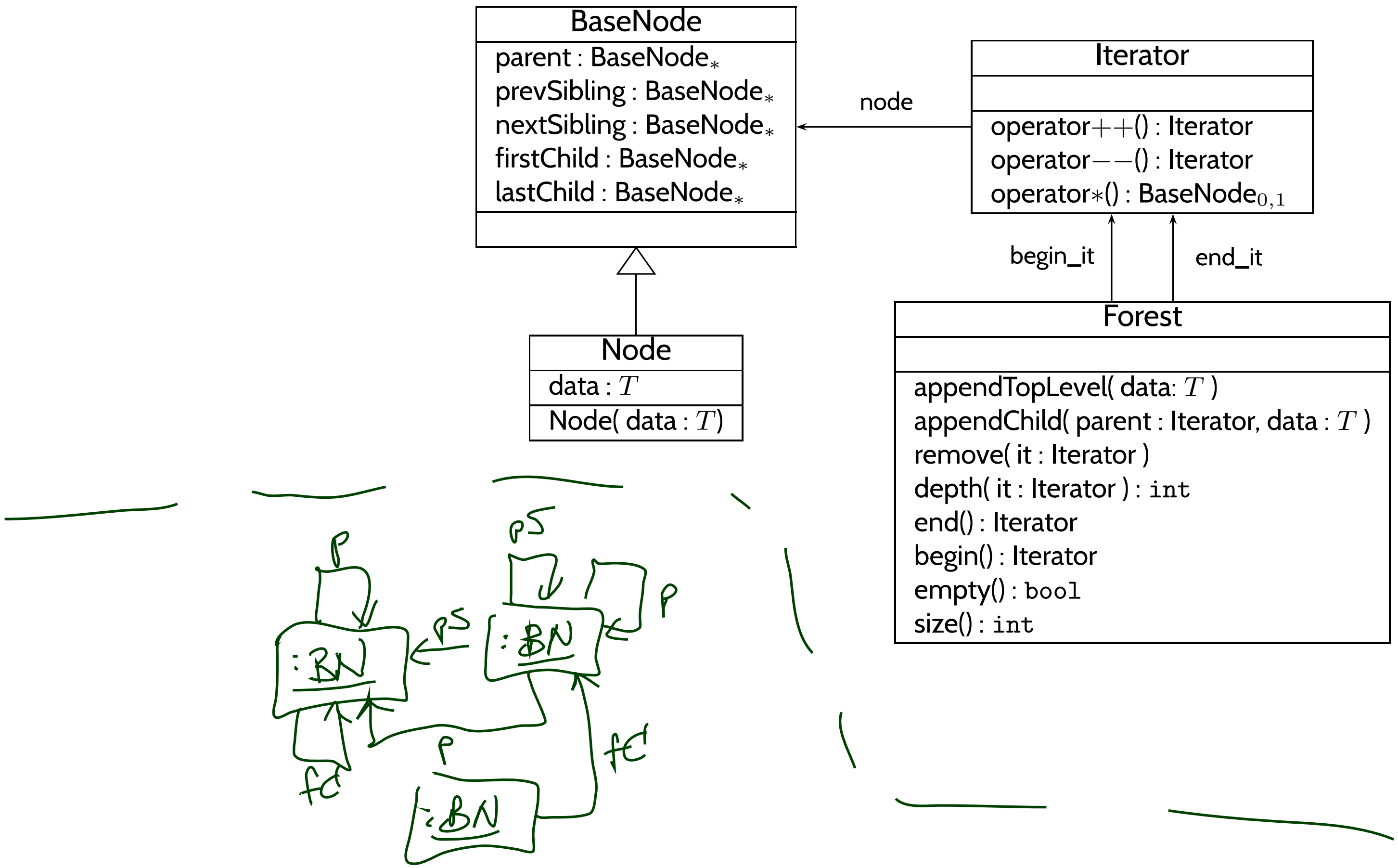
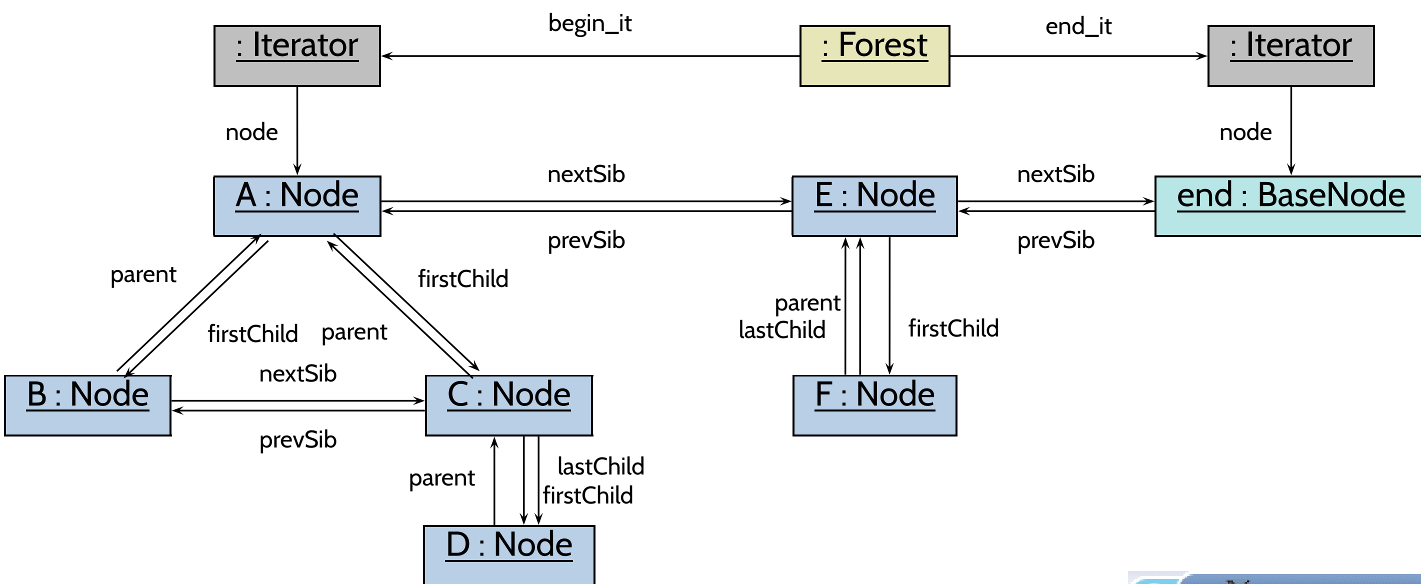
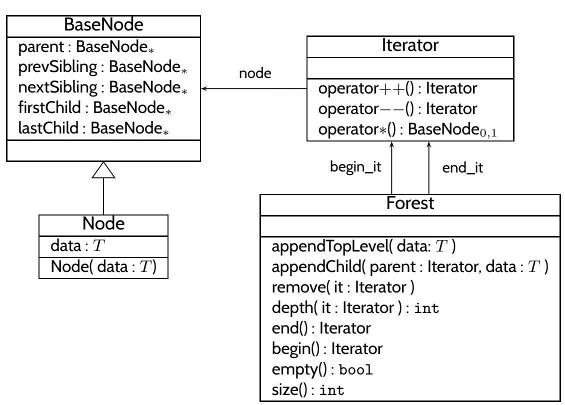
Object Diagrams at Work

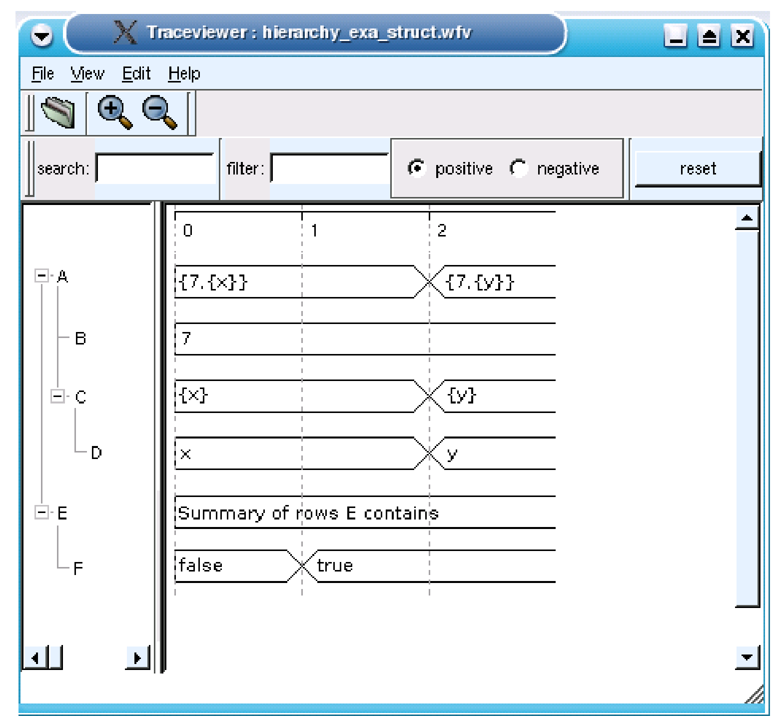
Example: Data Structure (Schumann et al., 2008)



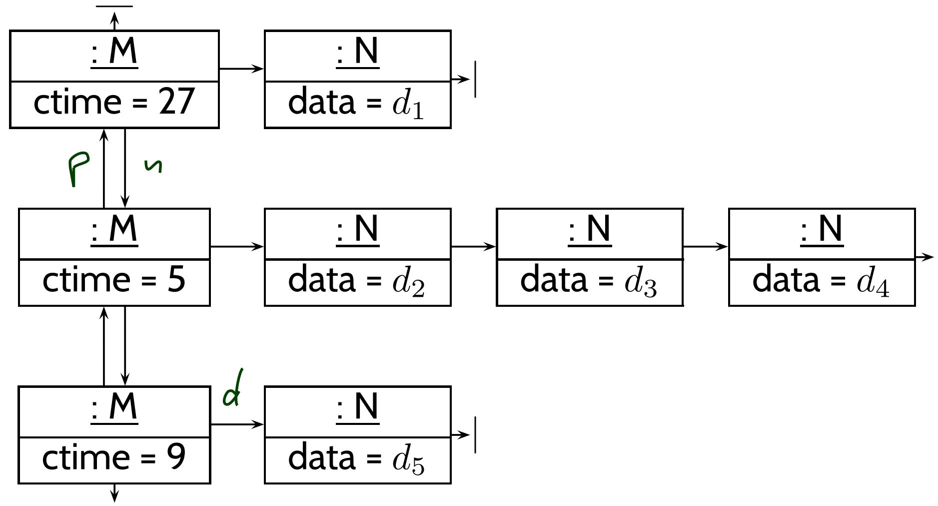
Example: Illustrative Object Diagram (Schumann et al., 2008)

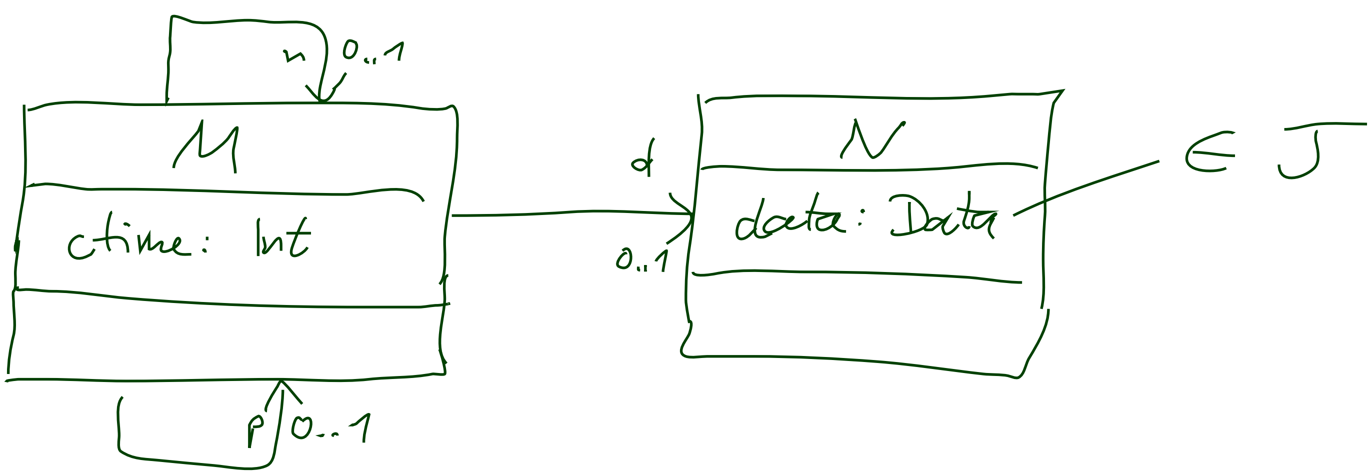




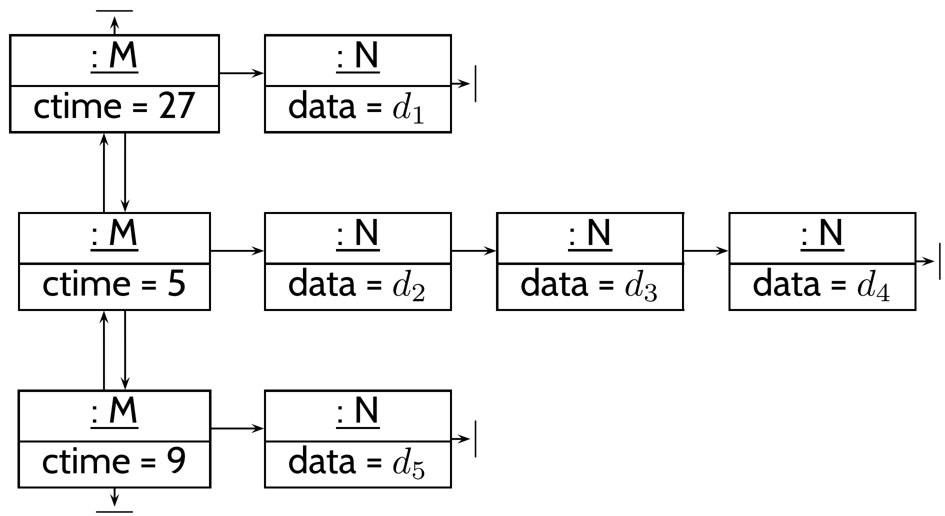


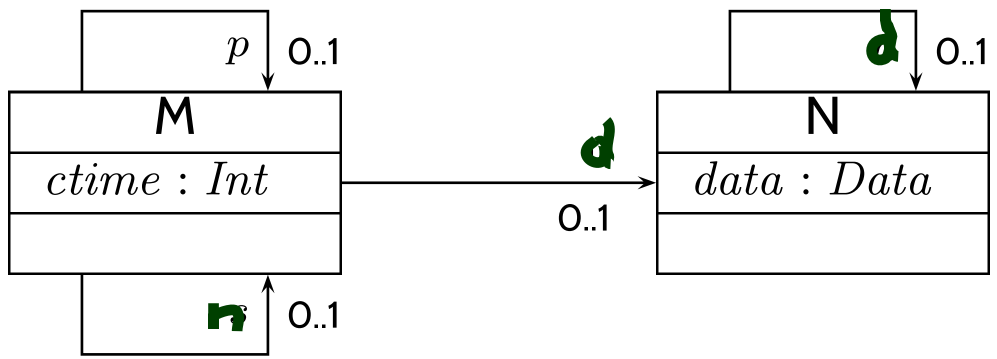
Object Diagrams for Structural Analysis





Object Diagrams for Structural Analysis





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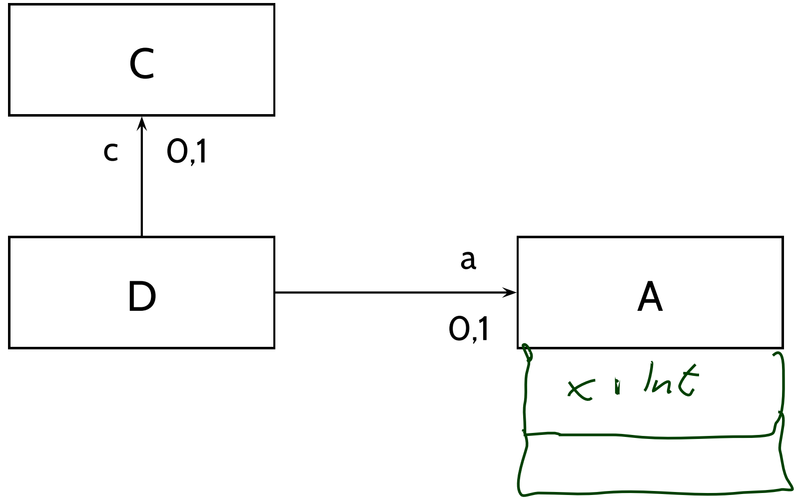
Towards Object Constraint Logic (OCL)
— “Proto-OCL” —

Motivation

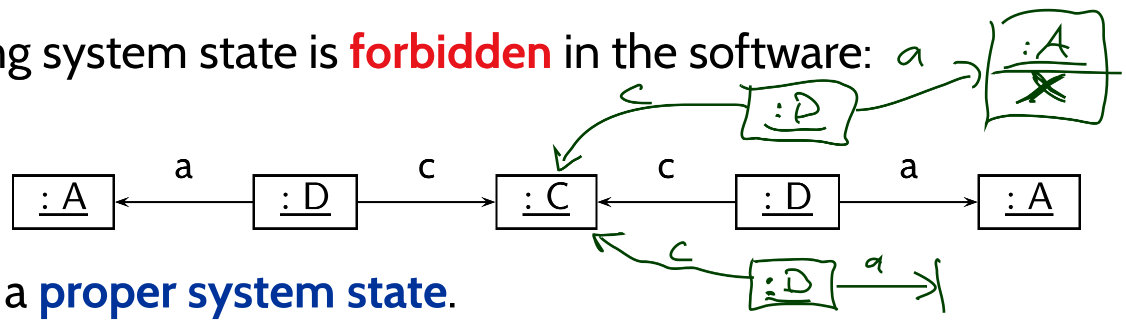


- How do I **precisely, formally** tell **my developers** that

  All D-instances having a link to the same C object must have links to the same A.

  $x(a(d_1))$

- That is, the following system state is **forbidden** in the software:

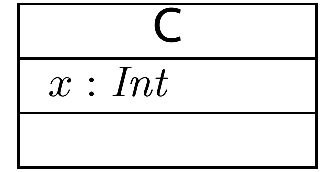


  Note: formally, it is a **proper system state**.

- Use **(Proto-)OCL**: “Dear developers, please only use system states which satisfy:”

$$\forall\, d_1 \in allInstances_D \bullet \forall\, d_2 \in allInstances_D \bullet c(d_1) = c(d_2) \implies a(d_1) = a(d_2)$$

Constraints on System States: Proto-OCL Syntax



- **Example**: for all C-instances, x should never have the value 27.

$$\forall\, c \in \mathit{allInstances}_C \bullet x(c) \neq 27$$

Definition. Proto-OCL Formulae wrt. signature $(\mathscr{T}, \mathscr{C}, V, \mathit{atr}, F, \mathit{mth})$
(c is a logical variable, $C \in \mathscr{C}$):

$$
\begin{array}{rlll}
F ::= & c & : \tau_C & \\
| & \mathit{allInstances}_C & : 2^{\tau_C} & \\
| & v(F) & : \tau_C \to \tau_\perp, & \text{if } v : \tau \in \mathit{atr}(C), \tau \in \mathscr{T} \\
| & v(F) & : \tau_C \to \tau_D, & \text{if } v : D_{0,1} \in \mathit{atr}(C) \\
| & v(F) & : \tau_C \to 2^{\tau_D}, & \text{if } v : D_* \in \mathit{atr}(C) \\
| & f(F_1, \dots, F_n) & : \tau_1 \times \dots \times \tau_n \to \tau, & \text{if } f : \tau_1 \times \dots \times \tau_n \to \tau \\
| & \forall\, c \in F_1 \bullet F_2 & : \tau_C \times 2^{\tau_C} \times \mathbb{B}_\perp \to \mathbb{B}_\perp &
\end{array}
$$

- The formula above in prefix normal form: $\forall\, c \in \mathit{allInstances}_C \bullet \neq(\ x(c),\ \ 27\ )$

Semantics

$\bot \neq true$, $\bot \neq false$

- **Proto-OCL Types:**
  - $\mathcal{I}[\![\tau_C]\!] = \mathscr{D}(C) \mathbin{\dot\cup} \{\bot\}$, $\mathcal{I}[\![\tau_\bot]\!] = \mathscr{D}(\tau) \mathbin{\dot\cup} \{\bot\}$, $\mathcal{I}[\![2^{\tau_C}]\!] = \mathscr{D}(C_*) \mathbin{\dot\cup} \{\bot\}$
  - $\mathcal{I}[\![\mathbb{B}_\bot]\!] = \{\textit{true}, \textit{false}\} \mathbin{\dot\cup} \{\bot\}$, $\mathcal{I}[\![\mathbb{Z}_\bot]\!] = \mathbb{Z} \mathbin{\dot\cup} \{\bot\}$

- **Functions:**
  - We assume $f_\mathcal{I}$ given for each function symbol f ($\to$ in a minute).

- **Proto-OCL Semantics** (interpretation function):

sys. state

valuation of logical variables

$$\mathcal{I}[\![\,\cdot\,]\!](\,\cdot\,,\,\cdot\,) : \text{Proto-OCL-Formulae} \times \Sigma_\mathscr{S}^\mathscr{D} \times B \to \{\textit{true}, \textit{false}, \bot\}$$

  - $\mathcal{I}[\![c]\!](\sigma, \beta) = \beta(c)$ (assuming β is a type-consistent valuation of the logical variables),
  - $\mathcal{I}[\![allInstances_C]\!](\sigma, \beta) = \mathrm{dom}(\sigma) \cap \mathscr{D}(C)$,
  - $\mathcal{I}[\![v(F)]\!](\sigma, \beta) = \begin{cases} \big(\sigma(\mathcal{I}[\![F]\!](\sigma, \beta))\big)(v) & \text{, if } \mathcal{I}[\![F]\!](\sigma, \beta) \in \mathrm{dom}(\sigma) \\ \bot & \text{, otherwise} \end{cases}$ (if not $v : C_{0,1}$)
  - $\mathcal{I}[\![v(F)]\!](\sigma, \beta) = \begin{cases} u' & \text{, if } \mathcal{I}[\![F]\!](\sigma, \beta) \in \mathrm{dom}(\sigma) \text{ and } \sigma(\mathcal{I}[\![F]\!](\sigma, \beta))(v) = \{u'\} \\ \bot & \text{, otherwise} \end{cases}$ (if $v : C_{0,1}$)
  - $\mathcal{I}[\![f(F_1, \dots, F_n)]\!](\sigma, \beta) = f_\mathcal{I}(\mathcal{I}[\![F_1]\!](\sigma, \beta), \dots, \mathcal{I}[\![F_n]\!](\sigma, \beta))$,
  - $\mathcal{I}[\![\forall\, c \in F_1 \bullet F_2]\!](\sigma, \beta) = \begin{cases} \textit{true} & \text{, if } \mathcal{I}[\![F_2]\!](\sigma, \beta[c := u]) = \textit{true} \text{ for all } u \in \mathcal{I}[\![F_1]\!](\sigma, \beta) \\ \textit{false} & \text{, if } \mathcal{I}[\![F_2]\!](\sigma, \beta[c := u]) = \textit{false} \text{ for some } u \in \mathcal{I}[\![F_1]\!](\sigma, \beta) \\ \bot & \text{, otherwise} \end{cases}$

Semantics Cont'd

- Proto-OCL is a **three-valued** logic: a formula evaluates to *true*, *false*, or $\bot$.
- **Example**: $\wedge_\mathcal{I}(\cdot,\cdot) : \{\mathit{true}, \mathit{false}, \bot\} \times \{\mathit{true}, \mathit{false}, \bot\} \to \{\mathit{true}, \mathit{false}, \bot\}$ is defined as follows:

x_1	*true*	*true*	*true*	*false*	*false*	*false*	$\bot$	$\bot$	$\bot$
x_2	*true*	*false*	$\bot$	*true*	*false*	$\bot$	*true*	*false*	$\bot$
$\wedge_\mathcal{I}(x_1, x_2)$	*true*	*false*	$\bot$	*false*	*false*	*false*	$\bot$	*false*	$\bot$

  We assume common logical connectives $\neg, \wedge, \vee, \ldots$ with canonical 3-valued interpretation.

- **Example**: $+_\mathcal{I}(\cdot,\cdot) : (\mathbb{Z} \dot\cup \{\bot\}) \times (\mathbb{Z} \dot\cup \{\bot\}) \to \mathbb{Z} \dot\cup \{\bot\}$

$$+_\mathcal{I}(x_1, x_2) = \begin{cases} x_1 + x_2 & \text{, if } x_1 \neq \bot \text{ and } x_2 \neq \bot \\ \bot & \text{, otherwise} \end{cases}$$

  We assume common arithmetic operations $-, /, *, \ldots$
  and relation symbols $>, <, \leq, \ldots$ with **monotone** 3-valued interpretation.

- And we assume the special unary function symbol $isUndefined$:

$$isUndefined_\mathcal{I}(x) = \begin{cases} \mathit{true} & \text{, if } x = \bot, \\ \mathit{false} & \text{, otherwise} \end{cases}$$

  $isUndefined_\mathcal{I}$ is **definite**: it never yields $\bot$.

Example: Evaluate Formula for System State

σ : $\underline{1_C : \mathsf{C}}$ | $x = 13$

C
$x : Int$

$$\forall\, c \in allInstances_C \bullet x(c) \neq 27$$

- Recall **prefix notation**: $\forall\, c \in allInstances_C \bullet \neq(x(c), 27)$

  Note: $\neq$ is a binary function symbol, 27 is a 0-ary function symbol.

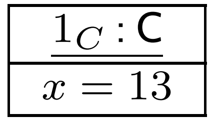
- **Example**:

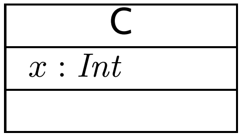
  $\mathcal{I}[\![\forall\, c \in allInstances_C \bullet \neq(x(c), 27)]\!](\sigma, \emptyset) =$ *true*, because…

  $\mathcal{I}[\![\neq(x(c), 27)]\!](\sigma, \beta)$, $\quad \beta := \emptyset[c := 1_C] = \{c \mapsto 1_C\}$

  $=$

Example: Evaluate Formula for System State

σ : 



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  $\mathcal{I}[\![\neq(x(c), 27)]\!](\sigma, \beta), \quad \beta := \emptyset[c := 1_C] = \{c \mapsto 1_C\}$

  $= \neq_{\mathcal{I}}(\, \mathcal{I}[\![x(c)]\!](\sigma, \beta),\ \mathcal{I}[\![27]\!](\sigma, \beta)\,)$

  $= \neq_{\mathcal{I}}(\, \sigma(\, \mathcal{I}[\![c]\!](\sigma, \beta)\,)(x),\ 27_{\mathcal{I}}\,)$

  $= \neq_{\mathcal{I}}(\, \underbrace{\sigma(\, \beta(c)\,)(x)}_{},\ 27_{\mathcal{I}}\,)$

  $= \quad (\sigma(1_C))(x)$

Example: Evaluate Formula for System State

σ : $\underline{1_C : \mathsf{C}}$ | $x = 13$

C
$x : Int$

$\underbrace{\forall\, c \in allInstances_C \bullet x(c) \neq 27}$

- Recall **prefix notation**: $\forall\, c \in allInstances_C \bullet \neq(x(c), 27)$

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- **Example**:

  $\mathcal{I}[\![\forall\, c \in allInstances_C \bullet \neq(x(c), 27)]\!](\sigma, \emptyset) =$ *true*, because...

  $\mathcal{I}[\![\neq(x(c), 27)]\!](\sigma, \beta), \quad \beta := \emptyset[c := 1_C] = \{c \mapsto 1_C\}$

  $= \neq_{\mathcal{I}}(\ \mathcal{I}[\![x(c)]\!](\sigma, \beta),\ \mathcal{I}[\![27]\!](\sigma, \beta)\ )$

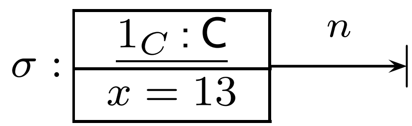
  $= \neq_{\mathcal{I}}(\ \sigma(\ \mathcal{I}[\![c]\!](\sigma, \beta)\ )(x),\ 27_{\mathcal{I}}\ )$

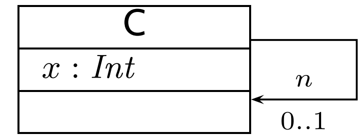
  $= \neq_{\mathcal{I}}(\ \sigma(\ \beta(c)\ )(x),\ 27_{\mathcal{I}}\ )$

  $= \neq_{\mathcal{I}}(\ \sigma(\ 1_C\ )(x),\ 27_{\mathcal{I}}\ )$

  $= \neq_{\mathcal{I}}(\ 13,\ 27\ ) =$ *true* ...and 1_C is the only C-object in σ: $\mathcal{I}[\![allInstances_C]\!](\sigma, \emptyset) = \{1_C\}$.

More Interesting Example





$\sigma = \{1_C \mapsto \{x \mapsto 13,\ n \mapsto \{\}\}\}$

$$\forall\, c \in allInstances_C \bullet x(n(c)) \neq 27$$

- Similar to the previous slide, we need the value of

$$\mathcal{I}[\![x(n(c))]\!](\sigma, \beta), \beta = \{c \mapsto 1_C\}$$

  - $\mathcal{I}[\![c]\!](\sigma, \beta) = \beta(c) = 1_C$
  - $\mathcal{I}[\![n(c)]\!](\sigma, \beta) = \bot$ since $\sigma(\ \mathcal{I}[\![c]\!](\sigma, \beta)\ )(n) = \emptyset \neq \{u'\}$ by rule

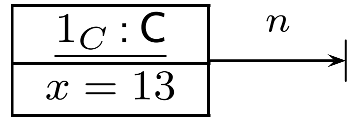
$$\mathcal{I}[\![v(F)]\!](\sigma, \beta) = \begin{cases} u' & \text{, if } \mathcal{I}[\![F]\!](\sigma, \beta) \in \mathrm{dom}(\sigma) \text{ and } \sigma(\mathcal{I}[\![F]\!](\sigma, \beta))(v) = \{u'\} \\ \bot & \text{, otherwise} \end{cases} \quad \text{(if } v : C_{0,1}\text{)}$$

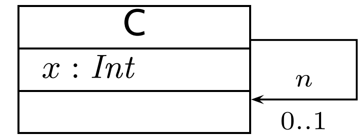
  - $\mathcal{I}[\![x(n(c))]\!](\sigma, \beta) = \bot$ since $\mathcal{I}[\![n(c)]\!](\sigma, \beta) = \bot$ by rule

$$\mathcal{I}[\![v(F)]\!](\sigma, \beta) = \begin{cases} \sigma\left(\mathcal{I}[\![F]\!](\sigma, \beta)\right)(v) & \text{, if } \mathcal{I}[\![F]\!](\sigma, \beta) \in \mathrm{dom}(\sigma) \\ \bot & \text{, otherwise} \end{cases} \quad \text{(if not } v : C_{0,1}\text{)}$$

More Interesting Example

σ :





$$\forall\, c \in allInstances_C \bullet x(n(c)) \neq 27$$

- Similar to the previous slide, we need the value of

$$\mathcal{I}[\![x(n(c))]\!](\sigma, \beta), \beta = \{c \mapsto 1_C\}$$

  - $\mathcal{I}[\![c]\!](\sigma, \beta) = \beta(c) = 1_C$
  - $\mathcal{I}[\![n(c)]\!](\sigma, \beta) = \bot$ since $\sigma(\,\mathcal{I}[\![c]\!](\sigma, \beta)\,)(n) = \emptyset \neq \{u'\}$ by rule

$$\mathcal{I}[\![v(F)]\!](\sigma, \beta) = \begin{cases} u' & \text{, if } \mathcal{I}[\![F]\!](\sigma, \beta) \in \mathrm{dom}(\sigma) \text{ and } \sigma(\mathcal{I}[\![F]\!](\sigma, \beta))(v) = \{u'\} \\ \bot & \text{, otherwise} \end{cases} \quad \text{(if } v : C_{0,1}\text{)}$$

  - $\mathcal{I}[\![x(n(c))]\!](\sigma, \beta) = \bot$ since $\mathcal{I}[\![n(c)]\!](\sigma, \beta) = \bot$ by rule

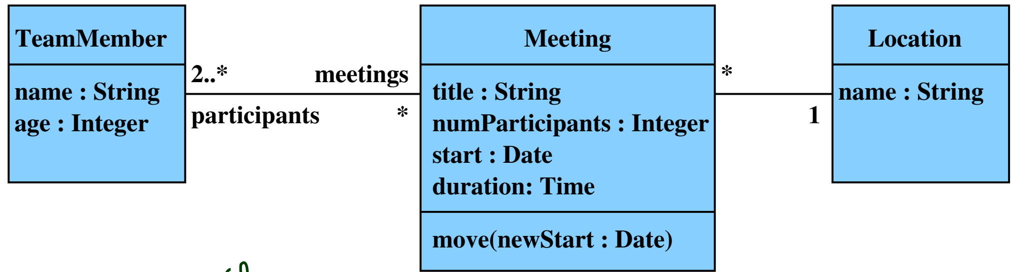
$$\mathcal{I}[\![v(F)]\!](\sigma, \beta) = \begin{cases} \sigma\left(\mathcal{I}[\![F]\!](\sigma, \beta)\right)(v) & \text{, if } \mathcal{I}[\![F]\!](\sigma, \beta) \in \mathrm{dom}(\sigma) \\ \bot & \text{, otherwise} \end{cases} \quad \text{(if not } v : C_{0,1}\text{)}$$

Object Constraint Language (OCL)

OCL is the same – just with less readable (?) syntax.

Literature: (OMG, 2006; Warmer and Kleppe, 1999).

Examples (from lecture “Softwaretechnik 2008”)



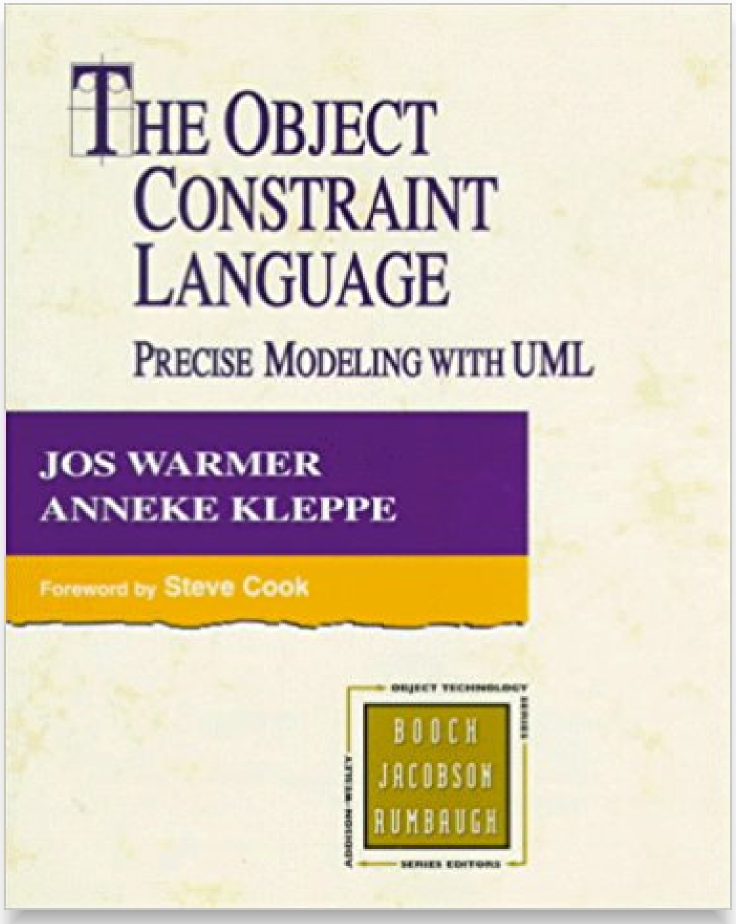
self :

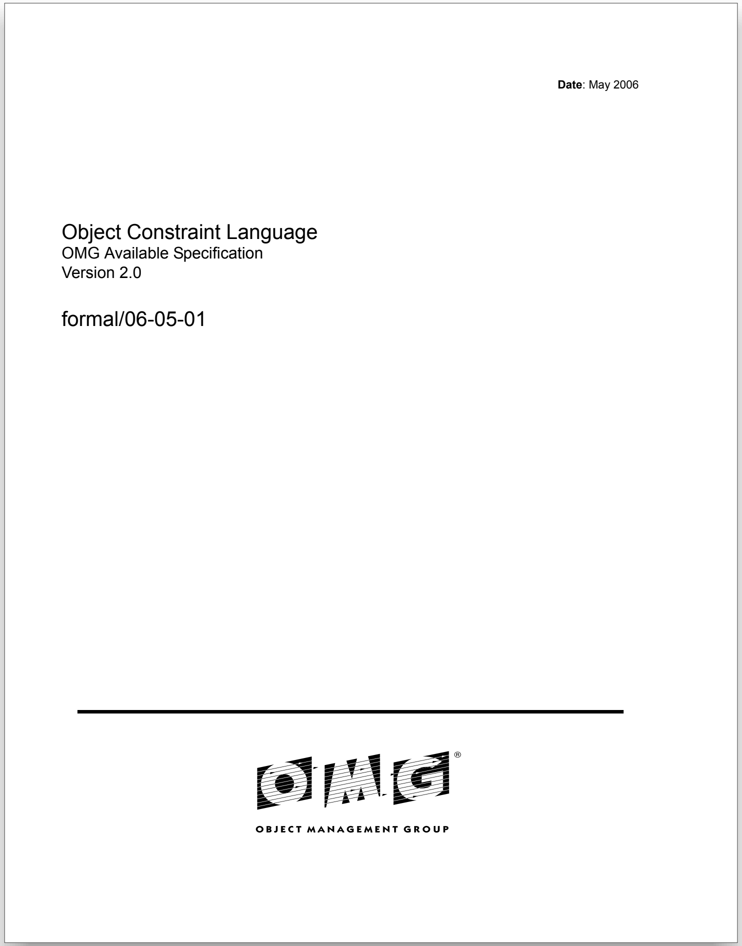
- **context** Meeting
  - **inv:** self.participants->size() = self. numParticipants
- **context** Location
  - **inv:** name="Lobby" **implies** meeting->isEmpty()

Prof. Dr. P. Thiemann, http://proglang.informatik.uni-freiburg.de/teaching/swt/2008/

$$\forall\, self \in allInstances_{Meeting} \bullet size(participants(self)) = numParticipants(self)$$

Literature




Date: May 2006

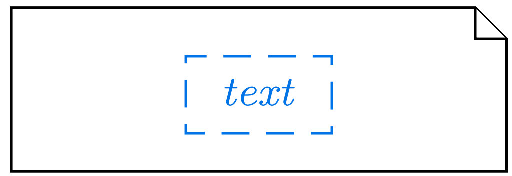
Object Constraint Language
OMG Available Specification
Version 2.0

formal/06-05-01

OBJECT MANAGEMENT GROUP

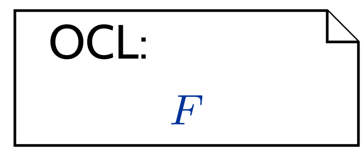
Where To Put OCL Constraints?

- **Notes**: A UML **note** is a diagram element of the form

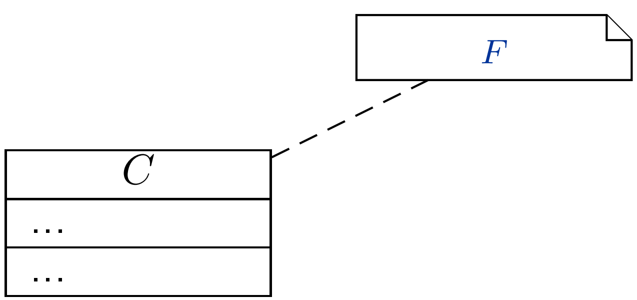


$text$ can principally be **everything**, in particular **comments** and **constraints**.

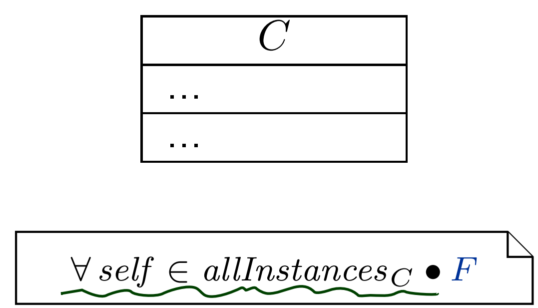
Sometimes, content is **explicitly classified** for clarity:



- Conventions:



stands for



Content

- **Object Diagrams Cont'd**
  - dangling references
  - partial vs. complete
  - object diagrams at work
- **Proto-OCL**
  - syntax, semantics
  - Proto-OCL vs. OCL
  - Putting It All Together: Proto-OCL vs. Software

Putting It All Together

Modelling Structure with Class Diagrams

Definition. Software is a finite description S of a (possibly infinite) set $[\![S]\!]$ of (finite or infinite) computation paths of the form $\sigma_0 \xrightarrow{\alpha_1} \sigma_1 \xrightarrow{\alpha_2} \sigma_2 \cdots$ where

- $\sigma_i \in \Sigma$, $i \in \mathbb{N}_0$, is called state (or configuration), and
- $\alpha_i \in A$, $i \in \mathbb{N}_0$, is called action (or event).

The (possibly partial) function $[\![\,\cdot\,]\!] : S \mapsto [\![S]\!]$ is called interpretation of S.

- The set of **states** Σ could be the set of **system states** as defined by a class diagram, e.g.

$$\Sigma := \Sigma_{\mathscr{S}}^{\mathscr{D}}$$

$\mathscr{S}$:

C
$x : Int$

- A corresponding **computation path** of a software S could be

$$\boxed{\begin{array}{c}\underline{27_C : \mathsf{C}}\\ x = 0\end{array}} \xrightarrow{\tau} \boxed{\begin{array}{c}\underline{27_C : \mathsf{C}}\\ x = 1\end{array}} \xrightarrow{\tau} \boxed{\begin{array}{c}\underline{27_C : \mathsf{C}}\\ x = 3\end{array}} \xrightarrow{\tau} \boxed{\begin{array}{c}\underline{27_C : \mathsf{C}}\\ x = 4\end{array}} \xrightarrow{\tau} \ldots$$

- If a requirement is formalised by the Proto-OCL constraint

$$F = \forall\, c \in allInstances_C \bullet x(c) < 4$$

then S **does not** satisfy the requirement.

More General: Software vs. Proto-OCL

- Let $\mathscr{S}$ be an **object system signature** and $\mathscr{D}$ a **structure**.
- Let S be a **software** with
  - states $\Sigma \subseteq \Sigma_{\mathscr{S}}^{\mathscr{D}}$, and
  - **computation paths** $[\![S]\!]$.
- Let F be a Proto-OCL constraint over $\mathscr{S}$.
- We say $[\![S]\!]$ **satisfies** F, denoted by $[\![S]\!] \models F$, if and only if for all

$$\sigma_0 \xrightarrow{\alpha_1} \sigma_1 \xrightarrow{\alpha_2} \sigma_2 \cdots \in [\![S]\!]$$

and all $i \in \mathbb{N}_0$,

$$\mathcal{I}[\![F]\!](\sigma_i, \emptyset) = \textit{true}.$$

- We say $[\![S]\!]$ **does not satisfy** F, denoted by $[\![S]\!] \not\models F$, if and only if there exists $\sigma_0 \xrightarrow{\alpha_1} \sigma_1 \xrightarrow{\alpha_2} \sigma_2 \cdots \in [\![S]\!]$ and $i \in \mathbb{N}_0$, such that $\mathcal{I}[\![F]\!](\sigma_i, \emptyset) = \textit{false}$.
- **Note**: $\neg([\![S]\!] \not\models F)$ does not imply $[\![S]\!] \models F$.

Topic Area Architecture & Design: Content

VL 10	• **Introduction and Vocabulary**
	• **Software Modelling**
⋮	– • model; views / viewpoints; 4+1 view
	• **Modelling structure**
VL 11	– • (simplified) Class & Object diagrams
⋮	– • (simplified) Object Constraint Logic (OCL)
VL 12	• **Principles of Design**
	– • modularity, separation of concerns
⋮	– • information hiding and data encapsulation
	– • abstract data types, object orientation
	• **Design Patterns**
VL 13	• **Modelling behaviour**
	– • Communicating Finite Automata (CFA)
⋮	– • Uppaal query language
	– • CFA vs. Software
VL 14	– • Unified Modelling Language (UML)
	– – • basic state-machines
⋮	– – • an outlook on hierarchical state-machines
	• **Model-driven/-based Software Engineering**

Tell Them What You've Told Them...

- **Class Diagrams** can be used to **graphically**
  - visualise code,
  - define an **object system structure** $\mathscr{S}$.
- An **Object System Structure** $\mathscr{S}$ (together with a structure $\mathscr{D}$)
  - defines a set of **system states** $\Sigma_{\mathscr{S}}^{\mathscr{D}}$.
- A **System State** $\sigma \in \Sigma_{\mathscr{S}}^{\mathscr{D}}$
  - can be **visualised** by an **object diagram**.
- **Proto-OCL** constraints can be evaluated on **system states**.
- A **software** over $\Sigma_{\mathscr{S}}^{\mathscr{D}}$ satisfies a Proto-OCL constraint F if and only if F evaluates to *true* in all system states of all the software's computation paths.

References

References

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OMG (2006). Object Constraint Language, version 2.0. Technical Report formal/06-05-01.

Schumann, M., Steinke, J., Deck, A., and Westphal, B. (2008). Traceviewer technical documentation, version 1.0. Technical report, Carl von Ossietzky Universität Oldenburg und OFFIS.

Warmer, J. and Kleppe, A. (1999). *The Object Constraint Language*. Addison-Wesley.