OwnerShip: invariants can depend on fields of other objects. For example, the invariant of list accesses node fields.

Callback: invariants can be temporarily violated. While invariant is violated we call a different method that calls back.

Atomicity: invariants can be temporarily violated. While invariant is violated another thread accesses object.
Temporarily Violating Invariants

```java
class Container {
    int[] content;
    int size;
    /*@ invariant 0 <= size && size <= content.length; @*/

    public void add(int v) {
        /* 1 */
        size++;
        /* 2 */
        if (size > content.length) {
            int[] newContent = new int[2*size+1];
            ...
            content = newContent;
        }
        ...
        /* 3 */
    }
}
```

When do Invariants Hold?

- Before a public method is called. /* 1 */
- After a public method returns. /* 3 */
- However, it may be violated in between. /* 2 */
Private Methods

```java
public class Container {
    int[] content;
    int size;
    /*@ invariant 0 <= size && size <= content.length; @*/

    private void growContent() {
        private /*@ helper @*/ void growContent() {
            ...
            content = newContent;
        }
    }

    public void add(int v) {
        /* invariant should hold */
        size++;
        /* invariant may be violated */
        if (size > content.length)
            growContent();
        ...
        /* invariant should hold, again */
    }
}
```

- Sometimes an invariant should not hold for a private method.
- JML has the keyword /*@ helper @*/.
Calling Methods of Other Classes

```java
public class Container {
    int[] content;
    int size;
    /*@ invariant 0 <= size && size <= content.length; @*/

    public void add(int v) {
        /* invariant should hold */
        size++;
        /* invariant may be violated */
        if (size > content.length) {
            newContent = new int[2*size+1];
            System.arraycopy(content, 0, newContent, 0, content.length);
            content = newContent;
        }

        ...  
        /* invariant should hold, again */
    }
}
```

- The invariant need not to hold, when calling other methods.
- However there is the callback problem.
The Callback Problem

```java
public class Log {
    public void log(String p) {
        logfile.write("Log: \""+p+"\" list\" is\""+Global.theList);
    }
}

public class Container {
    int[] content;
    int size;
    /*@ invariant 0 <= size && size <= content.length; @*/

    public void add(int v) {
        /* invariant should hold */
        size++;
        /* invariant may be violated */
        if (size > content.length) {
            Logger.log("growing\" array.\";
            ...
        }

    public String toString() {
        /* invariant should hold */
        ...
    }
}
```
The Callback Problem

- A method of a different class can be called while invariant is violated.
- This method may call a method of the first class.
- Who has to ensure that the invariant holds?
- jmlrac complains that invariant does not hold
- ESC/Java checks that most invariants hold at every method call, but not all invariants; this may lead to unsoundness.
A Ghost Variable for Invariants

Idea of David A. Naumann and Mike Barnett:

- Make the places where an invariant does not hold explicit.
- Add a ghost variable \textit{packed} that indicates if the invariant should hold.
- Before modifying an object set this variable to \textit{false}.
- When modification is finished, set it to \textit{true}.
- The following invariant should \textit{always} hold:
  \[
  \text{packed} \implies \text{invariants of object}
  \]
- The \textit{caller} has to ensure that the objects he uses are packed.
Example: A Ghost Variable for Invariants

```java
//@ public ghost boolean packed;
//@ private invariant packed ==> (size >= 0 && size <= content.length);

/*@ requires packed;
   @ ensures packed;
   @*/
public void add(int v) {
    unpack this;
    size++;
    ...
    pack this;
}
```

- The pre- and post-conditions explicitly states that invariant holds
- `unpack this` is an abbreviation for:
  ```java
  assert this.packed;
  set this.packed = false;
  ```
- `pack this` is an abbreviation for:
  ```java
  assert !this.packed;
  assert /*invariant of this holds*/;
  set this.packed = true;
  ```
The pack/unpack Mechanism

- An object must be unpacked before fields may be accessed.
- The invariant has to hold only while the object is packed.
- The invariant may only depend on fields of the object.
Static Checking with \textit{packed} ghost field:

- Fields may only be modified if \textit{packed} is false.
- For each \textit{pack} operation check that invariant holds again.
- Thus \textit{packed} $\Rightarrow$ \textit{invariants} holds for all states.
class TreeNode {
    int key, value;
    TreeNode left, right;
    /*@ invariant left != null ==> left.key <= key; @*/
    /*@ invariant right != null ==> right.key >= key; @*/

    public void add(Node n) {
        if (n.key < key) {
            if (left == null)
                left = n;
            else
                left.add(n);
        } else {
            ...
        }
    }
}
Adding Packed variable

class TreeNode {
    int key, value;
    TreeNode left, right;
    //@ public ghost boolean packed = false;

    /*@ invariant packed ==> (left != null ==> left.key <= key); @*/
    /*@ invariant packed ==> (right != null ==> right.key >= key); @*/

    //@ requires packed;
    //@ ensures packed;
    public void add(/*@non_null@*/ TreeNode n) {
        // unpack this
        if (n.key < key) {
            if (left == null)
                left = n;
            else
                left.add(n);
        } else {
            ...
        }
        // pack this
    }
}
Running ESC/Java gives:

> escjava2 -q TreeNode.java

TreeNode.java:19: Warning: Precondition possibly not established (Pre)
    left.add(n);

^{

Associated declaration is "TreeNode.java", line 9, col 8:
    //@ requires packed;

The nodes left and right must be packed!
class TreeNode {
    int key, value;
    TreeNode left, right;
    //@ public ghost boolean packed = false;

   /*@ invariant packed ==> (left != null ==> left.packed && left.key <= key); @*/
   /*@ invariant packed ==> (right != null ==> right.packed && right.key >= key); @*/

    //@ requires packed;
    //@ ensures packed;
    public void add(/*@ non_null @*/ TreeNode n) {
        ...
    }
}
Adding Ownership

There are still problems:

- The invariant also depends on fields of \textit{left} and \textit{right}. In particular the \textit{left.key} and \textit{left.packed}.
- Can \texttt{unpack this} violate the invariant of another TreeNode?
- How can we exclude undesired sharing, e.g., \texttt{left == this} or \texttt{left == n}?

Solution: Use the ownership principle
The owner must be unpacked before an owned object can be unpacked.

The invariant of owner may depend on owned objects.
How does pack/unpack work with ownership?

- To modify an object, you must **unpack** it first.
- To **unpack** an object, you must **unpack** the owner.
- To **pack** the owner again, its invariant must hold.

**unpack** `obj` is an abbreviation for:

```java
assert(obj.packed);
assert(obj.owner == null || !obj.owner.packed);
set obj.packed = false;
```

**pack** `obj` ensures that its owned classes are packed.

```java
assert(!obj.packed);
assert(left != null ==> (left.owner == this && left.packed));
assert(right != null ==> (right.owner == this && right.packed));
assert(/* other invariants of obj holds*/);
set obj.packed = true;
```
class TreeNode {
    int key, value;
    TreeNode left, right;
    //@ public ghost boolean packed = false;

    //@ invariant packed ==> (left != null ==> left.owner == this && left.packed && left.key <= key); @*/
    //@ invariant packed ==> (right != null ==> right.owner == this && right.packed && right.key >= key); @*/

    //@ requires packed && n.packed && n.owner == null;
    //@ ensures packed;
    public void add(//@non_null@*/ TreeNode n) {
        ...
    }
}
The ownership discipline has a few restrictions.

- An object invariant can only depend on fields of owned objects.
- An object can have at most one owner.
- A field may only be changed by the owner, or if the owner is unpacked.

Sometimes too restrictive!

Friendship offers another way to depend on other objects:

- An invariant can also depend on fields of granters.
- The class must define update guards for all fields it depends on.
- A granter object has a list of friends that depend on fields.
- A field may be changed if the update guards of all friends holds.
Friendship

Friendship is not symmetric. The allies are:

- Granter $G$ that gives rights to depend on a field.
- Friend $C$ whose invariant depends on a field.

Every class that changes a field of $G$ has to check the friend’s update condition.
Friend’s invariant can depend on granted fields.
Access to granted fields is checked against update guards.
A granter can have many friends.
All current friends must be checked.
The friend objects can be packed or unpacked.
class Object {
    /*@ spec_public */ int hashCode;
    //@ friend Map reads hashCode;
    //@ ghost JMLObjectSet deps;
}

class Map {
    JMLObjectSet buckets[];
    /*@ invariant
     \forall int i ; 0 <= i && i < buckets.length;
     (\forall Object o; buckets[i].has(o); o.deps.has(this) &&
      Math.abs(o.hashCode % buckets.length) == i); @*/

    /*@ guard obj.hashCode := val by
     val % buckets.length == obj.hashCode % buckets.length; @*/
}
Update Guard and Invariant

class FriendClass {
   //@ invariant friendInvariant(granter.field)
   //@ guard granter.field := val by updateGuardForField(granter, val);
}

The update guard must guarantee that the invariant is not invalidated:
friends.packed && friendInvariant(granter.field)
&& updateGuardForField(granter, val) ==> friendInvariant(val)
What May Appear in an Invariant

Only the following field accesses are allowed in an invariant:

- `this.field` for all fields.
- `x.field` if it appears in a subformula:
  \[ \texttt{forall Object x ; x.owner == this} \implies \ldots \]
- `object.field` if \texttt{object != null} \&\& \texttt{object.owner == this} can be proven.
- `x.field` if it appears in a subformula:
  \[ \texttt{forall Object x ; x.deps.has(this)} \implies \ldots \]
- `object.field` if \texttt{object != null} \&\& \texttt{object.deps.has(this)} can be proven.
A field access $obj.f = val$ only affects invariants of

- $obj$,
- $obj.owner$ if it is not null,
- and the objects in $obj.deps$.

$obj$ and $obj.owner$ must be unpacked if field is accessed. Thus their invariants need not to hold afterwards.

For the objects in $obj.deps$ the update guard must hold. Therefore, the invariant holds also with the new value $val$ for $obj.f$. 