

# Formal Methods for Java

## Lecture 25: Proving a JML-Program with KeY

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- Theorem Prover
- Developed at University of Karlsruhe
- <http://www.key-project.org/>.
- Theory specialized for Java(Card).
- Can generate proof-obligations from JML specification.
- Underlying theory: Sequent Calculus + Dynamic Logic

## Case Study: Euklid's Algorithm

Java code to compute gcd of non-negative numbers:

```
public static int gcd(int a, int b) {  
    while (a != 0 && b != 0) {  
        if (a > b)  
            a = a - b;  
        else  
            b = b - a;  
    }  
    return (a > b) ? a : b;  
}
```

Lets prove it with KeY-System.

We first need a specification.

## Definition (GCD)

Let  $a$  and  $b$  be natural numbers. A number  $d$  is the greatest common divisor (GCD) of  $a$  and  $b$  iff

- 1  $d|a$  and  $d|b$
- 2 If  $c|a$  and  $c|b$ , then  $c|d$ .

$d|a$  means  $d$  divides  $a$ .

$d|a : \Leftrightarrow \exists q. d * q = a$

# JML Specification

The specification can be converted to JML:

```
/*@
  @ requires a >= 0 &&& b >= 0;
  @ ensures \result >= 0;
  @ ensures (\exists int q; \result*q == a) &&&
  @         (\exists int q; \result*q == b) &&&
  @ (\forall int c;
  @   (\exists int q; c*q == a) &&& (\exists int q; c*q == b);
  @   (\exists int q; c*q == \result));
  @*/
public static int gcd(int a, int b)
```

## The rule `while_invariant_with_variant_dec`

The rule `while_invariant_with_variant_dec` takes an invariant  $inv$ , a modifies set  $\{m_1, \dots, m_k\}$  and a variant  $v$ . The following cases must be proven.

- Initially Valid:  $\implies inv \wedge v \geq 0$
- Body Preserves Invariant:

$$\begin{aligned} \implies \{m_1 := x_1 \parallel \dots \parallel m_k := x_k\} (inv \wedge [\{b = COND;\}] b = \mathbf{true}) \\ \rightarrow \langle BODY \rangle inv \end{aligned}$$

- Use Case:

$$\begin{aligned} \implies \{m_1 := x_1 \parallel \dots \parallel m_k := x_k\} (inv \wedge [\{b = COND;\}] b = \mathbf{false}) \\ \rightarrow \langle \dots \rangle \phi \end{aligned}$$

- Termination:

$$\begin{aligned} \implies \{m_1 := x_1 \parallel \dots \parallel m_k := x_k\} (inv \wedge v \geq 0 \wedge [\{b = COND;\}] b = \mathbf{true}) \\ \rightarrow \{old := v\} \langle BODY \rangle v \leq old \wedge v \geq 0 \end{aligned}$$

# Loop-Invariant

What is the loop invariant?

The algorithm changes  $a$  and  $b$ , but the gcd of  $a$  and  $b$  should stay the same.

In fact the set of common divisors of  $a$  and  $b$  never changes.

This suggests the following invariant:

$$\forall d. (d \mid \text{old}(a) \wedge d \mid \text{old}(b)) \leftrightarrow d \mid a \wedge d \mid b$$

In JML this can be specified as:

```
/*@ loop_invariant a >= 0 &&& b >= 0 &&&
  @   (\forall int d; true;
  @   (\exists int q; \old(a) == q*d)
  @   &&& (\exists int q; \old(b) == q*d)
  @   <==> (\exists int q; a == q*d) &&& (\exists int q; b == q*d)
  @   );
  @ assignable a, b;
  @ decreases a+b;
  @*/
```