

Software Design, Modelling and Analysis in UML

Lecture 02: Semantical Model

2012-10-24

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Contents & Goals

Last Lecture:

- Motivation: model-based development of things (houses, software) to cope with complexity, detect errors early
- Model-based (or -driven) Software Engineering
- UML Mode of the Lecture: Blueprint.

This Lecture:

- **Educational Objectives:** Capabilities for these tasks/questions:
 - Why is UML of the form it is?
 - Shall one feel bad if not using all diagrams during software development?
 - What is a signature, an object, a system state, etc.?
What's the purpose of signature, object, etc. in the course?
 - How do Basic Object System Signatures relate to UML class diagrams?

Content:

- Brief history of UML
- Course map revisited
- Basic Object System Signature, Structure, and System State

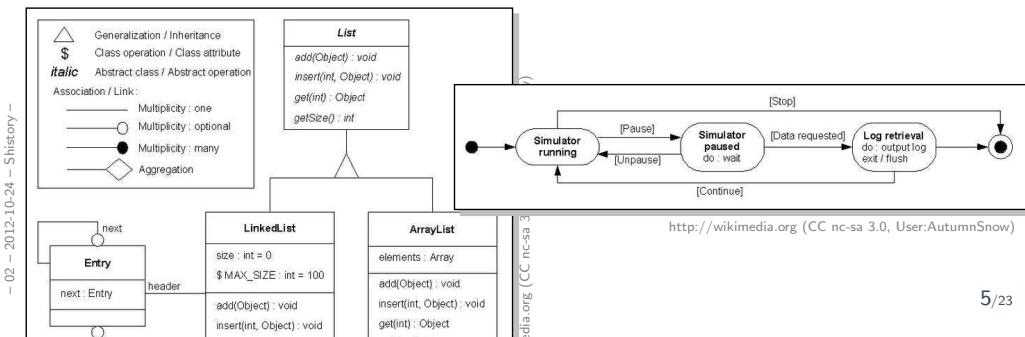
Why (of all things) UML?

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- Note: being a **modelling** languages doesn't mean being graphical (or: being a visual formalism [**Harel**]).
- For instance, [[Kastens and Büning, 2008](#)] also name:
 - Sets, Relations, Functions
 - Terms and Algebras
 - Propositional and Predicate Logic
 - Graphs
 - XML Schema, Entity Relation Diagrams, UML Class Diagrams
 - Finite Automata, Petri Nets, UML State Machines
- **Pro:** visual formalisms are found appealing and easier to **grasp**.
Yet they are not necessarily easier to **write!**
- **Beware:** you may meet people who dislike visual formalisms just for being graphical — maybe because it is easier to "trick" people with a meaningless picture than with a meaningless formula.
More serious: it's maybe easier to misunderstand a picture than a formula.

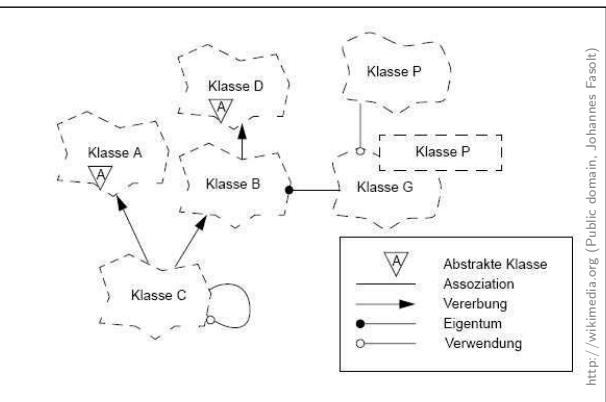
A Brief History of UML

- Boxes/lines and finite automata are used to visualise software **for ages**.
- **1970's, Software Crisis™**
 - Idea: learn from engineering disciplines to handle growing complexity.
 - Languages: **Flowcharts, Nassi-Shneiderman, Entity-Relation Diagrams**
- Mid **1980's**: **Statecharts** [Harel, 1987], **StateMate™** [Harel et al., 1990]
- Early **1990's**, advent of **Object-Oriented**-Analysis/Design/Programming
 - Inflation of notations and methods, most prominent:
 - **Object-Modeling Technique (OMT)** [Rumbaugh et al., 1990]



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 - **Booch Method and Notation** [Booch, 1993]



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 - **Booch Method and Notation** [Booch, 1993]
 - **Object-Oriented Software Engineering** (OOSE) [Jacobson et al., 1992]
- Each “persuasion” selling books, tools, seminars...
- Late **1990's**: joint effort **UML 0.x, 1.x**
 - Standards published by **Object Management Group** (OMG), “*international, open membership, not-for-profit computer industry consortium*”.
- Since **2005**: **UML 2.x**

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UML Overview [OMG, 2007b, 684]

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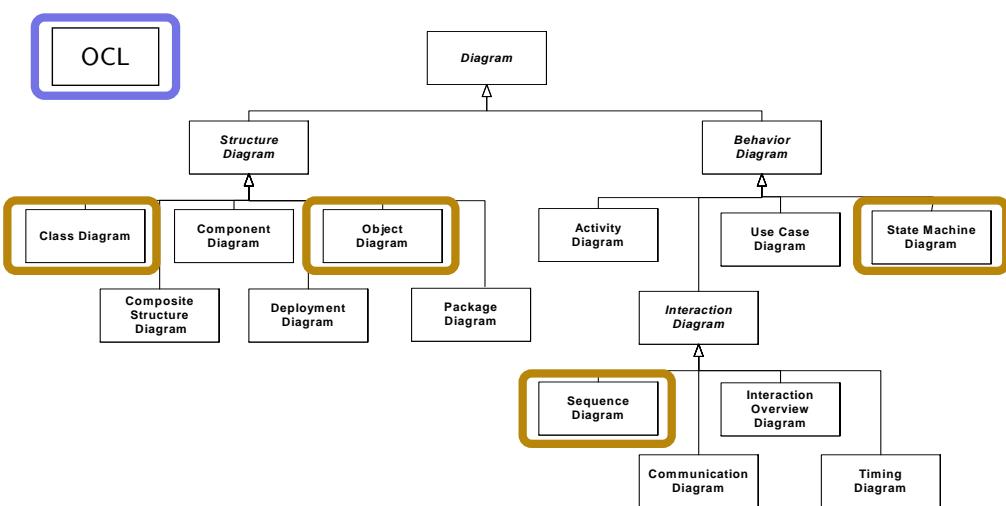


Figure A.5 - The taxonomy of structure and behavior diagram

[Dobing and Parsons, 2006]

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Common Expectations on UML

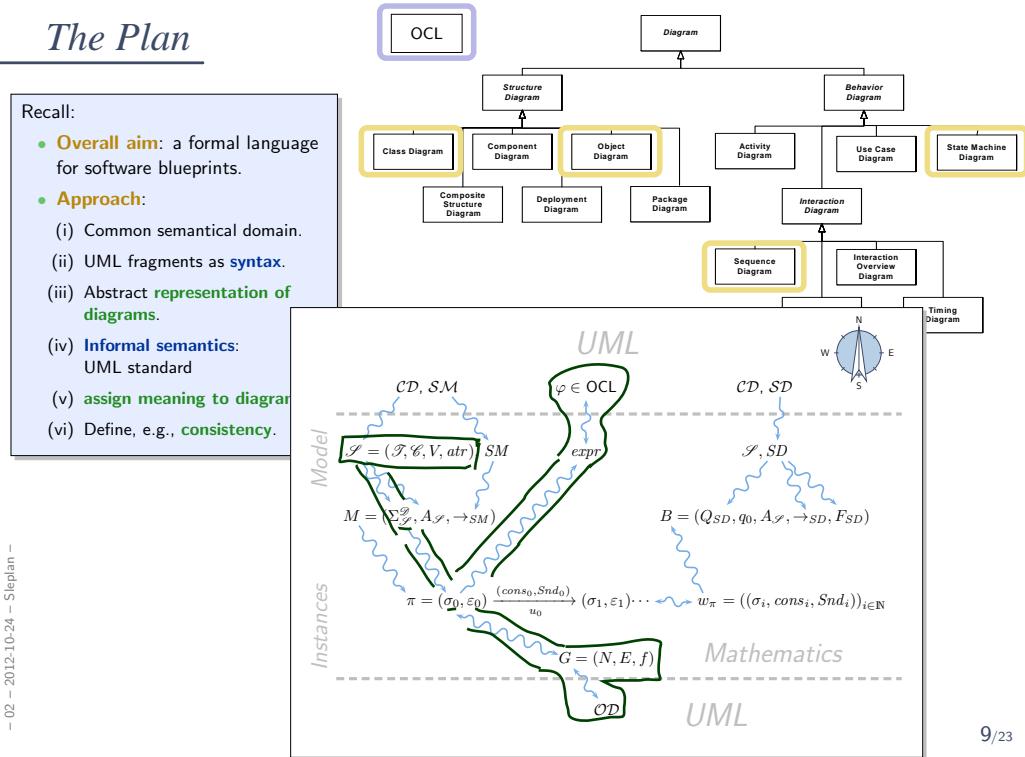
- Easily writeable, readable even by customers
- Powerful enough to bridge the gap between idea and implementation
- Means to tame complexity by separation of concerns ("views")
- Unambiguous
- Standardised, exchangeable between modelling tools
- UML standard says how to develop software
- Using UML leads to better software
- ...

We will see...

Seriously: After the course, you should have an own opinion on each of these claims.
In how far/in what sense does it hold? Why? Why not? How can it be achieved?
Which ones are really only hopes and expectations? ...?

Course Map Revisited

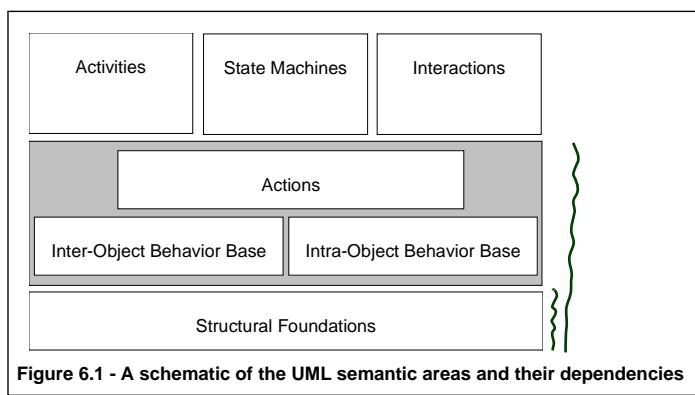
The Plan



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UML: Semantic Areas



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[OMG, 2007b, 11]

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Common Semantical Domain

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Basic Object System Signature

Definition. A (Basic) Object System **Signature** is a quadruple

$$\mathcal{S} = (\mathcal{T}, \mathcal{C}, V, atr)$$

where

- \mathcal{T} is a set of (basic) types,
 - \mathcal{C} is a finite set of classes,
 - V is a finite set of typed attributes, i.e., each $v \in V$ has type
 - $\tau \in \mathcal{T}$ or
 - $C_{0,1}$ or C_* , where $C \in \mathcal{C}$

(written $v : \tau$ or $v : C_{0,1}$ or $v : C_*$),
 - $atr : \mathcal{C} \rightarrow \bigcup_{\tau \in \mathcal{T}} V$ maps each class to its set of attributes.

Note: Inspired by OCL 2.0 standard [OMG, 2006], Annex A.

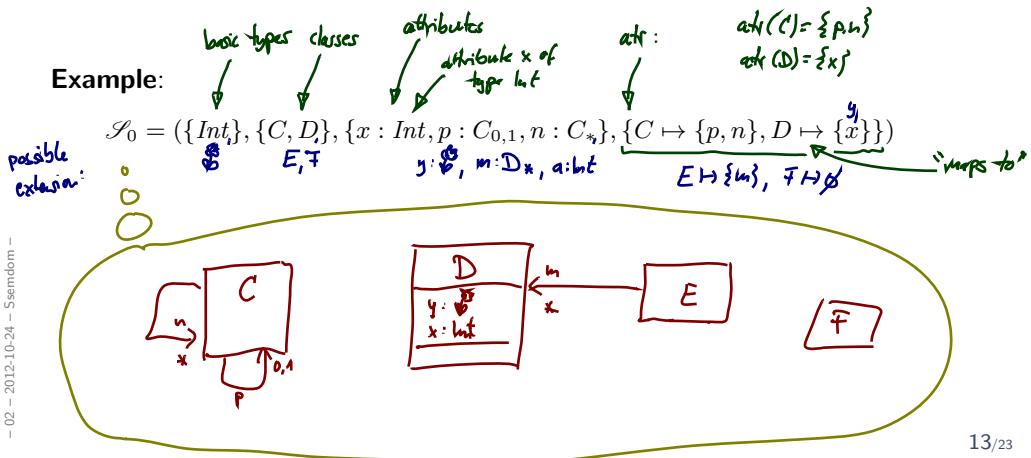
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Basic Object System Signature Example

$\mathcal{S} = (\mathcal{T}, \mathcal{C}, V, atr)$ where

- (basic) types \mathcal{T} and classes \mathcal{C} , (both finite),
- typed attributes V, τ from \mathcal{T} or $C_{0,1}$ or C_* , $C \in \mathcal{C}$,
- $atr : \mathcal{C} \rightarrow 2^V$ mapping classes to attributes.



Basic Object System Signature Another Example

$\mathcal{S} = (\mathcal{T}, \mathcal{C}, V, atr)$ where

- (basic) types \mathcal{T} and classes \mathcal{C} , (both finite),
- typed attributes V, τ from \mathcal{T} or $C_{0,1}$ or C_* , $C \in \mathcal{C}$,
- $atr : \mathcal{C} \rightarrow 2^V$ mapping classes to attributes.

Example:

$$\Psi = \left(\{\text{Int}, \text{Float}\}, \{C, D\}, \{ \begin{array}{l} C : x : \text{Int}, \\ D : x : \text{Float}, y : \text{Int} \end{array} \} \right)$$

(name of an attribute, could also be $C : x$)

Q: What about
 • class C with attribute x:Int
 • class D with attribute x:Float ?

A: Rename consistently.

Basic Object System Structure

Definition. A Basic Object System Structure of

$$\mathcal{S} = (\mathcal{T}, \mathcal{C}, V, atr)$$

is a domain function \mathcal{D} which assigns to each type a domain, i.e.

- $\tau \in \mathcal{T}$ is mapped to $\mathcal{D}(\tau)$,
 - $C \in \mathcal{C}$ is mapped to an infinite set $\mathcal{D}(C)$ of (object) identities.
- Note: Object identities only have the “=” operation;
object identities of different classes are disjoint, i.e.
 $\forall C, D \in \mathcal{C} : C \neq D \rightarrow \mathcal{D}(C) \cap \mathcal{D}(D) = \emptyset$.
- C_* and $C_{0,1}$ for $C \in \mathcal{C}$ are mapped to $2^{\mathcal{D}(C)}$.

We use $\mathcal{D}(\mathcal{C})$ to denote $\bigcup_{C \in \mathcal{C}} \mathcal{D}(C)$; analogously $\mathcal{D}(\mathcal{C}_*)$.

Note: We identify objects and object identities, because both uniquely determine each other (cf. OCL 2.0 standard).

Basic Object System Structure Example

Wanted: a structure for signature

$$\mathcal{S}_0 = (\{Int\}, \{C, D\}, \{x : Int, p : C_{0,1}, n : C_*\}, \{C \mapsto \{p, n\}, D \mapsto \{x\}\})$$

Recall: by definition, seek a \mathcal{D} which maps

- $\tau \in \mathcal{T}$ to **some** $\mathcal{D}(\tau)$,
- $c \in \mathcal{C}$ to **some** identities $\mathcal{D}(C)$ (infinite, disjoint for different classes),
- C_* and $C_{0,1}$ for $C \in \mathcal{C}$ to $\mathcal{D}(C_{0,1}) = \mathcal{D}(C_*) = 2^{\mathcal{D}(C)}$.

$$\begin{aligned}
 \mathcal{D}(Int) &= \mathbb{Z} && \text{also valid: } \mathcal{D}_2 \\
 \mathcal{D}(C) &= \mathbb{N}^+ \times \{C\} \simeq \{1_C, 2_C, 3_C, \dots\} && = \{1, 3, 5, \dots\} \\
 \mathcal{D}(D) &= \mathbb{N}^+ \times \{D\} \simeq \{1_D, 2_D, 3_D, \dots\} && \rightarrow \{2, 4, 6, \dots\} \\
 \mathcal{D}(C_{0,1}) = \mathcal{D}(C_*) &= 2^{\mathcal{D}(C)} && \\
 \mathcal{D}(D_{0,1}) = \mathcal{D}(D_*) &= 2^{\mathcal{D}(D)} &&
 \end{aligned}$$

System State

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Definition. Let \mathcal{D} be a structure of $\mathcal{S} = (\mathcal{T}, \mathcal{C}, V, \text{attr})$. A **system state** of \mathcal{S} wrt. \mathcal{D} is a **type-consistent** mapping

$$\sigma : \mathcal{D}(\mathcal{C}) \rightarrow (V \rightarrow (\mathcal{D}(\mathcal{T}) \cup \mathcal{D}(\mathcal{C}_*))).$$

That is, for each $u \in \mathcal{D}(C)$, $C \in \mathcal{C}$, if $u \in \text{dom}(\sigma)$

- $\text{dom}(\sigma(u)) = \text{attr}(C)$
- $(\sigma(u))(v) \in \mathcal{D}(\tau)$ if $v : \tau, \tau \in \mathcal{T}$
- $(\sigma(u))(v) \in \mathcal{D}(D_*)$ if $v : D_{0,1}$ or $v : D_*$ with $D \in \mathcal{C}$

We call $u \in \mathcal{D}(\mathcal{C})$ **alive** in σ if and only if $u \in \text{dom}(\sigma)$.

We use $\Sigma_{\mathcal{S}}^{\mathcal{D}}$ to denote the set of all system states of \mathcal{S} wrt. \mathcal{D} .

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System State Example

Signature, Structure:

$$\mathcal{S}_0 = (\{Int\}, \{C, D\}, \{x : Int, p : C_{0,1}, n : C_*\}, \{C \mapsto \{p, n\}, D \mapsto \{x\}\})$$

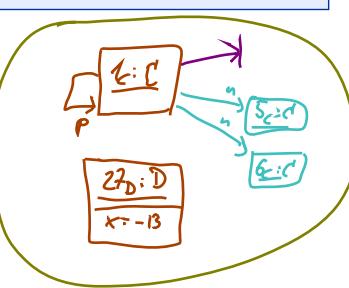
$$\mathcal{D}(Int) = \mathbb{Z}, \quad \mathcal{D}(C) = \{1_C, 2_C, 3_C, \dots\}, \quad \mathcal{D}(D) = \{1_D, 2_D, 3_D, \dots\}$$

Wanted: $\sigma : \mathcal{D}(\mathcal{C}) \rightarrow (V \rightarrow (\mathcal{D}(\mathcal{T}) \cup \mathcal{D}(\mathcal{C}_*)))$ such that

- $\text{dom}(\sigma(u)) = \text{attr}(C)$,
- $\sigma(u)(v) \in \mathcal{D}(\tau)$ if $v : \tau, \tau \in \mathcal{T}$,
- $\sigma(u)(v) \in \mathcal{D}(C_*)$ if $v : D_*$ with $D \in \mathcal{C}$.

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- $\sigma_1 = \emptyset$ *empty function* *or (p pointing nowhere):* $p \mapsto \emptyset$
 - $\sigma_2 = \{1_C \mapsto \{p \mapsto \{1_C\}, n \mapsto \{5_C, 6_C\}\}, 2_D \mapsto \{x \mapsto -13\}\}$ *OK*
 - $\sigma_3 = \{1_D \mapsto \emptyset\} \times$
 - $\sigma_4 = \{1_D \mapsto \{p \mapsto \emptyset, x \mapsto 3\}\} \times$
- Wsb. \mathcal{D}_2 :
- $\sigma_3 = \{1 \mapsto \{p \mapsto \{1\}, n \mapsto \{5, 7\}, 10 \mapsto \{x \mapsto \text{red}\}\}$



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System State Example

Signature, Structure:

$$\mathcal{S}_0 = (\{\text{Int}\}, \{C, D\}, \{x : \text{Int}, p : C_{0,1}, n : C_*\}, \{C \mapsto \{p, n\}, D \mapsto \{x\}\})$$

$$\mathcal{D}(\text{Int}) = \mathbb{Z}, \quad \mathcal{D}(C) = \{1_C, 2_C, 3_C, \dots\}, \quad \mathcal{D}(D) = \{1_D, 2_D, 3_D, \dots\}$$

Wanted: $\sigma : \mathcal{D}(\mathcal{C}) \nrightarrow (V \nrightarrow (\mathcal{D}(\mathcal{T}) \cup \mathcal{D}(\mathcal{C}_*)))$ such that

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- $\sigma(u)(v) \in \mathcal{D}(C_*)$ if $v : D_*$ with $D \in \mathcal{C}$.

• Concrete, explicit:

$$\sigma = \{1_C \mapsto \{p \mapsto \emptyset, n \mapsto \{5_C\}\}, 5_C \mapsto \underbrace{\{p \mapsto \emptyset, n \mapsto \emptyset\}}, 1_D \mapsto \{x \mapsto 23\}\}.$$

$$\text{dom}(\sigma(5_C)) = \{p, n\} = \text{atr}(C) \checkmark$$

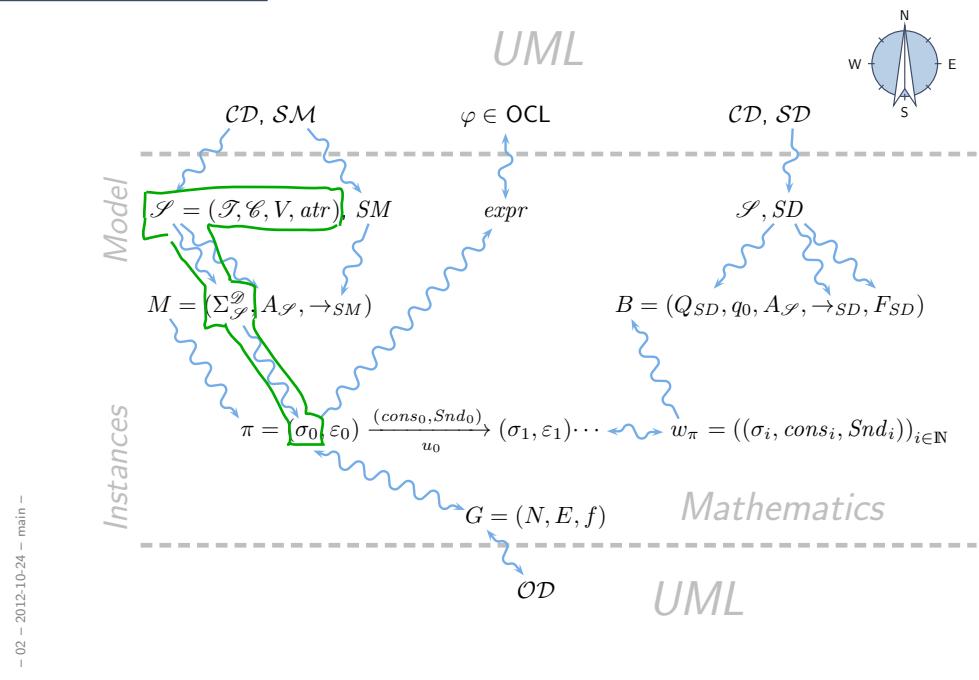
• Alternative: symbolic system state

$$\sigma = \{c_1 \mapsto \{p \mapsto \emptyset, n \mapsto \{c_2\}\}, c_2 \mapsto \{p \mapsto \emptyset, n \mapsto \emptyset\}, d \mapsto \{x \mapsto 23\}\}$$

assuming $c_1, c_2 \in \mathcal{D}(C), d \in \mathcal{D}(D), c_1 \neq c_2$.

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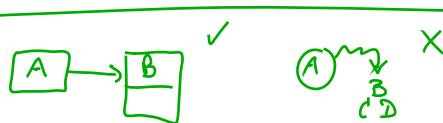
Course Map



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References



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