Software Design, Modelling and Analysis in UML

Lecture 13: Core State Machines III

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Contents & Goals

Last Lecture:

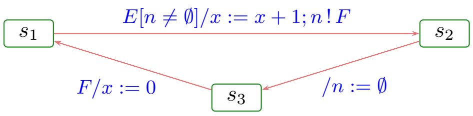
- Ether
- System configuration

This Lecture:

- **Educational Objectives:** Capabilities for following tasks/questions.
  - What does this State Machine mean? What happens if I inject this event?
  - Can you please model the following behaviour.
  - What is: Signal, Event, Ether, Transformer, Step, RTC.
- **Content:**
  - Transformer
  - Examples for transformer
  - Run-to-completion Step
  - Putting It All Together

System Configuration, Ether, Transformer

Where are we?



- **Wanted**: a labelled transition relation

$$(\sigma, \varepsilon) \xrightarrow[u_x]{(cons, Snd)} (\sigma', \varepsilon')$$

  on system configuration, labelled with the **consumed** and **sent** events, (σ', ε') being the result (or effect) of **one object** u_x taking a transition of **its** state machine from the current state mach. state $\sigma(u_x)(st_C)$.

- **Have**: system configuration (σ, ε) comprising current state machine state and stability flag for each object, and the ether.

- **Plan**:
  - (i) Introduce **transformer** as the semantics of action annotions. **Intuitively**, (σ', ε') is the effect of applying the transformer of the taken transition.
  - (ii) Explain how to choose transitions depending on ε and when to stop taking transitions — the **run-to-completion "algorithm"**.

Transformer

Definition.
Let $\Sigma_{\mathscr{S}}^{\mathscr{D}}$ the set of system configurations over some $\mathscr{S}_0$, $\mathscr{D}_0$, Eth.

We call a relation

$$t \subseteq \mathscr{D}(\mathscr{C}) \times (\Sigma_{\mathscr{S}}^{\mathscr{D}} \times Eth) \times (\Sigma_{\mathscr{S}}^{\mathscr{D}} \times Eth)$$

a (system configuration) **transformer**.

- In the following, we assume that each application of a transformer t to some system configuration (σ, ε) for object u_x is associated with a set of **observations**

$$Obs_t[u_x](\sigma, \varepsilon) \in 2^{\mathscr{D}(\mathscr{C}) \times \mathscr{D}(\mathscr{E}) \times Evs(\mathscr{E} \,\dot{\cup}\, \{*,+\}, \mathscr{D}) \times \mathscr{D}(\mathscr{C})}.$$

- An observation $(u_{src}, u_e, (E, \vec{d}), u_{dst}) \in Obs_t[u_x](\sigma, \varepsilon)$ represents the information that, as a "side effect" of u_x executing t, an event (!) $(E, \vec{d})$ has been sent from u_{src} to u_{dst}.

  Special cases: creation/destruction.

Why Transformers?

- **Recall** the (simplified) syntax of transition annotations:

$$annot ::= [\ \ \langle event \rangle\ \ [\ \text{'['}\ \langle guard \rangle\ \text{']'}\ ]\ \ [\ \text{'/'}\ \langle action \rangle]\ \ ]$$

- **Clear**: $\langle event \rangle$ is from $\mathscr{E}$ of the corresponding signature.
- **But:** What are $\langle guard \rangle$ and $\langle action \rangle$?
  - UML can be viewed as being **parameterized** in **expression language** (providing $\langle guard \rangle$) and **action language** (providing $\langle action \rangle$).
  - **Examples**:
    - **Expression Language**:
      - OCL
      - Java, C++, ... expressions
      - ...
    - **Action Language**:
      - UML Action Semantics, "Executable UML"
      - Java, C++, ... statements (plus some event send action)
      - ...

Transformers as Abstract Actions!

In the following, we assume that we're **given**

- an **expression language** $Expr$ for guards, and
- an **action language** Act for actions,

and that we're **given**

- a **semantics** for boolean expressions in form of a partial function

$$I[\![\,\cdot\,]\!](\cdot\,,\,\cdot\,) : Expr \to (\Sigma_{\mathscr{S}}^{\mathscr{D}} \times \mathscr{D}(\mathscr{C}) \nrightarrow \mathbb{B})$$

which evaluates expressions in a given system configuration,

Assuming I to be partial is a way to treat "undefined" during runtime. If I is not defined (for instance because of dangling-reference navigation or division-by-zero), we want to go to a designated "error" system configuration.

- a **transformer** for each action: for each $act \in Act$, we assume to have

$$t_{act} \subseteq \mathscr{D}(\mathscr{C}) \times (\Sigma_{\mathscr{S}}^{\mathscr{D}} \times Eth) \times (\Sigma_{\mathscr{S}}^{\mathscr{D}} \times Eth)$$

example OCL: $I[\![expr]\!](\sigma, u) :=$ true, if $I[\![expr]\!](\sigma, \{self \mapsto u\}) = true$; false, if $I[\![expr]\!](\sigma, \{self \mapsto u\}) = false$; and undefined otherwise

OCL interpretation

the same

object id to evaluate for

Expression/Action Language Examples

We can make the assumptions from the previous slide because **instances exist**:

- for OCL, we have the OCL semantics from Lecture 03. Simply remove the pre-images which map to "⊥".
- for Java, the operational semantics of the SWT lecture uniquely defines transformers for sequences of Java statements.

We distinguish the following kinds of transformers:

- **skip**: do nothing — recall: this is the default action   only skip
- **send**: modifies ε — interesting, because state machines are built around sending/consuming events   e.g. n!F
- **create**/**destroy**: modify domain of σ — not specific to state machines, but let's discuss them here as we're at it   e.g. new C, delete n
- **update**: modify own or other objects' local state — boring   e.g. x := x+1

Action Language

In the following we consider

$Act_{\mathscr{S}} := \{skip\}$
$\cup \{update(expr_1, v, expr_2) \mid expr_1, expr_2 \in OCLExpr, v \in V\}$
$\cup \{send(expr_1, E, expr_2) \mid expr_1, expr_2 \in OCLExpr, E \in \mathscr{E}\}$
$\cup \{create(C, expr, v) \mid C \in \mathscr{C}, expr \in OCLExpr, v \in V\}$
$\cup \{destroy(expr) \mid expr \in OCLExpr\}$

$Expr_{\mathscr{S}}$: OCL expressions over $\mathscr{S}$

Transformer Examples: Presentation

abstract syntax — op $update(e_1, v, e_2)$ — **concrete syntax** $e_1.v := e_2$

intuitive semantics

...

well-typedness

...

semantics

$((\sigma, \varepsilon), (\sigma', \varepsilon')) \in t_{\text{op}}[u_x]$ iff ...

or

$t_{\text{op}}[u_x](\sigma, \varepsilon) = \{(\sigma', \varepsilon')\}$ where ...

observables

$Obs_{\text{op}}[u_x] = \{\dots\}$, not a relation, depends on choice

(error) conditions

Not defined if ...

Transformer: Skip

abstract syntax — skip
concrete syntax — *skip*

intuitive semantics

do nothing

well-typedness

./.

semantics

$$t[u_x](\sigma, \varepsilon) = \{(\sigma, \varepsilon)\}$$

"if u_x executes skip on (σ, ε), then the result is (σ, ε)"

observables

$$Obs_{\texttt{skip}}[u_x](\sigma, \varepsilon) = \emptyset$$

(error) conditions

Transformer: Update

abstract syntax — $\texttt{update}(expr_1, v, expr_2)$
concrete syntax — $expr_1 . v := expr_2$

intuitive semantics

Update attribute v in the object denoted by $expr_1$ to the value denoted by $expr_2$.

well-typedness

$expr_1 : \tau_C$ and $v : \tau \in atr(C)$; $expr_2 : \tau$;
$expr_1, expr_2$ obey visibility and navigability

semantics

$$t_{\texttt{update}(expr_1, v, expr_2)}[u_x](\sigma, \varepsilon) = \{(\sigma', \varepsilon)\}$$

where $\sigma' = \sigma[u \mapsto \sigma(u)[v \mapsto I[\![expr_2]\!](\sigma, u_x)]]$ with
$u = I[\![expr_1]\!](\sigma, u_x)$

either does not change value denoted by $expr_2$ in σ for object u

change local state of object u

change value of v in $\sigma(u)$

object denoted by $expr_1$ (relative to u_x)

observables

$$Obs_{\texttt{update}(expr_1, v, expr_2)}[u_x] = \emptyset$$

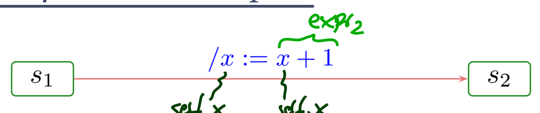
(error) conditions

Not defined if $I[\![expr_1]\!](\sigma, u_x)$ or $I[\![expr_2]\!](\sigma, u_x)$ not defined.

i.e. $t_{update(e_1, v, e_2)}[u_x](\sigma, \varepsilon) = \emptyset$

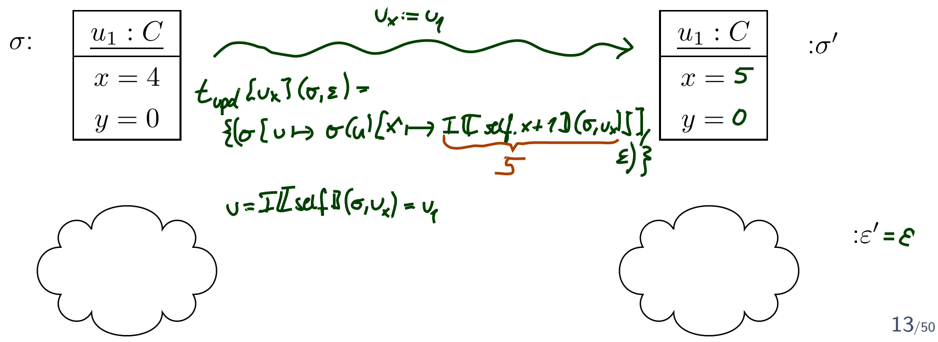
Update Transformer Example

$\mathcal{SM}_C$:



update$(expr_1, v, expr_2)$

$$t_{\mathtt{update}(expr_1, v, expr_2)}[u_x](\sigma, \varepsilon) = (\sigma[u \mapsto \sigma(u)[v \mapsto I[\![expr_2]\!](\sigma, \beta)]], \varepsilon),$$
$$u = I[\![expr_1]\!](\sigma, \beta)$$



Transformer: Send

abstract syntax | **concrete syntax**

send$(E(expr_1, ..., expr_n), expr_{dst})$

intuitive semantics

Object $u_x : C$ sends event E to object $expr_{dst}$, i.e. create a fresh signal instance, fill in its attributes, and place it in the ether.

well-typedness

$$expr_{dst} : \tau_D,\ C, D \in \mathscr{C} \setminus \mathscr{E};\ E \in \mathscr{E};$$
$$atr(E) = \{v_1 : \tau_1, \ldots, v_n : \tau_n\};\ expr_i : \tau_i,\ 1 \leq i \leq n;$$

all expressions obey visibility and navigability in C

semantics

$$t_{\mathtt{send}(E(expr_1, \ldots, expr_n), expr_{dst})}[u_x](\sigma, \varepsilon) = (\sigma', \varepsilon')$$

where $\sigma' = \sigma \dot\cup \{u \mapsto \{v_i \mapsto d_i \mid 1 \leq i \leq n\}\}$; $\varepsilon' = \varepsilon \oplus (u_{dst}, u)$; if $u_{dst} = I[\![expr_{dst}]\!](\sigma, \beta) \in \text{dom}(\sigma)$; $d_i = I[\![expr_i]\!](\sigma, \beta)$ for $1 \leq i \leq n$;

$u \in \mathscr{D}(E)$ a fresh identity, i.e. $u \notin \text{dom}(\sigma)$,

and where $(\sigma', \varepsilon') = (\sigma, \varepsilon)$ if $u_{dst} \notin \text{dom}(\sigma)$

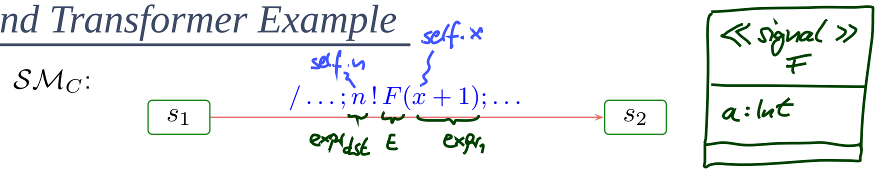
observables

$$Obs_{\mathtt{send}}[u_x] = \{(u_x, u, (E, d_1, \ldots, d_n), u_{dst})\}$$

(error) conditions

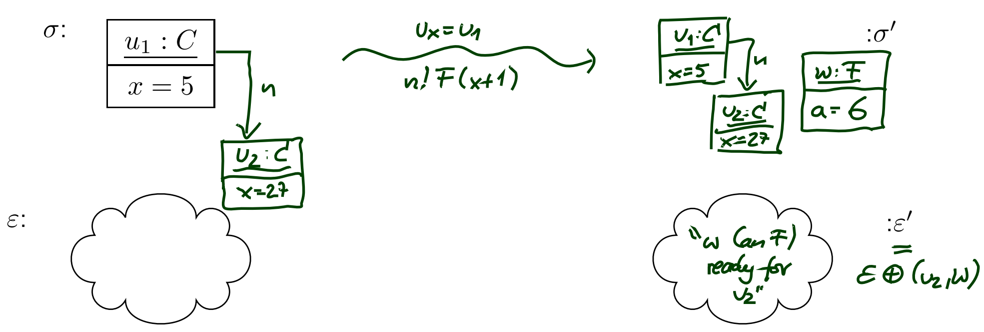
$I[\![expr]\!](\sigma, \beta)$ not defined for any $expr \in \{expr_{dst}, expr_1, \ldots, expr_n\}$

Send Transformer Example



$$\texttt{send}(E(expr_1, ..., expr_n), expr_{dst})$$

$$t_{\texttt{send}(expr_{src}, E(expr_1,...,expr_n), expr_{dst})}[u_x](\sigma, \varepsilon) = ...$$



References

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