

Software Design, Modelling and Analysis in UML

Lecture 15: ~~Hierarchical State Machines I~~ State Machines V

2014-01-13

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Contents & Goals

Last Lecture:

- RTC-Rules: Discard, Dispatch, Commence.

This Lecture:

- **Educational Objectives:** Capabilities for following tasks/questions.
 - What does this State Machine mean? What happens if I inject this event?
 - Can you please model the following behaviour.
 - What is: initial state.
 - What does this **hierarchical** State Machine mean? What **may happen** if I inject this event?
 - What is: AND-State, OR-State, pseudo-state, entry/exit/do, final state, . . .
- **Content:**
 - Step, RTC, Divergence
 - Putting It All Together
 - Rhapsody Demo
 - Hierarchical State Machines Syntax

Step and Run-to-completion Step

Notions of Steps: The Step

Note: we call one evolution $(\sigma, \varepsilon) \xrightarrow[u]{(cons, Snd)} (\sigma', \varepsilon')$ a **step**.

Thus in our setting, a **step directly corresponds** to *in its state machine*
one object (namely u) takes a **single transition** between regular states.

(We have to extend the concept of “single transition” for hierarchical state machines.)

That is: We’re going for an interleaving semantics without true parallelism.

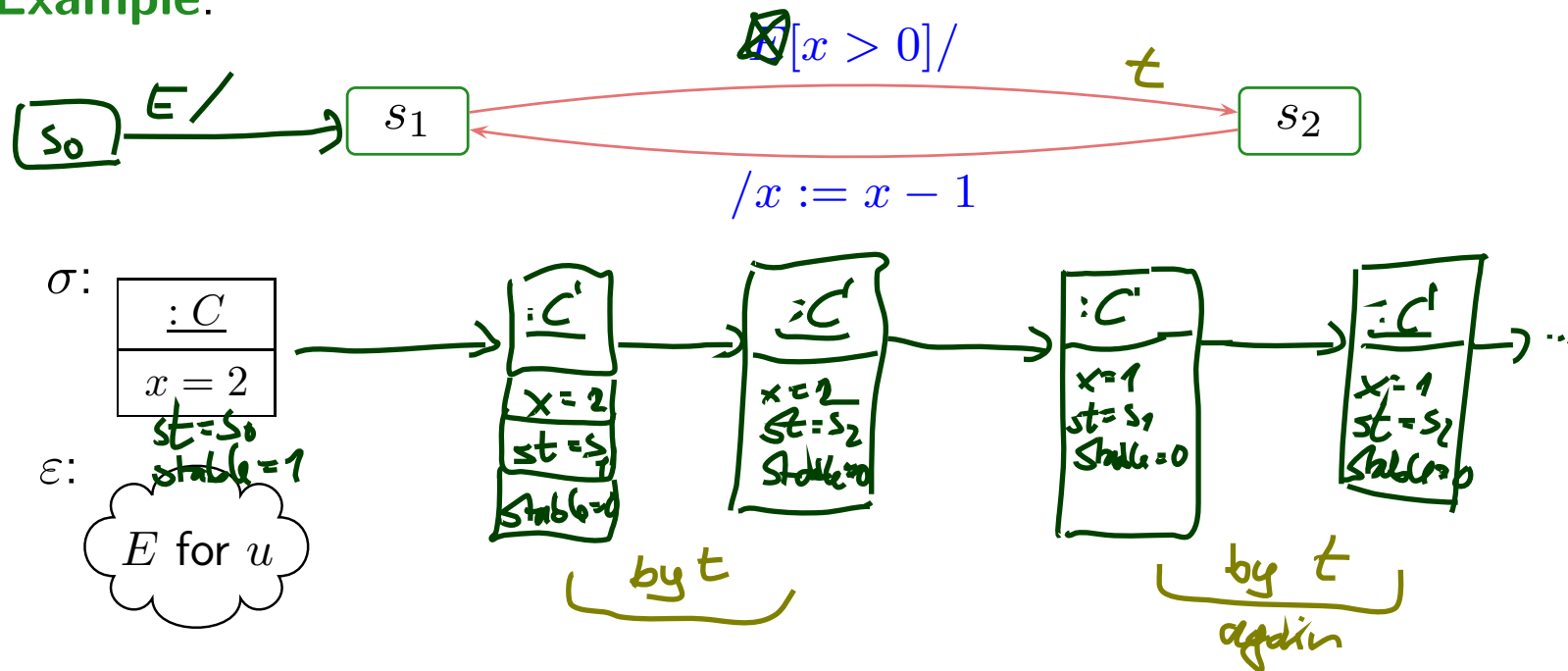
Notions of Steps: The Run-to-Completion Step

What is a **run-to-completion** step...?

- **Intuition:** a maximal sequence of steps, where the first step is a **dispatch** step and all later steps are **commence** steps.
- **Note:** one step corresponds to one transition in the state machine.

A run-to-completion step is in general not syntactically definable — one transition may be taken multiple times during an RTC-step.

Example:



Notions of Steps: The Run-to-Completion Step Cont'd

Proposal: Let

$$(\sigma_0, \varepsilon_0) \xrightarrow[u_0]{(cons_0, Snd_0)} \dots \xrightarrow[u_{n-1}]{(cons_{n-1}, Snd_{n-1})} (\sigma_n, \varepsilon_n), \quad n > 0,$$

we cannot extend the sequence and still satisfy the properties below

other objects may do interleaving steps

be a finite (!), non-empty, maximal, consecutive sequence such that

- object u is alive in σ_0 ,
- $u_0 = u$ and $(cons_0, Snd_0)$ indicates dispatching to u , i.e. $cons = \{(u, \vec{v} \mapsto \vec{d})\}$,
- there are no receptions by u in between, i.e.

$(\sigma_i, \varepsilon_i), (\sigma_{i+1}, \varepsilon_{i+1})$ are in transition relation

$$cons_i \cap \{u\} \times Evs(\mathcal{E}, \mathcal{D}) = \emptyset, i > 1,$$

- $u_{n-1} = u$ and u is stable only in σ_0 and σ_n , i.e.

$$\sigma_0(u)(stable) = \sigma_n(u)(stable) = 1 \text{ and } \sigma_i(u)(stable) = 0 \text{ for } 0 < i < n,$$

Let $0 = k_1 < k_2 < \dots < k_N = n$ be the maximal sequence of indices such that $u_{k_i} = u$ for $1 \leq i \leq N$. Then we call the sequence

$$(\sigma_0(u) =) \quad \sigma_{k_1}(u), \sigma_{k_2}(u) \dots, \sigma_{k_N}(u) \quad (= \sigma_{n-1}(u))$$

a (!) **run-to-completion computation** of u (from (local) configuration $\sigma_0(u)$).

Divergence

We say, object u **can diverge** on reception $cons$ from (local) configuration $\sigma_0(u)$ if and only if there is an infinite, consecutive sequence

$$(\sigma_0, \varepsilon_0) \xrightarrow{(cons_0, Snd_0)} (\sigma_1, \varepsilon_1) \xrightarrow{(cons_1, Snd_1)} \dots$$

such that u doesn't become stable again.

- **Note:** disappearance of object not considered in the definitions.
By the current definitions, it's neither divergence nor an RTC-step.

Run-to-Completion Step: Discussion.

What people may **dislike** on our definition of RTC-step is that it takes a **global** and **non-compositional** view. That is:

- In the projection onto a single object we still **see** the effect of interaction with other objects.
- Adding classes (or even objects) may change the divergence behaviour of existing ones.
- Compositional would be: the behaviour of a set of objects is determined by the behaviour of each object “in isolation”.

Our semantics and notion of RTC-step doesn't have this (often desired) property.

Can we give (syntactical) criteria such that any global run-to-completion step is an interleaving of local ones?

Maybe: Strict interfaces.

(Proof left as exercise...)

- **(A)**: Refer to private features only via “self” .
(Recall that other objects of the same class can modify private attributes.)
- **(B)**: Let objects only communicate by events, i.e.
don't let them modify each other's local state via links **at all**.

Putting It All Together

The Missing Piece: Initial States

Recall: a labelled transition system is (S, \rightarrow, S_0) . We **have**

- S : system configurations (σ, ε)
- \rightarrow : labelled transition relation $(\sigma, \varepsilon) \xrightarrow[u]{(cons, Snd)} (\sigma', \varepsilon')$.

Wanted: initial states S_0 .

Proposal:

Require a (finite) set of **object diagrams** OD as part of a UML model

class diagrams, induce signature, and system states $(\mathcal{C}\mathcal{D}, \mathcal{S}\mathcal{M}, \mathcal{O}\mathcal{D})$. *object diagrams, induce system states, the initial ones*
(statemachine diagrams, induce core state machines, and $\xrightarrow[u]{(cons, Snd)}$)

And set

$$S_0 = \{(\sigma, \varepsilon) \mid \sigma \in G^{-1}(OD), OD \in \mathcal{O}\mathcal{D}, \varepsilon \text{ empty}\}.$$

Other Approach: (used by Rhapsody tool) multiplicity of classes.

We can read that as an abbreviation for an object diagram.

Semantics of UML Model — So Far

The **semantics** of the **UML model**

$$\mathcal{M} = (\mathcal{C}\mathcal{D}, \mathcal{SM}, \mathcal{O}\mathcal{D})$$

where

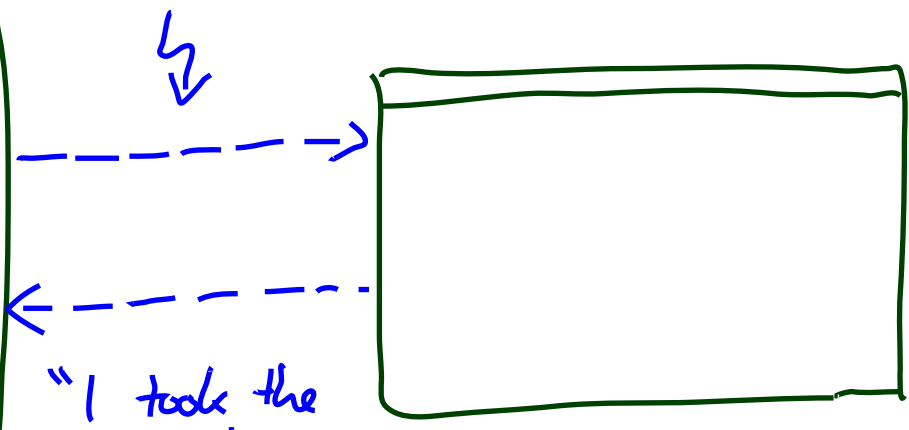
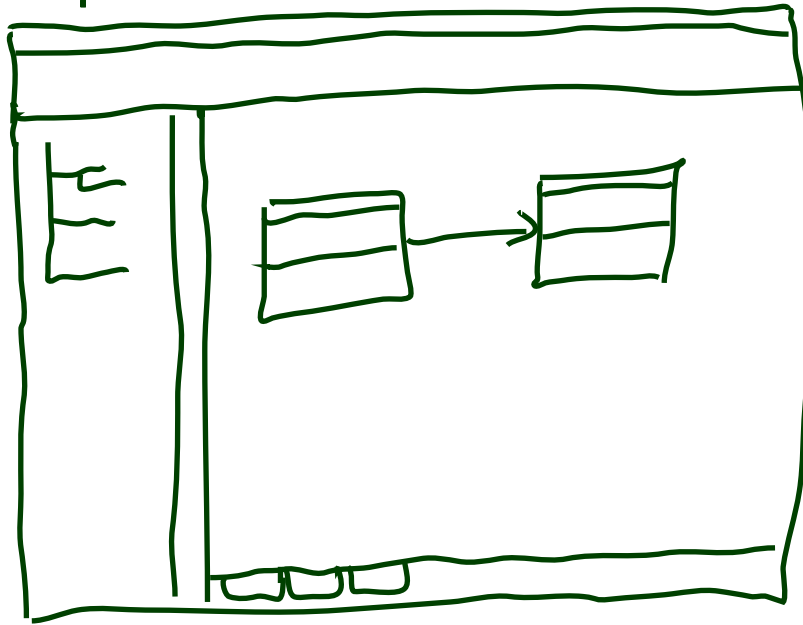
- some classes in $\mathcal{C}\mathcal{D}$ are stereotyped as 'signal' (standard), some signals and attributes are stereotyped as 'external' (non-standard),
- there is a 1-to-1 relation between classes and state machines,
- $\mathcal{O}\mathcal{D}$ is a set of object diagrams over $\mathcal{C}\mathcal{D}$,

is the **transition system** (S, \rightarrow, S_0) constructed on the previous slide.

The **computations of** \mathcal{M} are the computations of (S, \rightarrow, S_0) .

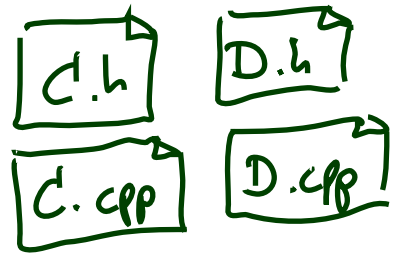
Contemporary UML Modelling Tools

Rhapsody



"I took the transition from s_0 to s_1 "

generate



...



build/make (compiling)



Run

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