Software Design, Modelling and Analysis in UML

Lecture 02: Semantical Model

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Why (of all things) UML?

[Kastens and Büning, 2008] sider as examples:

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A Brief History of UML

 1970's, Software CrisisTM
— Idea: learn from engineering disciplines to handle growing complexity. Boxes/lines and finite automata are used to visualise software for ages.

Mid 1980's: Statecharts [Harel, 1987], StateMateTM [Harel et al., 1990]

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Languages: Flowcharts, Nassi-Shneiderman, Entity-Relation Diagrams

- being a modelling languages doesn't mean being graphical (or: being a visual formalism [Harel]).

 Propositional and Predicate Logic Terms and Algebras Sets, Relations, Functions

- XML Schema, Entity Relation Diagrams, UML Class Diagrams
 Finite Automata, Petri Nets, UML State Machines

- Pro: visual formalisms are found appealing and easier to grasp.
 Yet they are not necessarily easier to write!
- Beware: you may meet people who dislike visual formalisms just for being graphical maybe because it is easier to "trick" people with a meaningless picture than with a meaningless formula.
- More serious: it's maybe easier to misunderstand a picture than a formula. $_{4,\rm 22}$

Contents & Goals

Last Lecture:

- Motivation: model-based development of things (houses, software) to cope with complexity, detect errors early
 Model-based (or-driven) Software Engineering
 UML Mode of the Lecture: Blueprint.

This Lecture:

- Educational Objectives: Capabilities for these tasks/questions:
- Why is UML of the form it is?
- What is a signature, an object, a system state, etc.?
 What's the purpose of signature, object, etc. in the course? Shall one feel bad if not using all diagrams during software development?

How do Basic Object System Signatures relate to UML class diagrams?

- Brief history of UML
- Basic Object System Signature, Structure, and System State

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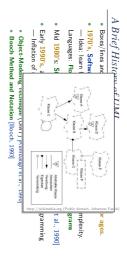
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- Early 1990's, advent of Object-Oriented-Analysis/Design/Programming - Inflation of notations and methods, most prominent:
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 Booch Method and Notation [Booch, 1993]
- Object-Oriented Software Engineering (OOSE) [Jacobson et al., 1992]

Each "persuasion" selling books, tools, seminars...

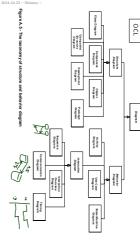
 Late 1990's: joint effort UML 0.x, 1.x Standards published by **Object Management Group** (OMG), "international, open membership, not-for-profit computer industry consortium".

• Since 2005: UML 2.x

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OCL Diagram

UML Overview [OMG, 2007b, 684]



Common Expectations on UML

UML Overview [OMG, 2007b, 684]

OCL

Diagram

Baharahr
Dayen

Use Case
Diagram

Stagram

Interaction Overview Diagram

Timing Diagram

- Powerful enough to bridge the gap between idea and implementation Easily writeable, readable even by customers
- Means to tame complexity by separation of concerns ("views")
- Unambiguous
- UML standard says how to develop software Standardised, exchangeable between modelling tools
- Using UML leads to better software

We will see...

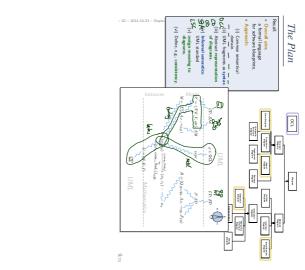
Seriously: After the course, you should have an own opinion on each of these claims. In how far/in what sense does it hold? Why? Why not? How can it be achieved? Which ones are really only hopes and expectations? ...?

Figure A.5 - The taxonomy of structure and behavior diagram

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Course Map Revisited

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Common Semantical Domain

Basic Object System Signature Example

Basic Object System Signature

Definition. A (Basic) Object System Signature is a quadruple

where $\mathcal{S}=(\mathcal{R} \otimes V_{i} atr) \quad \begin{array}{c} \text{for each dass $C \in \mathcal{C}$}\\ \text{where} \\ \text{of C_{i} as et of (basic) types,} \\ \text{of C_{i} a finite set of (classes,} \\ \text{of V is a finite set of typed attributes, i.e., each $v \in V$ has type} \end{array}$

(written $v:\tau$ or $v:C_{0,1}$ or $v:C_*$), • $C_{0,1}$ or C_* , where $C \in \mathscr{C}$

 $atr:\mathscr{C} \rightarrow 2$ maps each class to its set of attributes. Lotel function processed of V

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 (basic) types 𝒯 and classes 𝒞, (both finite),
 typed attributes V, 𝒯 from 𝒯 or C_{0.1} or C_{*}, C ∈ 𝒞,
 atr : 𝒞 → 2^V mapping classes to attributes. $\mathcal{S} = (\mathcal{T}, \mathcal{C}, V, atr)$ where

 $\mathcal{S}_0 = (\{Int\}, \{C, D\}, \{\underline{x}: \underline{Int}, D: C_{0.1}, n: C_r\}, \{C \mapsto \{p, n\}, D \mapsto \{x\}\}\})$ $= \{fint\}, \{C, D\}, \{\underline{x}: \underline{Int}, D: C_{0.1}, n: C_r\}, \{C \mapsto \{p, n\}, D \mapsto \{x\}\}\}\}$ $= \{fint\}, \{C, D\}, \{\underline{x}: \underline{Int}, D: C_{0.1}, n: C_r\}, \{C \mapsto \{p, n\}, D \mapsto \{x\}\}\}\}$ $= \{fint\}, \{C, D\}, \{\underline{x}: \underline{Int}, D: C_{0.1}, n: C_r\}, \{C \mapsto \{p, n\}, D \mapsto \{x\}\}\}\}$ $= \{fint\}, \{C, D\}, \{\underline{x}: \underline{Int}, D: C_{0.1}, n: C_r\}, \{C \mapsto \{p, n\}, D \mapsto \{x\}\}\}\}$ $= \{fint\}, \{C, D\}, \{\underline{x}: \underline{Int}, D: C_{0.1}, n: C_r\}, \{C \mapsto \{p, n\}, D \mapsto \{x\}\}\}\}$ $= \{fint\}, \{C, D\}, \{\underline{x}: \underline{Int}, D: C_{0.1}, n: C_r\}, \{C \mapsto \{p, n\}, D \mapsto \{x\}\}\}$ $= \{fint\}, \{C, D\}, \{\underline{x}: \underline{Int}, D: C_{0.1}, n: C_r\}, \{C \mapsto \{p, n\}, D \mapsto \{x\}\}\}$

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Basic Object System Signature Another Example

Example: \(\{ \frac{\Phi}{\Phi} \}, \{ A, B, \overline{\Omega}, \{ \frac{\Phi}{\Phi} \}, \{ \frac{\Phi}{\Phi}, \{ \frac{\Phi}{\Phi} \}, \{ \frac{\Phi}{\Phi} \}, \{ \frac{\Phi}{\Phi}, \{ \frac{\Phi}{\Phi} \}, \{ \frac{\Ph $\begin{tabular}{ll} & (basic) \ types \ \mathcal{T} \ and \ classes \ \mathscr{C}, \ (both \ finite), \\ & \ typed \ attributes \ V, \ \tau \ from \ \mathcal{T} \ or \ C_{0,1} \ or \ C_{*}, \ C \in \mathscr{C}, \\ & \ atr : \mathscr{C} \to 2^{V} \ mapping \ classes \ to \ attributes. \\ \end{tabular}$ $\mathcal{S} = (\mathcal{T}, \mathcal{C}, V, atr)$ where

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Basic Object System Structure

Definition. A Basic Object System Structure of $\mathscr{S}=(\mathscr{T}\otimes V, dt^*)$ is a domain function \mathscr{D} which assigns to each type a domain, i.e. $*\tau\in\mathscr{T}$ is mapped to $\mathscr{D}(\tau)$, $*C\in\mathscr{C}$ is mapped to $\mathscr{D}(\tau)$, Note: Object identities only have the "—" operation: Object identities of different classes are disjoint, i.e. $\forall C,D\in\mathscr{C}:C\neq D\to\mathscr{D}(\cap\mathscr{D})=\emptyset$ are mapped to $\mathscr{D}(C):D=\emptyset$. We use $\mathscr{D}(\mathscr{C})$ to denote $\bigcup_{C\in\mathscr{C}}\mathscr{D}(C):D=\emptyset$ analogously $\mathscr{D}(\mathscr{C}_*)$.

Note: We identify objects and object identities, because both uniquely determine each other (cf. OCL 2.0 standard).

System State deport of all V to the form for a factor of the state of V to the form of the state of the system state of V and V to the system state of V and V and V. A system state of V and V are V and V as system state of V and V are V are V and V are

We call $u \in \mathscr{D}(\mathscr{C})$ alive in σ if and only if $u \in \mathrm{dom}(\sigma)$. We use $\Sigma_{\mathscr{S}}^{\mathscr{D}}$ to denote the set of all system states of \mathscr{S} wrt. \mathscr{D} .

Basic Object System Structure Example

Wanted: a structure for signature

 $\mathscr{S}_0 = (\{Int\}, \{C, D\}, \{x: Int, p: C_{0,1}, n: C_*\}, \{C \mapsto \{p, n\}, D \mapsto \{x\}\})$

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Recall: by definition, seek a \mathscr{D} which maps r \in \mathscr{T} to some \mathscr{D}(r), r \in \mathscr{C} to some identities \mathscr{D}(C) (infinite, disjoint for different classes), C_{i_1} and C_{i_2} for C \in \mathscr{C} to \mathscr{D}(C_{i_1}) = \mathscr{D}(C_{i_2}) = \mathscr{D}^{\otimes (C)}.
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$$\begin{split} \mathscr{D}(Int) &= \mathbb{Z} \\ \mathscr{D}(C) &= \mathbb{N}^{l_{\infty}} \{ \mathcal{E}_{\delta}^{l_{\infty}} \mathcal{E}_{\delta_{c}, l_{\infty}} \} \\ \mathscr{D}(C) &= \mathbb{N}^{l_{\infty}} \{ \mathcal{E}_{\delta}^{l_{\infty}} \mathcal{E}_{\delta_{c}, l_{\infty}} \} \\ \mathscr{D}(D) &= \mathbb{N}^{l_{\infty}} \mathcal{E}_{\delta}^{l_{\infty}} \mathcal{E}_{\delta_{c}, l_{\infty}, l_{\infty}} \} \\ \mathscr{D}(C_{0,1}) &= \mathscr{D}(C) \\ \mathscr{D}(C_{0,1}) &= \mathscr{D}(C) \\ \mathscr{D}(D_{0,1}) &= \mathscr{D}(D) \\ \mathscr{D}(D$$

System State Example

Signature, Structure:

$$\mathcal{G}_0 = (\{Int\}, \{C, D\}, \{x: Int, p: C_{0,1}, n: C_*\}, \{C \mapsto \{p, n\}, D \mapsto \{x\}\}))$$

$$\mathcal{G}(Int) = \mathbf{Z}, \quad \mathcal{G}(C) \neq \{1_C, 2_C, 3_C, \dots\}, \quad \mathcal{G}(D) = \{1_D, 2_D, 3_D, \dots\}$$

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Wanted: \sigma: \mathcal{D}(\mathcal{C}) \rightarrow (V \rightarrow (\mathcal{D}(\mathcal{C}) \cup \mathcal{D}(\mathcal{C}_n))) such that f_n all v \in d^m(\mathcal{C})

• \sigma(u)(v) \in \mathcal{D}(v) if v: \tau, \tau \in \mathcal{F}, • \sigma(u)(v) \in \mathcal{D}(C_n) if v: D, with D \in \mathcal{C}.

• \sigma_{\mathcal{C}}(u)(v) \in \mathcal{D}(v) if v: \tau, \tau \in \mathcal{F}, • \sigma(u)(v) \in \mathcal{D}(C_n) if v: D, with D \in \mathcal{C}.

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System State Example

Signature, Structure:

$$\begin{split} \mathcal{S}_{0} &= (\{Int\}, \{C, D\}, \{x: Int, p: C_{0,1}, n: C_{\star}\}, \{C \mapsto \{p, n\}, D \mapsto \{x\}\}) \\ &\mathcal{D}(Int) = \mathbf{Z}, \quad \mathcal{D}(C) = \{1_{C}, 2_{C}, 3_{C}, \dots\}, \quad \mathcal{D}(D) = \{1_{D}, 2_{D}, 3_{D}, \dots\} \end{split}$$

```
Wanted: \sigma: \mathscr{D}(\mathscr{C}) \to (V \to (\mathscr{D}(\mathscr{T}) \cup \mathscr{D}(\mathscr{C}_*))) such that dom(\sigma(u)) = dir(C), \sigma(u)(v) \in \mathscr{D}(r) \text{ if } v: \tau, \tau \in \mathscr{T}, \sigma(u)(v) \in \mathscr{D}(r) \text{ if } v: D, \text{ with } D \in \mathscr{C} \text{ .}
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Concrete, explicit:

 $\sigma = \{1_C \mapsto \{p \mapsto \emptyset, n \mapsto \{5_C\}\}, 5_C \mapsto \{p \mapsto \emptyset, n \mapsto \emptyset\}, 1_D \mapsto \{x \mapsto 23\}\}.$

Alternative: symbolic system state

 $\sigma = \{c_1 \mapsto \{p \mapsto \emptyset, n \mapsto \{c_2\}\}, c_2 \mapsto \{p \mapsto \emptyset, n \mapsto \emptyset\}, d \mapsto \{x \mapsto 23\}\}$

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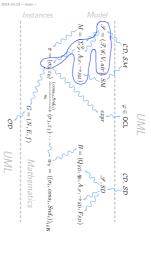
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Course Map



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References

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