

Lecture 09: Class Diagrams III

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Contents & Goals

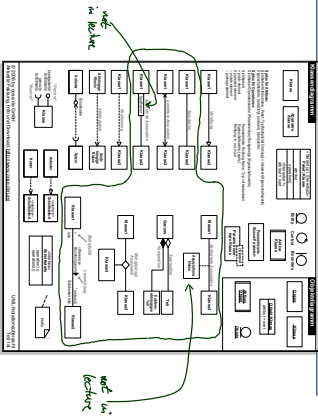
Last Lectures:

- completed class diagrams... except for associations

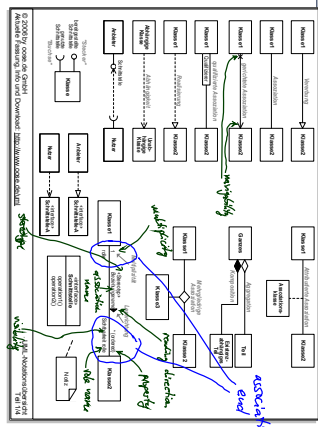
This Lecture:

- **Educational Objectives:** Capabilities for following tasks/questions:
 - Please explain this class diagram with associations.
 - Which annotations of an association arrow are semantically relevant?
 - What's a role name? What's it good for?
 - What is "multiplicity"? How did we treat them semantically?
 - What is "reading direction"? "navigability"? "ownership"? ...?
 - What's the difference between "aggregation" and "composition"?
- **Content:**
 - Study concrete syntax for "associations".
 - (**Temporarily**) extend signature, define mapping from diagram to signature.
 - Study effect on OCL.
 - **Box:** where do we put OCL constraints?

UML Class Diagram Syntax [Oesterich, 2006]

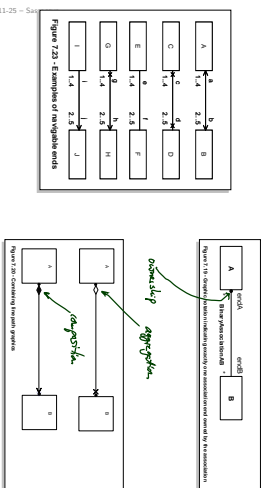


UML Class Diagram Syntax [Oesterich, 2006]



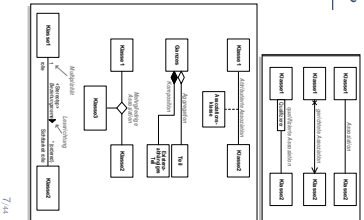
Associations: Syntax

UML Class Diagram Syntax [OMG, 2007b, 61:43]



What Do We (Have to) Cover?

- An association has
 - a name.
 - a reading direction and *just a link*
 - at least two ends.
 - a set of *Stakeholders* *by the design*
- Each end has
 - a role name.
 - a multiplicity.
 - a set of properties, such as *unique, ordered, etc.*
- a qualifier. (*leaf in UML*)
- Visibility.
- a navigability.
- an ownership.
- and possibly a diamond (*associates*)



(Temporarily) Extend Signature: Associations

- Only for the course of Lectures 9/10 we assume that each attribute in \mathcal{V}
 - either is $(v : \tau, \xi, \text{expr}_{\sigma_1}, P_1)$ with $\tau \in \mathcal{S}$ (as before).
 - or is an association of the form

$$(r : \langle \text{role}_1 : C_1, \mu_1, P_1, \xi_1, \nu_1, \sigma_1 \rangle, \dots, \langle \text{role}_n : C_n, \mu_n, P_n, \xi_n, \nu_n, \sigma_n \rangle)$$

the class refers this association end is lower

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where $n \geq 2$ (at least two ends), r, role_i are just names, $C_i \in \mathcal{C}$, $1 \leq i \leq n$, the multiplicity μ_i is an expression of the form $\mu ::= * | N | N..M | N..* | \mu, \mu$

P_i is a set of properties (as before), $\xi \in \{+, -, \#, \sim\}$ (as before), $\nu_i \in \{X, \rightarrow, \sim\}$ is the navigability, $\sigma_i \in B$ is the ownership.

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Alternative syntax for multiplicities: $\mu ::= N..M | N..* | \mu, \mu$ and define * and N as abbreviations.

Note: N could abbreviate 0, N, 1..N or N..N. We use last one.

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(Temporarily) Extend Signature: Basic Type Attributes

- Also only for the course of this lecture
 - we only consider basic type attributes to "belong" to a class (to appear in $\text{atr}(C)$).
 - associations are not "owned" by a particular class (do not appear in $\text{atr}(C)$) but live on their own.

Formally: we only call $(\mathcal{S}, \mathcal{C}, \mathcal{V}, \text{atr})$ a signature (extended for associations) if

$$\text{atr} : \mathcal{C} \rightarrow 2^{\text{[set of } \tau \in \mathcal{V}]}$$

(Temporarily) Extend Signature: Associations

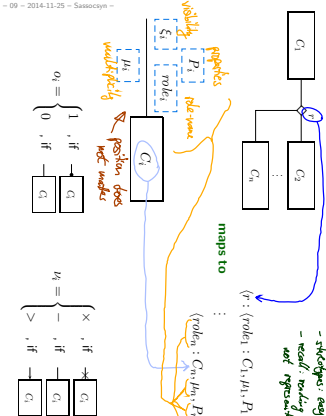
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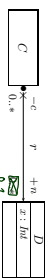
From Association Lines to Extended Signatures



maps to $(r : \langle \text{role}_1 : C_1, \mu_1, P_1, \xi_1, \nu_1, \sigma_1 \rangle, \dots, \langle \text{role}_n : C_n, \mu_n, P_n, \xi_n, \nu_n, \sigma_n \rangle)$

stakeholders: easy to add - recall reading function not represented

$\mu_i = \begin{cases} X, & \text{if } \rightarrow \\ -, & \text{if } \leftarrow \\ 0, & \text{if } \leftarrow \end{cases}$



Signature:

$S = \{ \langle C, r \rangle, \langle D, x \rangle \}$
 $\langle r: \langle \langle C, 0..*, r \rangle, \langle D, 0..1, x \rangle \rangle \}$
 $\langle x: \langle \langle C, 0..*, r \rangle, \langle D, 0..1, x \rangle \rangle \}$
only basic type allowed

Most components of associations or association end may be omitted. For instance [OMG, 2007b, 17], Section 6.4.2, proposes the following rules:

- Name: Use $A_i(C_1) \dots (C_n)$

if the name is missing

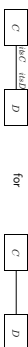


- Reading Direction: no default.

- Role Name: Use the class name at that end in lower-case letters



- Other convention: (used e.g. by modeling tool Riprapd)



- Multiplicity: 1 ^{A*} In my opinion, it's safer to assume 0..1 or * if there are no fixed, written, agreed conventions ("expect the worst").
- Properties: \emptyset
- Visibility: public
- Navigability and Ownership: not so easy; [OMG, 2007b, 43]

Various options may be chosen for showing navigation arrows on a diagram. In practice, it is often convenient to suppress some of the arrows and crosses and just show exceptional situations:

- Show all arrows and 'x's. Navigation and its absence are made completely explicit.
- Suppress all arrows and 'x's. No inference can be drawn about navigation. This is similar to any situation in which information is suppressed from a view.
- Suppress arrows for associations with navigability in both directions, and show arrows only for associations with one-way navigability.
- In this case, the two-way navigability cannot be distinguished from situations where there is no navigation at all; however, the latter case occurs rarely in practice.

Wait, If Omitting Things...

- ...is causing so much trouble (e.g. leading to misunderstanding), why does the standard say "In practice, it is often convenient..."?
- Is it a good idea to trade convenience for precision/unambiguity?

It depends

- Convenience as such is a legitimate goal.
- In UML-As-Sketch mode, precision "doesn't matter", so convenience (for writer) can even be a primary goal.
- In UML-As-Blueprint mode, precision is the primary goal. And misunderstandings are in most cases annoying.
- But: (even in UML-As-Blueprint mode) If all associations in your model have multiplicity *, then probably a good idea to suppress *'s. So, tell the reader about it and leave out the *'s.

Association Semantics

Overview

What's left? Named association with at least two typed ends, each having

- a role name, a set of properties, a navigability, and a multiplicity, a visibility, an ownership.

The Plan:

- Extend system states, introduce so-called links as instances of associations — depends on name and on type and number of ends.
- Integrate role name and multiplicity into OCL syntax/semantics.
- Extend typing rules to care for visibility and navigability.
- Consider multiplicity also as part of the constraints set $InV(C,D)$.
- Properties: for now assume $P_2 = \{ \text{int} : \text{int} \}$.
- Properties: (in general) and ownership: later.

Associations in General

Recall: We consider associations of the following form:

$$\langle r : \langle role_1 : C_1, role_2 : P_1, \dots, role_n : C_n, role_{n+1} : P_n, role_{n+2} : C_{n+1}, \dots \rangle \rangle$$

Only these parts are relevant for extended system states:

$$\langle r : \langle role_1 : C_1, \dots, role_n : C_n, \dots \rangle \rangle$$

(recall: we assume $P_i = P_i = \{\text{unique}\}$).

The UML standard thinks of associations as **n-ary relations** which **"live on their own"** in a system state.

- That is, **links** (= association instances)
- do not belong** (in general) to certain objects (in contrast to pointers, e.g.)
- are "first-class citizens" **next to objects**.
- are (in general) **not directed** (in contrast to pointers).

Links in System States

$$\langle r : \langle role_1 : C_1, \dots, role_n : P_n, \dots \rangle \rangle$$

Only for the course of Lectures 9/10 we change the definition of system states:

Definition. Let \mathcal{D} be a structure of the (extended) signature $\mathcal{S} = (\mathcal{S}, \mathcal{R}, V, \text{attr})$.

A system state of \mathcal{S} wrt. \mathcal{D} is a pair (σ, λ) consisting of

- a type-consistent mapping $\sigma : \mathcal{S}(\mathcal{C}) \rightarrow (\text{attr}(\mathcal{D}) \rightarrow \mathcal{D}(\mathcal{C}))$, *under the basic object only*
- a mapping λ which assigns each association $\langle r : \langle role_1 : C_1, \dots, role_n : C_n \rangle \rangle \in \mathcal{V}$ a relation $\lambda(r) \subseteq \mathcal{D}(C_1) \times \dots \times \mathcal{D}(C_n)$ (i.e. a set of type-consistent n -tuples of identities)

Association Semantics: The System State Aspect

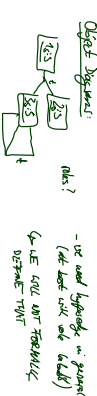


Example

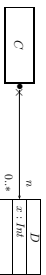
$\sigma = \{ \langle \text{id} : \text{id}(a), \text{id} : \text{id}(b) \rangle, \langle \text{id} : \text{id}(a), \text{id} : \text{id}(c) \rangle, \langle \text{id} : \text{id}(b), \text{id} : \text{id}(c) \rangle, \langle \text{id} : \text{id}(c), \text{id} : \text{id}(d) \rangle, \langle \text{id} : \text{id}(d), \text{id} : \text{id}(e) \rangle \}$

$\lambda = \{ \langle \text{id} : \text{id} \rangle \rightarrow \{ \langle \text{id} : \text{id}(a), \text{id} : \text{id}(b) \rangle, \langle \text{id} : \text{id}(a), \text{id} : \text{id}(c) \rangle, \langle \text{id} : \text{id}(b), \text{id} : \text{id}(c) \rangle \}, \langle \text{id} : \text{id}(b) \rangle \rightarrow \{ \langle \text{id} : \text{id}(b), \text{id} : \text{id}(c) \rangle \}, \langle \text{id} : \text{id}(c) \rangle \rightarrow \{ \langle \text{id} : \text{id}(c), \text{id} : \text{id}(d) \rangle \}, \langle \text{id} : \text{id}(d) \rangle \rightarrow \{ \langle \text{id} : \text{id}(d), \text{id} : \text{id}(e) \rangle \} \}$

one student may assume all able (add an O2 constraint if not obvious)



Association/Link Example



Signature

$$\mathcal{S} = \{ \langle \text{id} : \text{id} \rangle, \langle C, D \rangle, \langle x : \text{int} \rangle \}$$

$$\langle AC, D : \langle c : C, 0..* \rangle, \langle \text{unique} \rangle, \langle x : \text{int} \rangle \rangle$$

$$\langle m : D, 0..* \rangle, \langle \text{unique} \rangle, \langle > : 0 \rangle \rangle$$

$$\langle C \rightarrow 0..D \rightarrow \{ \} \rangle \rangle$$

A system state of \mathcal{S} (some reasonable \mathcal{D}) is (σ, λ) with

$$\sigma = \{ \langle \text{id} : \text{id} \rightarrow 0..3 \rangle \rightarrow \{ x \rightarrow 1 \}, \langle \text{id} : \text{id} \rightarrow \{ x \rightarrow 2 \} \}$$

$$\lambda = \{ \langle AC, D \rangle \rightarrow \{ \langle \text{id} : \text{id}, 3 \rangle, \langle \text{id} : \text{id}, 2 \rangle \} \}$$

Extended System States and Object Diagrams

Legitimate question: how do we represent system states such as

$$\sigma = \{ \langle \text{id} : \text{id} \rightarrow 0..3 \rangle \rightarrow \{ x \rightarrow 1 \}, \langle \text{id} : \text{id} \rightarrow \{ x \rightarrow 2 \} \}$$

$$\lambda = \{ \langle AC, D \rangle \rightarrow \{ \langle \text{id} : \text{id}, 3 \rangle, \langle \text{id} : \text{id}, 2 \rangle \} \}$$

as object diagram?

References

43/44

[Osterweich, 2006] Osterweich, B. (2006). *Analyse und Design mit UML 2.1. 8. Auflage*. Oldenbourg, 8. edition.

[OMG, 2007a] OMG (2007a). *Unified modeling language: Infrastructure, version 2.1.2*. Technical Report formal/07-11-04.

[OMG, 2007b] OMG (2007b). *Unified modeling language: Superstructure, version 2.1.2*. Technical Report formal/07-11-02.

44/44