Software Design, Modelling and Analysis in UML

Lecture 04: OCL Semantics

2014-10-30

Prof. Dr. Andreas Podelski, **Dr. Bernd Westphal**

Albert-Ludwigs-Universität Freiburg, Germany

Contents & Goals

Last Lecture:

- OCL Syntax

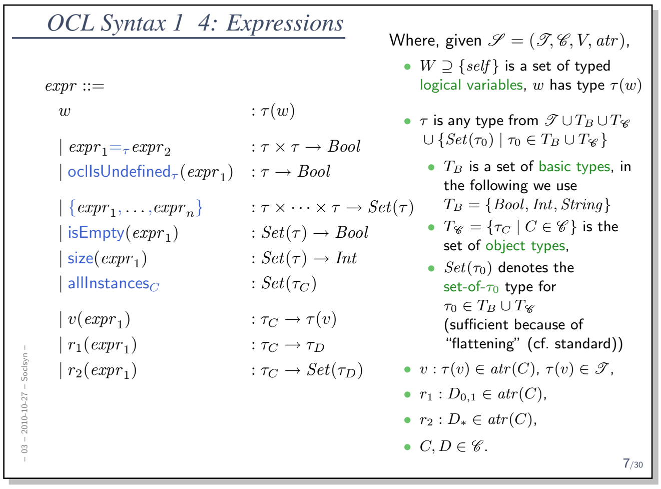
This Lecture:

- **Educational Objectives:** Capabilities for following tasks/questions.
  - What does it mean that an OCL expression is satisfiable?
  - When is a set of OCL constraints said to be consistent?
  - Can you think of an object diagram which violates this OCL constraint?

    next time
- **Content:**
  - OCL Semantics
  - maybe: OCL consistency and satisfiability

OCL Semantics [OMG, 2006]

The Task



OCL Syntax 1 4: Expressions

$$
\begin{array}{ll}
expr ::= & \\
\quad w & : \tau(w) \\
\quad \mid expr_1 =_\tau expr_2 & : \tau \times \tau \to Bool \\
\quad \mid \mathsf{oclIsUndefined}_\tau(expr_1) & : \tau \to Bool \\
\quad \mid \{expr_1, \dots, expr_n\} & : \tau \times \dots \times \tau \to Set(\tau) \\
\quad \mid \mathsf{isEmpty}(expr_1) & : Set(\tau) \to Bool \\
\quad \mid \mathsf{size}(expr_1) & : Set(\tau) \to Int \\
\quad \mid \mathsf{allInstances}_C & : Set(\tau_C) \\
\quad \mid v(expr_1) & : \tau_C \to \tau(v) \\
\quad \mid r_1(expr_1) & : \tau_C \to \tau_D \\
\quad \mid r_2(expr_1) & : \tau_C \to Set(\tau_D)
\end{array}
$$

Where, given $\mathscr{S} = (\mathscr{T}, \mathscr{C}, V, atr)$,

- $W \supseteq \{self\}$ is a set of typed logical variables, w has type $\tau(w)$
- τ is any type from $\mathscr{T} \cup T_B \cup T_\mathscr{C}$ $\cup \{Set(\tau_0) \mid \tau_0 \in T_B \cup T_\mathscr{C}\}$
  - T_B is a set of basic types, in the following we use $T_B = \{Bool, Int, String\}$
  - $T_\mathscr{C} = \{\tau_C \mid C \in \mathscr{C}\}$ is the set of object types,
  - $Set(\tau_0)$ denotes the set-of-τ_0 type for $\tau_0 \in T_B \cup T_\mathscr{C}$ (sufficient because of “flattening” (cf. standard))
- $v : \tau(v) \in atr(C)$, $\tau(v) \in \mathscr{T}$,
- $r_1 : D_{0,1} \in atr(C)$,
- $r_2 : D_* \in atr(C)$,
- $C, D \in \mathscr{C}$.

– 03 – 2010-10-27 – Soclsyn –

7/30

$I(\cdot,\cdot,\cdot)$

- Given an OCL expression $expr$, a system state $\sigma \in \Sigma_\mathscr{S}^\mathscr{D}$, and a valuation of logical variables β, define

$$I[\![\,\cdot\,]\!](\,\cdot\,,\,\cdot\,) : OCLExpressions(\mathscr{S}) \times \Sigma_\mathscr{S}^\mathscr{D} \times \underbrace{(W \to I(\mathscr{T} \cup T_B \cup T_\mathscr{C}))} \to \underbrace{I(Bool)}_{:= \{true,\ false,\ \bot_{Bool}\}}$$

such that

$$I[\![expr]\!](\sigma, \beta) \in \{\mathit{true}, \mathit{false}, \bot_{Bool}\}.$$

Basically business as usual...

(i) Equip each OCL (!) **basic type** with a reasonable **domain**, i.e. define function

$$I_{(i)} \text{ with } \operatorname{dom}(I_{(i)}) = T_B = \{Int, Bool, String\}$$

$$\text{e.g. } I_{(i)}(Bool) = \{true, false, \bot_{Bool}\}$$

Basically business as usual...

(i) Equip each OCL (!) **basic type** with a reasonable **domain**, i.e. define function

$$I_{(i)} \text{ with } \operatorname{dom}(I_{(i)}) = T_B$$

(ii) Equip each **object type** τ_C with a reasonable **domain**, i.e. define function

$$I_{(ii)} \text{ with } \operatorname{dom}(I_{(ii)}) = \tau_C$$

(most reasonable: $\mathscr{D}(C)$ determined by structure $\mathscr{D}$ of $\mathscr{S}$).

(iii) Equip each **set type** $Set(\tau_0)$ with reasonable **domain**, i.e. define function

$$I_{(iii)} \text{ with } \operatorname{dom}(I_{(iii)}) = \{Set(\tau_0) \mid \tau_0 \in T_B \cup T_{\mathscr{C}}\}$$

(iv) Equip each **arithmetical operation** with a reasonable **interpretation** (that is, with a **function** operating on the corresponding **domains**).

$$I_{(iv)} \text{ with } \operatorname{dom}(I_{(iv)}) = \{+, -, \leq, \dots\}, \text{ e.g., } I(+) \in I(Int) \times I(Int) \to I(Int)$$

$$expr_1 + expr_2 : Int \times Int \to Int$$

Basically business as usual...

(i) Equip each OCL (!) **basic type** with a reasonable **domain**, i.e. define function

$$I \text{ with } \operatorname{dom}(I) = T_B$$

(ii) Equip each **object type** τ_C with a reasonable **domain**, i.e. define function

$$I \text{ with } \operatorname{dom}(I) = \tau_C$$

(most reasonable: $\mathscr{D}(C)$ determined by structure $\mathscr{D}$ of $\mathscr{S}$).

(iii) Equip each **set type** $Set(\tau_0)$ with reasonable **domain**, i.e. define function

$$I \text{ with } \operatorname{dom}(I) = \{Set(\tau_0) \mid \tau_0 \in T_B \cup T_{\mathscr{C}}\}$$

(iv) Equip each **arithmetical operation** with a reasonable **interpretation** (that is, with a **function** operating on the corresponding **domains**).

$$I \text{ with } \operatorname{dom}(I) = \{+, -, \leq, \dots\}, \text{ e.g., } I(+) \in I(Int) \times I(Int) \to I(Int)$$

(v) **Set operations** similar: I with $\operatorname{dom}(I) = \{\text{isEmpty}, \dots\}$

(vi) Equip each **expression** with a reasonable **interpretation**, i.e. define function

$$I : Expr \times \Sigma_{\mathscr{S}}^{\mathscr{D}} \times (W \to I(\mathscr{T} \cup T_B \cup T_{\mathscr{C}})) \to I(Bool)$$

...except for OCL being a **three-valued logic**, and the "iterate" expression.

(i) Domains of Basic Types of OCL

Recall:

- $T_B = \{Bool, Int, String\}$

assume both sets disjoint

We set:

read "bottom" or "undefined"

- $I_{(i)}(Bool) := \{true, false\} \,\dot\cup\, \{\bot_{Bool}\}$
- $I_{(i)}(Int) := \mathbb{Z} \,\dot\cup\, \{\bot_{Int}\}$
- $I_{(i)}(String) := \underbrace{\dots}_{\text{finite sequences of characters}} \dot\cup\, \{\bot_{String}\}$

We may omit index τ of $\bot_\tau$ if it is clear from context.

(ii) Domains of Object and (iii) Set Types

- Now we need a structure $\mathscr{D}$ of our signature $\mathscr{S} = (\mathscr{T}, \mathscr{C}, V, atr)$.
- **Recall:** $\mathscr{D}$ assigns an (infinite) domain $\mathscr{D}(C)$ to each class $C \in \mathscr{C}$.

- Let τ_C be an (OCL) **object type** for a class $C \in \mathscr{C}$.
- We set

$$I_{(ii)}(\tau_C) := \mathscr{D}(C) \,\dot{\cup}\, \{\bot_{\tau_C}\}$$

- Let τ be a type from $T_B \cup T_{\mathscr{C}}$.
- We set

$$I_{(iii)}(Set(\tau)) := \overset{\text{powerset of } I(\tau)}{\overset{\downarrow}{2^{I(\tau)}}} \,\dot{\cup}\, \{\bot_{Set(\tau)}\}$$

Note: in the OCL standard, only **finite** subsets of $I(\tau)$.
But infinity doesn't scare **us**, so we simply allow it.

(iv) Interpretation of Arithmetic Operations

- **Literals** map to fixed values:

$$\underset{(iv)}{I}(\underset{\underset{Bool}{\because}}{\mathtt{true}}) := \underset{\underset{I_{(i)}(Bool)}{\cap}}{true}, \quad I(\mathtt{false}) := \overset{\overset{I_{(i)}(Bool)}{\Downarrow}}{false}, \quad I(\mathtt{0}) := \overset{\overset{I_{(i)}(Int) = \mathbb{Z} \,\dot{\cup}\, \{\bot_{Int}\}}{\Downarrow}}{0}, \quad I(\mathtt{1}) := 1, \ldots$$

$$I(\mathtt{OclUndefined}_\tau) := \bot_\tau$$

(iv) Interpretation of Arithmetic Operations

- **Literals** map to fixed values:

$$I(\mathsf{true}) := \mathit{true}, \quad I(\mathsf{false}) := \mathit{false}, \qquad I(0) := 0, \quad I(1) := 1, \ldots$$
$$I(\mathsf{OclUndefined}_\tau) := \perp_\tau$$

- **Boolean operations** (defined point-wise for $x_1, x_2 \in I(\tau)$):

$$\overset{\tau \times \tau \to Bool}{\underbrace{I(=_\tau)}_{I_{(iv)}(=_\tau) : I(\tau) \times I(\tau) \to I(Bool) = \{true, false, \perp\}}}(x_1, x_2) := \begin{cases} true & , \text{if } x_1 = x_2 \text{ and } x_1 \neq \perp_\tau \text{ and } x_2 \neq \perp_\tau \\ false & , \text{if } x_1 \neq x_2 \text{ and } x_1 \neq \perp_\tau \text{ and } x \neq \perp_\tau \\ \perp_{Bool} & , \text{otherwise} \end{cases}$$

(iv) Interpretation of Arithmetic Operations

- **Literals** map to fixed values:

$$I(\mathsf{true}) := \mathit{true}, \quad I(\mathsf{false}) := \mathit{false}, \qquad I(0) := 0, \quad I(1) := 1, \ldots$$
$$I(\mathsf{OclUndefined}_\tau) := \perp_\tau$$

- **Boolean operations** (defined point-wise for $x_1, x_2 \in I(\tau)$):

$$I(=_\tau)(x_1, x_2) := \begin{cases} \mathit{true} & , \text{if } x_1 \neq \perp_\tau \neq x_2 \text{ and } x_1 = x_2 \\ \mathit{false} & , \text{if } x_1 \neq \perp_\tau \neq x_2 \text{ and } x_1 \neq x_2 \\ \perp_{Bool} & , \text{otherwise} \end{cases}$$

- **Integer operations** (defined point-wise for $x_1, x_2 \in I(Int)$):

$$\overset{Int \times Int \to Int}{\underbrace{I(+)}_{I(Int) \times I(Int) \to I(Int) = \mathbb{Z} \cup \{\perp\}}}(x_1, x_2) := \begin{cases} x_1 + x_2 & , \text{if } x_1 \neq \perp \text{ and } x_2 \neq \perp \\ \perp & , \text{otherwise} \end{cases}$$

(iv) Interpretation of Arithmetic Operations

- **Literals** map to fixed values:

$$I(\mathsf{true}) := \mathit{true}, \quad I(\mathsf{false}) := \mathit{false}, \qquad I(0) := 0, \quad I(1) := 1, \dots$$
$$I(\mathsf{OclUndefined}_\tau) := \bot_\tau$$

- **Boolean operations** (defined point-wise for $x_1, x_2 \in I(\tau)$):

$$I(=_\tau)(x_1, x_2) := \begin{cases} \mathit{true} & \text{, if } x_1 \neq \bot_\tau \neq x_2 \text{ and } x_1 = x_2 \\ \mathit{false} & \text{, if } x_1 \neq \bot_\tau \neq x_2 \text{ and } x_1 \neq x_2 \\ \bot_{Bool} & \text{, otherwise} \end{cases}$$

- **Integer operations** (defined point-wise for $x_1, x_2 \in I(Int)$):

$$I(+)(x_1, x_2) := \begin{cases} x_1 + x_2 & \text{, if } x_1 \neq \bot \neq x_2 \\ \bot & \text{, otherwise} \end{cases}$$

$\omega(expr_1, \dots, expr_n)$ e.g. $+(expr_1, expr_2)$

Note: There is a **common principle**.
Namely, the **interpretation** of an operation $\omega : \tau_1 \times \dots \tau_n \to \tau$ is a function $I(\omega) : I(\tau_1) \times \dots \times I(\tau_n) \to I(\tau)$ on corresponding semantical domain(s).

$0 + 27 = 13$

$= (+(0, 27), 13)$

$I(0) = 0$
$I(27) = 27$
$I(13) = 13$
$I(+) : I(Int) \times I(Int) \to I(Int)$
$I(=) : I(Int) \times I(Int) \to I(Bool)$

actually $=_{Int}$ here

see previous slide

$$I[\![\omega(expr_1, \dots, expr_n)]\!](\sigma, \beta) := (I(\omega))(I[\![expr_1]\!](\sigma,\beta), \dots, I[\![expr_n]\!](\sigma, \beta))$$

$$I[\![+(0,27)]\!](\sigma,\beta) = \underbrace{(I(+))(\underbrace{I(0)}_{0}, \underbrace{I(27)}_{27})}_{27} = 27$$

(iv) Interpretation of OclIsUndefined

- The **is-undefined** predicate (defined point-wise for $x \in I(\tau)$):

$$I(\text{oclIsUndefined}_\tau)(x) := \begin{cases} \mathit{true} & \text{, if } x = \bot_\tau \\ \mathit{false} & \text{, otherwise} \end{cases}$$

(v) Interpretation of Set Operations

Basically the same principle as with arithmetic operations...

Let $\tau \in T_B \cup T_{\mathscr{C}}$.

- **Set comprehension** ($x_1, \dots, x_n \in I(\tau)$):

$$I(\{\}^\tau_n)(x_1, \dots, x_n) := \{x_1, \dots, x_n\}$$

  for all $n \in \mathbb{N}_0$

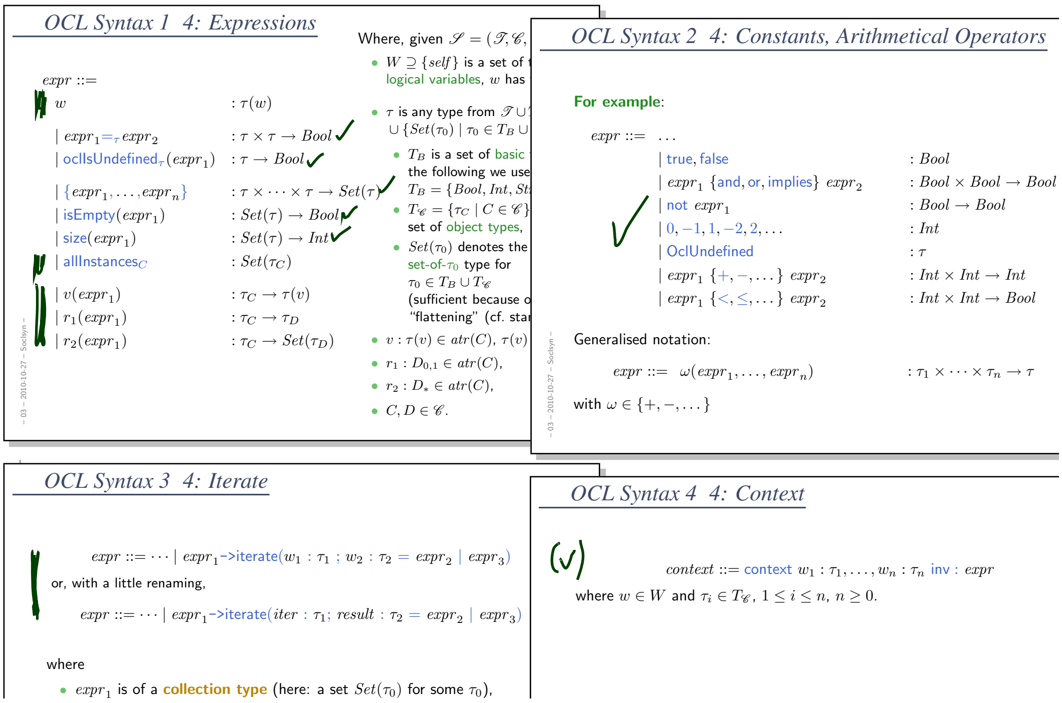
- **Empty-ness check** ($x \in I(Set(\tau))$):

$$I(\text{isEmpty}^\tau)(x) := \begin{cases} \mathit{true} & \text{, if } x = \emptyset \\ \bot_{Bool} & \text{, if } x = \bot_{Set(\tau)} \\ \mathit{false} & \text{, otherwise} \end{cases}$$

- **Counting** ($x \in I(Set(\tau))$):

$$I(\text{size}^\tau)(x) := \overset{\text{cardinality}}{|x|} \text{ if } x \neq \bot_{Set(\tau)} \text{ and } \bot_{Int} \text{ otherwise}$$

(vi) Putting It All Together



OCL Syntax 1 4: Expressions

$expr ::=$

w	$: \tau(w)$
$\mid expr_1 =_\tau expr_2$	$: \tau \times \tau \to Bool$ ✓
$\mid$ oclIsUndefined$_\tau(expr_1)$	$: \tau \to Bool$ ✓
$\mid \{expr_1, \dots, expr_n\}$	$: \tau \times \dots \times \tau \to Set(\tau)$ ✓
$\mid$ isEmpty$(expr_1)$	$: Set(\tau) \to Bool$ ✓
$\mid$ size$(expr_1)$	$: Set(\tau) \to Int$ ✓
$\mid$ allInstances$_C$	$: Set(\tau_C)$
$\mid v(expr_1)$	$: \tau_C \to \tau(v)$
$\mid r_1(expr_1)$	$: \tau_C \to \tau_D$
$\mid r_2(expr_1)$	$: \tau_C \to Set(\tau_D)$

Where, given $\mathscr{S} = (\mathscr{T}, \mathscr{C},$

- $W \supseteq \{self\}$ is a set of t logical variables, w has
- τ is any type from $\mathscr{T} \cup$ $\cup \{Set(\tau_0) \mid \tau_0 \in T_B \cup$
  - T_B is a set of basic the following we use $T_B = \{Bool, Int, St$
  - $T_{\mathscr{C}} = \{\tau_C \mid C \in \mathscr{C}\}$ set of object types,
  - $Set(\tau_0)$ denotes the set-of-τ_0 type for $\tau_0 \in T_B \cup T_{\mathscr{C}}$ (sufficient because o "flattening" (cf. star
- $v : \tau(v) \in atr(C)$, $\tau(v)$
- $r_1 : D_{0,1} \in atr(C)$,
- $r_2 : D_* \in atr(C)$,
- $C, D \in \mathscr{C}$.

OCL Syntax 2 4: Constants, Arithmetical Operators

For example:

$expr ::= \dots$	
$\mid$ true, false	$: Bool$
$\mid expr_1$ {and, or, implies} $expr_2$	$: Bool \times Bool \to Bool$
$\mid$ not $expr_1$	$: Bool \to Bool$
$\mid 0, -1, 1, -2, 2, \dots$	$: Int$
$\mid$ OclUndefined	$: \tau$
$\mid expr_1 \{+, -, \dots\}\ expr_2$	$: Int \times Int \to Int$
$\mid expr_1 \{<, \leq, \dots\}\ expr_2$	$: Int \times Int \to Bool$

✓

Generalised notation:

$$expr ::= \omega(expr_1, \dots, expr_n) \qquad : \tau_1 \times \dots \times \tau_n \to \tau$$

with $\omega \in \{+, -, \dots\}$

OCL Syntax 3 4: Iterate

$$expr ::= \dots \mid expr_1\text{->iterate}(w_1 : \tau_1 \ ; \ w_2 : \tau_2 = expr_2 \mid expr_3)$$

or, with a little renaming,

$$expr ::= \dots \mid expr_1\text{->iterate}(iter : \tau_1; \ result : \tau_2 = expr_2 \mid expr_3)$$

where

- $expr_1$ is of a **collection type** (here: a set $Set(\tau_0)$ for some τ_0),

OCL Syntax 4 4: Context

(✓)

$$context ::= \text{context } w_1 : \tau_1, \dots, w_n : \tau_n \text{ inv : } expr$$

where $w \in W$ and $\tau_i \in T_{\mathscr{C}}$, $1 \leq i \leq n$, $n \geq 0$.

Valuations of Logical Variables

$\{self_C \mid C \in \mathscr{C}\}$

$x \quad : \tau_C$

- **Recall**: we have typed logical variables $(w \in)\ W$, $\tau(w)$ is the type of w.
- By β, we denote a valuation of the logical variables, i.e. for each $w \in W$,

$$\beta : W \longrightarrow \left.\begin{matrix} I(Int) \\ \cup I(Bool) \\ \cup \dots \\ \cup I(\tau_C) \\ \cup \dots \end{matrix}\right\} \bigcup_{w \in W} I(\tau(w))$$

$$\beta(w) \in I(\tau(w)).$$

$$W = \{x : Int, \ self_C : \tau_C\}$$

$$\beta : W \longrightarrow I(Int) \cup I(\tau_C) = \mathbb{Z} \dot\cup \{\bot\} \cup \mathscr{D}(C)$$

Examples:

- $\beta(x) = 27 \in I(Int)$
  $\beta(self_C) = 1_C \in I(\tau_C) = \mathscr{D}(C)$
- $\beta(x) = \bot_{Int}$
  $\beta(self_C) = 5_C$

(vi) Putting It All Together...

$$expr ::= w \mid \omega(expr_1, \dots, expr_n) \mid \mathsf{allInstances}_C \mid v(expr_1) \mid r_1(expr_1) \mid r_2(expr_1) \mid expr_1\text{->}\mathsf{iterate}(v_1 : \tau_1 \,;\, v_2 : \tau_2 = expr_2 \mid expr_3)$$

- $I[\![w]\!](\sigma, \beta) := \beta(\omega)$
- $I[\![\omega(expr_1, \dots, expr_n)]\!](\sigma, \beta) := (I(\omega))(I[\![expr_1]\!](\sigma,\beta), \dots, I[\![expr_n]\!](\sigma,\beta))$
- $I[\![\mathsf{allInstances}_C]\!](\sigma, \beta) := \mathrm{dom}(\sigma) \cap \mathscr{D}(C)$

  Note: in the OCL standard, $\mathrm{dom}(\sigma)$ is assumed to be **finite**.
  Again: doesn't scare us.

(vi) Putting It All Together...

$$expr ::= w \mid \omega(expr_1, \dots, expr_n) \mid \mathsf{allInstances}_C \mid v(expr_1) \mid r_1(expr_1) \mid r_2(expr_1) \mid expr_1\text{->}\mathsf{iterate}(v_1 : \tau_1 \,;\, v_2 : \tau_2 = expr_2 \mid expr_3)$$

Assume $expr_1 : \tau_C$ for some $C \in \mathscr{C}$. Set $u_1 := I[\![expr_1]\!](\sigma, \beta) \in \mathscr{D}(\tau_C)$.

- $I[\![\underset{T}{v}(expr_1)]\!](\sigma, \beta) := \begin{cases} (\sigma(u_1))(v) & \text{, if } u_1 \in \mathrm{dom}(\sigma) \\ \bot & \text{, otherwise} \end{cases}$
- $I[\![\underset{D_{0,1}}{r_1}(expr_1)]\!](\sigma, \beta) := \begin{cases} u & \text{, if } u_1 \in \mathrm{dom}(\sigma) \text{ and } \sigma(u_1)(r_1) = \{u\} \\ \bot & \text{, otherwise} \end{cases}$
- $I[\![\underset{D_*}{r_2}(expr_1)]\!](\sigma, \beta) := \begin{cases} \sigma(u_1)(r_2) & \text{, if } u_1 \in \mathrm{dom}(\sigma) \\ \bot & \text{, otherwise} \end{cases}$

  (Recall: σ evaluates r_2 of type C_* to a set)

$\mathscr{S} = (\{Nat\}, \{TeamMember, Meeting\}, \{age : Nat, m : M_{0,1}, p : TM_*\},$
$\{TM \mapsto \{age, m\}, M \mapsto \{p\}\})$

$\mathscr{D}(Nat) = \mathbb{N}_0$

$W = \{x, self_M, self_{TM}\}$

$\sigma = \{1_{TM} \mapsto (27, 5_M), 2_{TM} \mapsto (17, 5_M), 5_M \mapsto (\{1_{TM}, 2_{TM}\})\}$

- $I[\![allInstances_{TM}]\!](\sigma,\beta) = dom(\sigma) \cap \mathscr{D}(TM) = \{1_{TM}, 2_{TM}, 5_M\} \cap \mathscr{D}(TM) = \{1_{TM}, 2_{TM}\}$
- $\beta_1 : x \mapsto 10, \ldots$

  $I[\![x > allInstances_{TM} \to size]\!](\sigma,\beta)$

  $$= I[\![> (x, size(allInstances_{TM}))]\!](\sigma,\beta) = \underbrace{I(>)\Big(\underbrace{I[\![x]\!](\sigma,\beta_1)}_{\beta_1(x)=10}, \underbrace{\overset{|\cdot|}{I(size)}\big(\underbrace{I[\![allInst_{TM}]\!](\sigma,\beta_1)}_{\{1_{TM},2_{TM}\}}\big)}_{2}\Big)}_{true}$$

- $\beta_2 : self_C \mapsto 2_{TM}, \ldots$ ← and some value for x

  $I[\![self_C . age]\!](\sigma,\beta_2) = I[\![age(self_C)]\!](\sigma,\beta_2)$

  $= \sigma(u_1)(age) = \sigma(2_{TM})(age) = 17$

  $u_1 = I[\![self_C]\!](\sigma,\beta_2) = \beta_2(self_C) = 2_{TM}$

- $\beta_3 : self_C \mapsto 7_{TM}, \ldots$ $I[\![self_C . age]\!](\sigma,\beta_3) = \bot$ because $7_{TM} \notin dom(\sigma)$

(vi) Putting It All Together...

$expr ::= w \mid \omega(expr_1, \ldots, expr_n) \mid \mathsf{allInstances}_C \mid v(expr_1) \mid r_1(expr_1)$
$\mid r_2(expr_1) \mid expr_1\text{->}\mathsf{iterate}(v_1 : \tau_1 \; ; v_2 : \tau_2 = expr_2 \mid expr_3)$

- $I[\![expr_1\text{->}\mathsf{iterate}(v_1 : \tau_1 \; ; v_2 : \tau_2 = expr_2 \mid expr_3)]\!](\sigma,\beta)$

  $$:= \begin{cases} I[\![expr_2]\!](\sigma,\beta) & \text{, if } I[\![expr_1]\!](\sigma,\beta) = \emptyset \\ iterate(hlp, v_1, v_2, expr_3, \sigma, \beta') & \text{, otherwise} \end{cases}$$

  where $\beta' = \beta[hlp \mapsto I[\![expr_1]\!](\sigma,\beta), v_2 \mapsto I[\![expr_2]\!](\sigma,\beta)]$ and
- $iterate(hlp, v_1, v_2, expr_3, \sigma, \beta')$

  $$:= \begin{cases} I[\![expr_3]\!](\sigma, \beta'[v_1 \mapsto x]) & \text{, if } \beta'(hlp) = \{x\} \\ I[\![expr_3]\!](\sigma, \beta'') & \text{, if } \beta'(hlp) = X \mathbin{\dot\cup} \{x\} \text{ and } X \neq \emptyset \end{cases}$$

  where $\beta'' = \beta'[v_1 \mapsto x, v_2 \mapsto iterate(hlp, v_1, v_2, expr_3, \sigma, \beta'[hlp \mapsto X])]$

(vi) Putting It All Together...

$$expr ::= w \mid \omega(expr_1, \dots, expr_n) \mid \mathsf{allInstances}_C \mid v(expr_1) \mid r_1(expr_1) \mid r_2(expr_1) \mid expr_1\text{->}\mathsf{iterate}(v_1 : \tau_1 \,;\, v_2 : \tau_2 = expr_2 \mid expr_3)$$

- $I[\![expr_1\text{->}\mathsf{iterate}(v_1 : \tau_1 \,;\, v_2 : \tau_2 = expr_2 \mid expr_3)]\!](\sigma, \beta)$

$$:= \begin{cases} I[\![expr_2]\!](\sigma, \beta) & \text{, if } I[\![expr_1]\!](\sigma, \beta) = \emptyset \\ iterate(hlp, v_1, v_2, expr_3, \sigma, \beta') & \text{, otherwise} \end{cases}$$

where $\beta' = \beta[hlp \mapsto I[\![expr_1]\!](\sigma, \beta), v_2 \mapsto I[\![expr_2]\!](\sigma, \beta)]$ and

- $iterate(hlp, v_1, v_2, expr_3, \sigma, \beta')$

$$:= \begin{cases} I[\![expr_3]\!](\sigma, \beta'[v_1 \mapsto x]) & \text{, if } \beta'(hlp) = \{x\} \\ I[\![expr_3]\!](\sigma, \beta'') & \text{, if } \beta'(hlp) = X \,\dot{\cup}\, \{x\} \text{ and } X \neq \emptyset \end{cases}$$

where $\beta'' = \beta'[v_1 \mapsto x, v_2 \mapsto iterate(hlp, v_1, v_2, expr_3, \sigma, \beta'[hlp \mapsto X])]$

Quiz: Is (our) I a function?

References

[Cabot and Clarisó, 2008] Cabot, J. and Clarisó, R. (2008). UML-OCL verification in practice. In Chaudron, M. R. V., editor, *MoDELS Workshops*, volume 5421 of *Lecture Notes in Computer Science*. Springer.

[Cengarle and Knapp, 2001] Cengarle, M. V. and Knapp, A. (2001). On the expressive power of pure OCL. Technical Report 0101, Institut für Informatik, Ludwig-Maximilians-Universität München.

[Cengarle and Knapp, 2002] Cengarle, M. V. and Knapp, A. (2002). Towards OCL/RT. In Eriksson, L.-H. and Lindsay, P. A., editors, *FME*, volume 2391 of *Lecture Notes in Computer Science*, pages 390–409. Springer-Verlag.

[Flake and Müller, 2003] Flake, S. and Müller, W. (2003). Formal semantics of static and temporal state-oriented OCL constraints. *Software and Systems Modeling*, 2(3):164–186.

[Jackson, 2002] Jackson, D. (2002). Alloy: A lightweight object modelling notation. *ACM Transactions on Software Engineering and Methodology*, 11(2):256–290.

[OMG, 2006] OMG (2006). Object Constraint Language, version 2.0. Technical Report formal/06-05-01.

[OMG, 2007a] OMG (2007a). Unified modeling language: Infrastructure, version 2.1.2. Technical Report formal/07-11-04.

[OMG, 2007b] OMG (2007b). Unified modeling language: Superstructure, version 2.1.2. Technical Report formal/07-11-02.