

# *Software Design, Modelling and Analysis in UML*

## *Lecture 7: Class Diagrams II*

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# Contents & Goals

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## Last Lecture:

- Representing class diagrams as (extended) signatures — for the moment without associations: later.

## This Lecture:

- **Educational Objectives:** Capabilities for following tasks/questions.
  - Could you please map this class diagram to a signature?
  - What if things are missing?
  - Could you please map this signature to a class diagram?
  - What is the semantics of 'abstract'?
  - What is visibility good for?
- **Content:**
  - Map class diagram to (extended) signature cont'd.
  - Stereotypes – for documentation.
  - Visibility as an extension of well-typedness.

## *Mapping UML CDs to Extended Signatures*

## Example Cont'd

Dinglars

|  |
|--|
| $\langle\langle S_1, \dots, S_k \rangle\rangle$<br>$C$                                 |
| $\xi_1 v_1 : T_1 = \text{expr}_0^1 \{P_{1,1}, \dots, P_{1,m_1}\}$                      |
| $\vdots$   |
| $\xi_\ell v_\ell : T_\ell = \text{expr}_0^\ell \{P_{\ell,1}, \dots, P_{\ell,m_\ell}\}$ |

$\Downarrow$

$\Downarrow$

$C(n) := \langle C, \{S_1, \dots, S_k\}, a(n), t(n) \rangle$

$V(n) := \{ \langle v_1 : T_1, \xi_1, \text{expr}_0^1, \{P_{1,1}, \dots, P_{1,m_1}\} \rangle, \dots, \langle v_\ell : T_\ell, \xi_\ell, \text{expr}_0^\ell, \{P_{\ell,1}, \dots, P_{\ell,m_\ell}\} \rangle \}$

$\text{atr}(n) := \{ C \mapsto \{v_1, \dots, v_\ell\} \}$

$\vdots$

$\langle S: D_*, +, ?, \{ordered\} \rangle$

$\Downarrow$

$\langle S: D_*, +, \text{D}, \{ordered\} \rangle$

$CD_1$

$\langle\langle \text{Stereotype}_1, \dots, \text{Stereotype}_n \rangle\rangle$

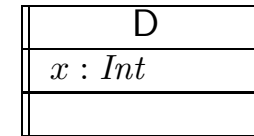
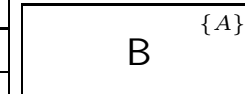
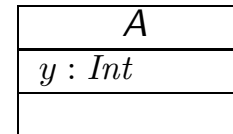
Package::C

+ r :  $C_{0,1} = \text{expr}$

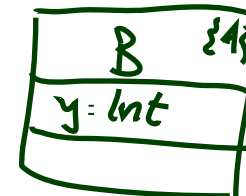
s :  $D_* \{ordered\}$

- v :  $Int = 27$

w :  $Float \{readOnly\}$



$CD_2$



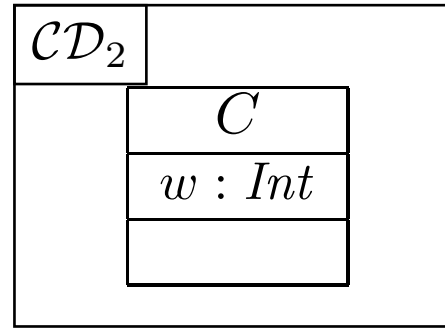
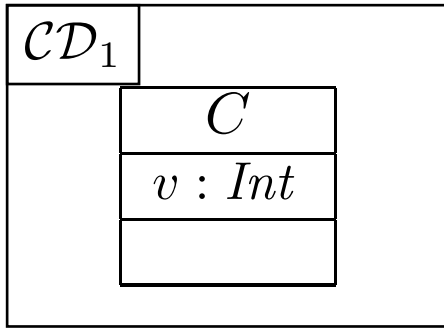
# Is the Mapping a Function?

**Question:** Is  $\mathcal{S}(\mathcal{CD})$  **well-defined**?

There are two possible **sources for problems**:

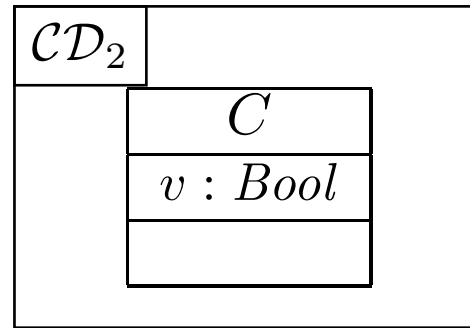
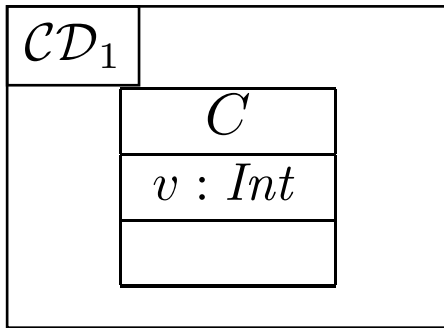
(1) A **class**  $C$  may appear in **multiple** class **diagrams**:

(i)



$$\mathcal{S} = (\{Int\}, \{C\}, \{v: Int, w: Int\}, \{C \mapsto \{v, w\}\})$$

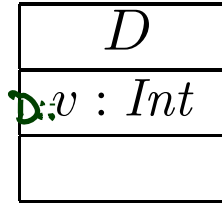
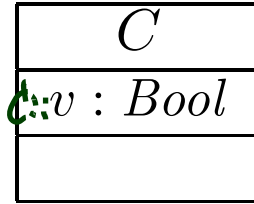
(ii)



Simply **forbid** the case (ii) — easy syntactical check on diagram.

# Is the Mapping a Function?

(2) An **attribute**  $v$  may appear in **multiple classes** with different type:



Two approaches:

$$V = (\{Bool, Int\}, \{C, D\}, \{v:Bool, v:Int\})$$

$$\{C \mapsto \{v\} \dots \text{which?}\}$$

- Require **unique** attribute names.

This requirement can easily be established (implicitly, behind the scenes) by viewing  $v$  as an abbreviation for

$$C::v \quad \text{or} \quad D::v$$

depending on the context. ( $C::v : Bool$  and  $D::v : Int$  are then unique.)

- Subtle, formalist's approach: observe that

$$\langle v : Bool, \dots \rangle \quad \text{and} \quad \langle v : Int, \dots \rangle$$

$$\langle v:Bool, +, \downarrow, \emptyset \rangle \neq \langle v:Int, +, \downarrow, \emptyset \rangle$$

are **different things** in  $V$ . We don't follow that path...

$$M: C \mapsto \{ \langle v:Bool, +, \downarrow, \emptyset \rangle \}$$

$$D \mapsto \{ \langle v:Int, +, \downarrow, \emptyset \rangle \}$$

# *Class Diagram Semantics*

# Semantics

The semantics of a set of **class diagrams**  $\mathcal{C}\mathcal{D}$  is the induced **signature**  $\mathcal{I}(\mathcal{C}\mathcal{D})$ .

The **signature** induces a set of **system states**  $\Sigma_{\mathcal{I}}^{\mathcal{D}}$  (given a **structure**  $\mathcal{D}$ ).

- Do we need to redefine/extend  $\mathcal{D}$ ? **No.**

(Would be different if we considered the definition of enumeration types in class diagrams. Then the domain of an enumeration type  $T$ , i.e. the set  $\mathcal{D}(T)$ , would be determined by the class diagram, and not free for choice.)



$\mathcal{I} = (\{\text{Color}\}, \dots$   
from diagram

$\mathcal{D}(\text{Color}) = \{\text{red}, \text{green}, \text{blue}\}$   
from diagram



# Semantics

The semantics of a set of **class diagrams**  $\mathcal{C}\mathcal{D}$  is the induced **signature**  $\mathcal{I}(\mathcal{C}\mathcal{D})$ .

The **signature** induces a set of **system states**  $\Sigma_{\mathcal{D}}^{\mathcal{D}}$  (given a **structure**  $\mathcal{D}$ ).

- Do we need to redefine/extend  $\mathcal{D}$ ? **No.**

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- What is the effect on  $\Sigma_{\mathcal{D}}^{\mathcal{D}}$ ? **Little.**

For now, we only **remove** abstract class instances, i.e.

$$\sigma : \mathcal{D}(\mathcal{C}) \rightarrow (V \rightarrow (\mathcal{D}(\mathcal{I}) \cup \mathcal{D}(\mathcal{C}_*)))$$

*abstract*

is now **only** called **system state** if and only if, for all  $\langle C, S_C, 1, t \rangle \in \mathcal{C}$ ,

$$\underbrace{\text{dom}(\sigma) \cap \mathcal{D}(C) = \emptyset.}$$

With  $a = 0$  as default “abstractness”, the earlier definitions apply directly. (We’ll revisit this when discussing inheritance.)

# What About The Rest?

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- **Classes:**
  - **Active:** not represented in  $\sigma$ .  
**Later:** relevant for behaviour, i.e., how system states evolve over time.
  - **Stereotypes:** in a minute.
- **Attributes:**
  - **Initial value expression:** not represented in  $\sigma$ .  
**Later:** provides an initial value as effect of “creation action”.
  - **Visibility:** not represented in  $\sigma$ .  
**Later:** viewed as additional **typing information** for well-formedness of actions; and with inheritance.
  - **Properties:** such as `readOnly`, `ordered`, `composite` (**Deprecated** in the standard.)
    - `readOnly` — **later** treated similar to visibility.
    - `ordered` — not considered in our UML fragment ( $\rightarrow$  sets vs. sequences).
    - `composite` — cf. lecture on associations.

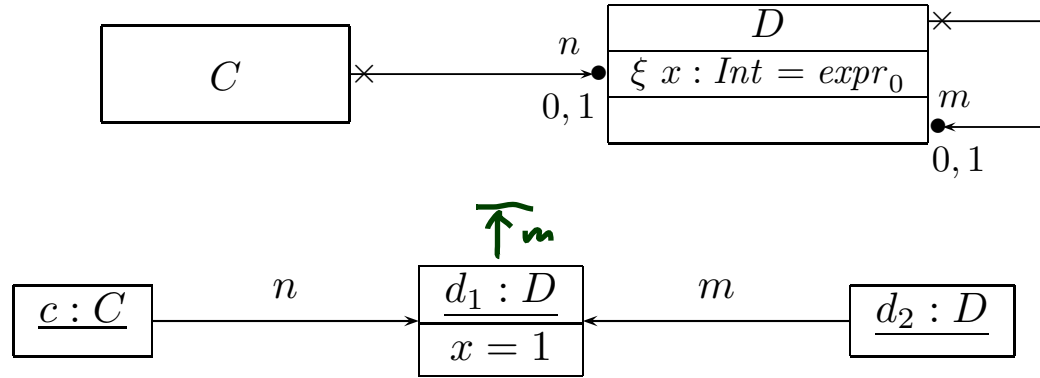
## *Rhapsody Demo I*

RECALL: SEND ME YOUR POOL-ACCOUNT NAME  
( meyerp, NOT: a124, xh102 (R2) )

# *Visibility*

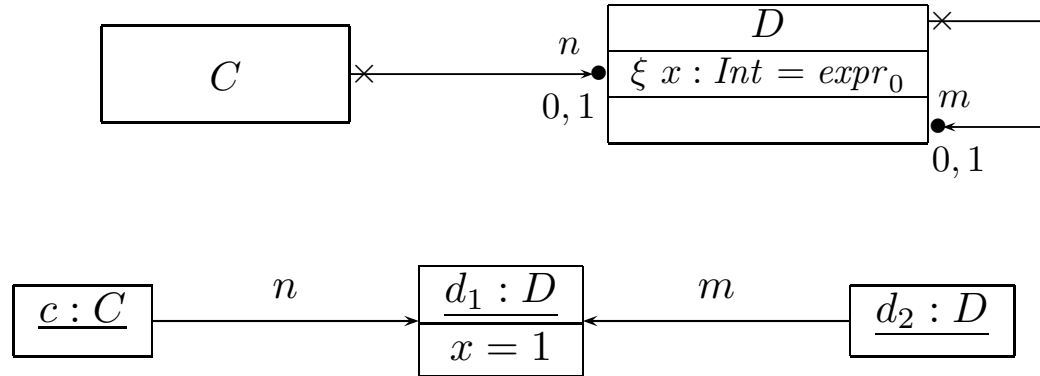
# The Intuition by Example

$$\mathcal{S} = (\{Int\}, \{C, D\}, \{n : D_{0,1}, m : D_{0,1}, \langle x : Int, \xi, expr_0, \emptyset \rangle\}, \{C \mapsto \{n\}, D \mapsto \{x, m\}\})$$



# The Intuition by Example

$$\mathcal{S} = (\{Int\}, \{C, D\}, \{n : D_{0,1}, m : D_{0,1}, \langle x : Int, \xi, expr_0, \emptyset \rangle\}, \{C \mapsto \{n\}, D \mapsto \{x, m\}\})$$



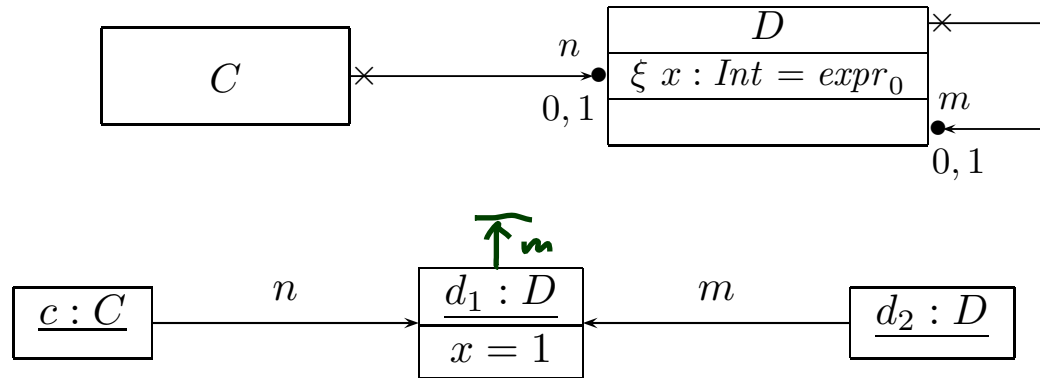
Assume  $w_1 : \tau_C$  and  $w_2 : \tau_D$  are logical variables.

Which of the following **syntactically correct** (?) OCL expressions **should** we consider to be **well-typed**?

| $\xi$ of $x$ :    | public | private | protected | package |
|-------------------|--------|---------|-----------|---------|
| $w_1 . n . x = 0$ |        |         | later     | not     |
| $w_2 . m . x = 0$ |        |         | later     | not     |

# The Intuition by Example

$$\mathcal{S} = (\{Int\}, \{C, D\}, \{n : D_{0,1}, m : D_{0,1}, \langle x : Int, \xi, expr_0, \emptyset \rangle\}, \{C \mapsto \{n\}, D \mapsto \{x, m\}\})$$



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Which of the following **syntactically correct** (?) OCL expressions **should** we consider to be **well-typed**?

| $\xi$ of $x$ :    | <u>public</u>                          | private                                     | protected | package |
|-------------------|--|---|-----------|---------|
| $w_1 . n . x = 0$ | ✓ <del>HT</del><br>✗ <del>‡</del><br>? | ✓ <del>‡</del><br>✗ <del>   </del><br>?     | later     | not     |
| $w_2 . m . x = 0$ | ✓ <del>HT</del><br>✗<br>?              | ✓ <del>HT  </del><br>✗ <del>HT  </del><br>? | later     | not     |

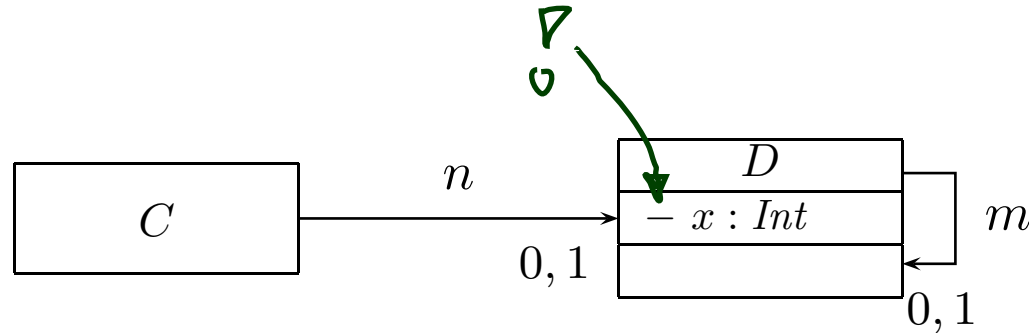
by class (OCL, C++, Java, ...)

by object

# Context

$$\mathcal{S} = (\{\text{Int}\}, \{C, D\}, \{n : D_{0,1}, m : D_{0,1}, \langle x : \text{Int}, \xi, \text{expr}_0, \emptyset \rangle\}, \{C \mapsto \{n\}, D \mapsto \{x, m\}\})$$

- By example:



$$\underbrace{\text{self}_D . x}_{: \tau_D} > 0 \quad \checkmark$$

$$\underbrace{\text{self}_D . m . x}_{: \tau_D} > 0 \quad \checkmark$$

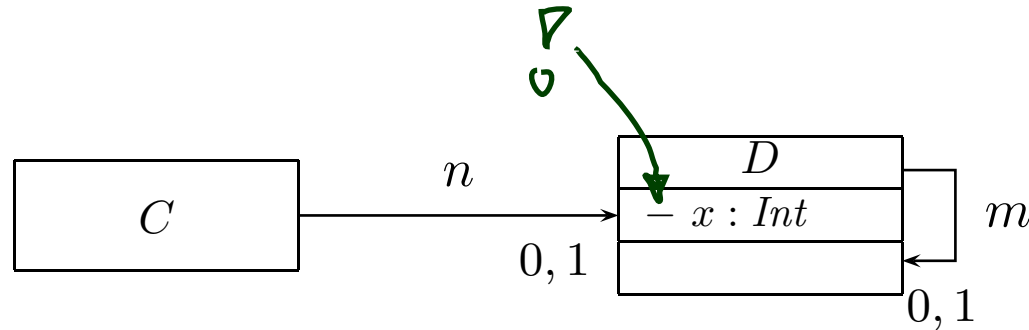
$$\underbrace{\text{self}_C . n . x}_{: \tau_C} > 0 \quad \times$$



# Context

$$\mathcal{S} = (\{\text{Int}\}, \{C, D\}, \{n : D_{0,1}, m : D_{0,1}, \langle x : \text{Int}, \xi, \text{expr}_0, \emptyset \rangle\}, \{C \mapsto \{n\}, D \mapsto \{x, m\}\})$$

- By example:



$$\underbrace{\text{self}_D . x}_{\tau_D} > 0 \quad \checkmark$$

$$\underbrace{\text{self}_D . m . x}_{\tau_D} > 0 \quad \checkmark$$

$$\underbrace{\text{self}_C . n . x}_{\tau_C} > 0 \quad \times$$

- That is, whether an expression involving attributes with visibility is well-typed **depends** on the class of objects for which it is evaluated.
- Visibility is **'by class'** — **not** 'by object'.

# Attribute Access in Context

**Recall:** attribute access in OCL Expressions,  $C, D \in \mathcal{C}$ .

- |               |  |   |
|---------------|--|---|
| $v(expr_1)$   | : $\tau_C \rightarrow \cancel{\tau_C} T$ | • $v : T \in atr(C), T \in \mathcal{T}$ , |
| $r_1(expr_1)$ | : $\tau_C \rightarrow \tau_D$            | • $r_1 : D_{0,1} \in atr(C)$ ,            |
| $r_2(expr_1)$ | : $\tau_C \rightarrow Set(\tau_D)$       | • $r_2 : D_* \in atr(C)$ ,                |

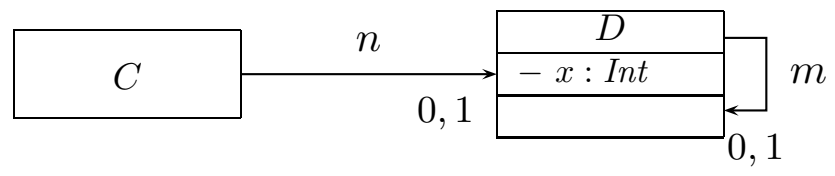
## New rules for well-typedness **considering visibility:**

- |                    |                                    |   |
|--------------------|------------------------------------|---|
| • $v(w)$           | : $\tau_C \rightarrow T$           | $w : \tau_C, v : T \in atr(C), T \in \mathcal{T}$   |
| • $r_1(w)$         | : $\tau_C \rightarrow \tau_D$      | $w : \tau_C, r_1 : D_{0,1} \in atr(C)$  |
| • $r_2(w)$         | : $\tau_C \rightarrow Set(\tau_D)$ | $w : \tau_C, r_1 : D_* \in atr(C)$  |
| • $v(expr_1(w))$   | : $\tau_C \rightarrow T$           | $\langle v : T, \xi, expr_0, P \rangle \in atr(C), T \in \mathcal{T},$<br>$\underbrace{expr_1(w) : \tau_C}, \underbrace{w : \tau_{C_1} \text{ and } C_1 = C}, \text{ or } \underbrace{\xi = +}$ |
| • $r_1(expr_1(w))$ | : $\tau_C \rightarrow \tau_D$      | $\langle r_1 : D_{0,1}, \xi, expr_0, P \rangle \in atr(C),$<br>$expr_1(w) : \tau_C, w : \tau_{C_1} \text{ and } C_1 = C, \text{ or } \xi = +$   |
| • $r_2(expr_1(w))$ | : $\tau_C \rightarrow Set(\tau_D)$ | $\langle r_2 : D_*, \xi, expr_0, P \rangle \in atr(C),$<br>$expr_1(w) : \tau_C, w : \tau_{C_1} \text{ and } C_1 = C, \text{ or } \xi = +$   |

# Example

- (i)  $v(w) : \tau_C \rightarrow T$   $w : \tau_C, v : T \in atr(C), T \in \mathcal{T}$
- (ii)  $r_1(w) : \tau_C \rightarrow \tau_D$   $w : \tau_C, r_1 : D_{0,1} \in atr(C)$
- (iii)  $v(expr_1(w))$   $\langle v : T, \xi, expr_0, P \rangle \in atr(C), T \in \mathcal{T},$   
 $expr_1(w) : \tau_C, w : \tau_{C_1}$  and  $C_1 = C,$  or  $\xi = +$
- (iv)  $r_1(expr_1(w))$   $\langle r_1 : D_{0,1}, \xi, expr_0, P \rangle \in atr(C),$   
 $expr_1(w) : \tau_C, w : \tau_{C_1}$  and  $C_1 = C,$  or  $\xi = +$

$v(r_1^1(\dots(r_1^n(w))))$



- $self_D . x > 0 \rightsquigarrow x(self_D) > 0$  ok, by (i)
  - $self_D . m . x > 0 \rightsquigarrow x(m(self_D)) > 0$  ok, by ~~(ii)~~, (iii)
  - $self_C . n . x > 0 \rightsquigarrow x(n(self_C)) > 0$  not ok by (iii) [and (i), (ii), and (iv) obviously]
- $: \tau_D$   
 $\underbrace{\hspace{2em}}$   
 $m(self_D)$  ok by (iii)  
 $\underbrace{\hspace{2em}}$   
 $: \tau_D$   
 $\underbrace{\hspace{2em}}$   
 $n(self_C)$   $\neq$  and  
 $\underbrace{\hspace{2em}}$   
 $: \tau_C$   $\neq$

# The Semantics of Visibility

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- **Observation:**

- Whether an expression **does** or **does not** respect visibility is **a matter of well-typedness only**.

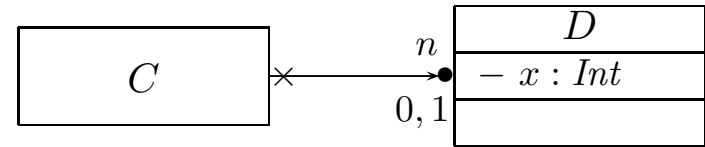
- We only evaluate (= apply  $I$  to) **well-typed** expressions.

→ We **need not** adjust the interpretation function  $I$  to support visibility.

Just decide: should we take visibility into account yes / no, and check well-typedness by the new / old rules.

# What is Visibility Good For?

- Visibility is a property of attributes — is it useful to consider it in OCL?



- In other words: given the diagram above, **is it useful** to state the following invariant (even though  $x$  is private in  $D$ )

context  $C$  inv :  $n.x > 0$  ?

**It depends.**

(cf. [OMG \(2006\)](#), Sect. 12 and 9.2.2)

- **Constraints and pre/post conditions:**

- Visibility is **sometimes not** taken into account. To state “global” requirements, it may be adequate to have a “global view”, i.e. be able to “look into” all objects.
- But: visibility supports “narrow interfaces”, “information hiding”, and similar **good design practices**. To be more robust against changes, try to state requirements only in the terms which are visible to a class.

**Rule-of-thumb:** if attributes are important to state requirements on design models, leave them public or provide get-methods (later).

- **Guards and operation bodies:**

- If in doubt, **yes** (= do take visibility into account).

Any so-called **action language** typically takes visibility into account.

# *Stereotypes*

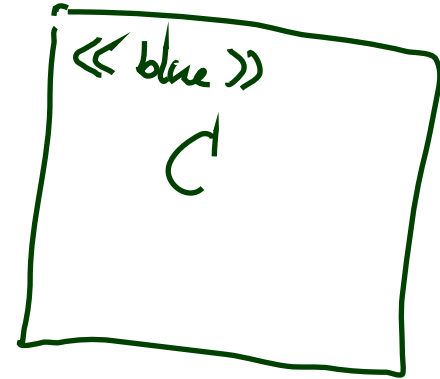
# Stereotypes as Labels or Tags

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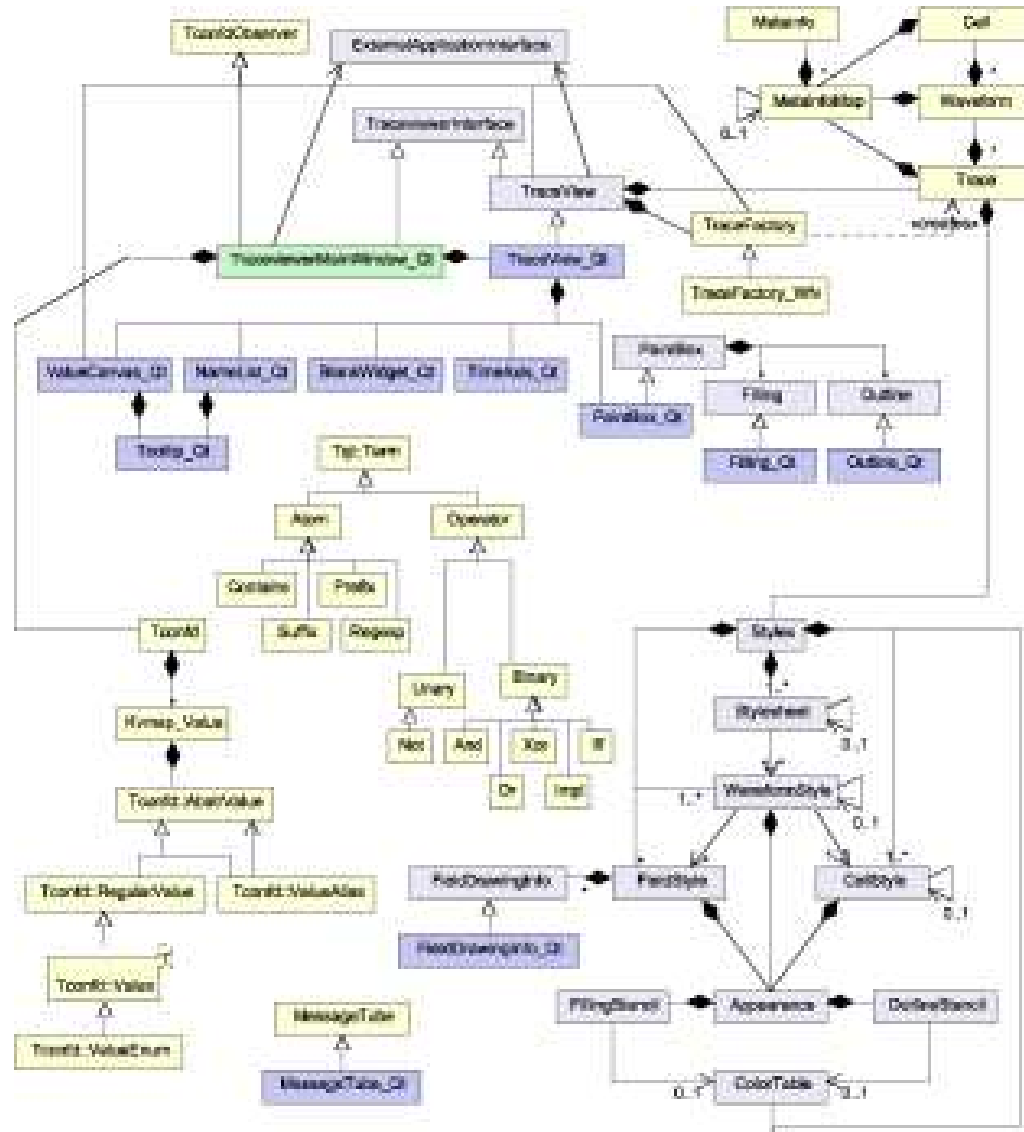
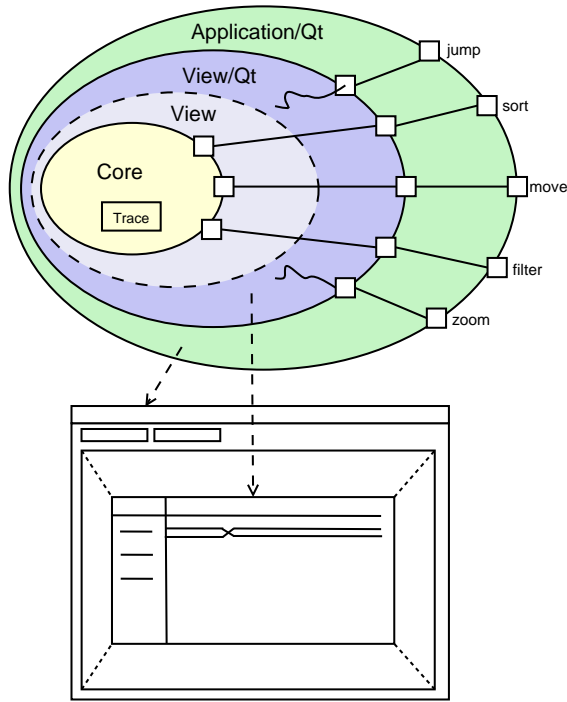
- What are Stereotypes?
  - **Not** represented in system states.
  - **Not** contributing to typing rules / well-formedness.
- Oestereich (2006):

View stereotypes as (additional) “**labelling**” (“tags”) or as “**grouping**”.
- Useful for documentation and model-driven development, e.g. code-generation:
  - **Documentation**: e.g. layers of an architecture.

Sometimes, packages (cf. [OMG \(2011a,b\)](#)) are sufficient and “right”.
  - **Model Driven Architecture (MDA)**: **later**.



# Example: Stereotypes for Documentation



- **Example:** Timing Diagram Viewer Schumann et al. (2008)
- Architecture has four layers:
  - core, data layer
  - abstract view layer
  - toolkit-specific view layer/widget
  - application using widget

Stereotype “=” layer “=” colour.



## Other Examples

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- Use stereotypes 'Team<sub>1</sub>', 'Team<sub>2</sub>', 'Team<sub>3</sub>' and assign stereotype Team<sub>*i*</sub> to class *C* if Team<sub>*i*</sub> is responsible for class *C*.
- Use stereotypes to label classes with licensing information (e.g., LGPL vs. proprietary).
- Use stereotypes 'Server<sub>A</sub>', 'Server<sub>B</sub>' to indicate where objects should be stored.
- Use stereotypes to label classes with states in the development process like “under development”, “submitted for testing”, “accepted”.
- etc. etc.

**Necessary:** a **common idea** of what each stereotype stands for.

(To be defined / agreed on by the team, not the job of the UML consortium.)

# *References*

# References

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OMG (2011a). Unified modeling language: Infrastructure, version 2.4.1. Technical Report formal/2011-08-05.

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