

Software Design, Modelling and Analysis in UML

Lecture 9: Class Diagrams IV

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Contents & Goals

Last Lecture:

- Associations syntax and semantics.
- Associations in OCL syntax.

This Lecture:

- **Educational Objectives:** Capabilities for following tasks/questions.
 - Compute the value of a given OCL constraint in a system state with links.
 - How did we treat “multiplicity” semantically?
 - What does “navigability”, “ownership”, . . . mean?
 - ...
- **Content:**
 - Associations and OCL: semantics.
 - Associations: the rest.

Associations and OCL Cont'd

Recall: Associations and OCL Syntax

Recall: OCL syntax as introduced in Lecture 3, interesting part:

$$\begin{array}{l} \text{expr} ::= \dots \quad | \quad r_1(\text{expr}_1) : \tau_C \rightarrow \tau_D \\ \quad | \quad r_2(\text{expr}_1) : \tau_C \rightarrow \text{Set}(\tau_D) \end{array} \quad \begin{array}{l} r_1 : D_{0,1} \in \text{atr}(C) \\ r_2 : D_* \in \text{atr}(C) \end{array}$$

Now becomes

$$\begin{array}{l} \text{expr} ::= \dots \quad | \quad \text{role}(\text{expr}_1) : \tau_C \rightarrow \tau_D \\ \quad | \quad \text{role}(\text{expr}_1) : \tau_C \rightarrow \text{Set}(\tau_D) \end{array} \quad \begin{array}{l} \mu = 0..1 \text{ or } \mu = 1..1 \\ \text{otherwise} \end{array}$$

if there is

$$\begin{aligned} & \langle r : \dots, \langle \text{role} : D, \mu, -, -, -, - \rangle, \dots, \langle \text{role}' : C, -, -, -, -, - \rangle, \dots \rangle \in V \text{ or} \\ & \langle r : \dots, \langle \text{role}' : C, -, -, -, -, - \rangle, \dots, \langle \text{role} : D, \mu, -, -, -, - \rangle, \dots \rangle \in V, \quad \text{role} \neq \text{role}'. \end{aligned}$$

Note:

- Association name as such **does not occur** in OCL syntax, role names do.
- expr_1 has to denote an object of a class which “participates” in the association.

OCL and Associations: Semantics

Recall:

Assume $expr_1 : \tau_C$ for some $C \in \mathcal{C}$. Set $u_1 := I[\![expr_1]\!](\sigma, \beta) \in \mathcal{D}(T_C)$.

- $I[\![r_1(expr_1)]\!](\sigma, \beta) := \begin{cases} u & , \text{ if } u_1 \in \text{dom}(\sigma) \text{ and } \sigma(u_1)(r_1) = \{u\} \\ \perp & , \text{ otherwise} \end{cases}$
- $I[\![r_2(expr_1)]\!](\sigma, \beta) := \begin{cases} \sigma(u_1)(r_2) & , \text{ if } u_1 \in \text{dom}(\sigma) \\ \perp & , \text{ otherwise} \end{cases}$

Now needed:

$$I[\![role(expr_1)]\!]((\sigma, \lambda), \beta)$$

- We cannot simply write $\sigma(u)(role)$.

Recall: $role$ is (**for the moment**) not an attribute of object u (not in $atr(C)$).

- What we have is $\lambda(r)$ (with association name r , not with role name $role$!).

$$\langle r : \dots, \langle role : D, \mu, -, -, -, - \rangle, \dots, \langle role' : C, -, -, -, -, - \rangle, \dots \rangle$$

But it yields a set of n -tuples, of which **some** relate u and some instances of D .

- $role$ denotes the position of the D 's in the tuples constituting the value of r .

OCL and Associations: Semantics Cont'd

Assume $expr_1 : \tau_C$ for some $C \in \mathcal{C}$. Set $u_1 := I[\![expr_1]\!](\sigma, \lambda), \beta) \in \mathcal{D}(T_C)$.

- $I[\![role(expr_1)]\!](\sigma, \lambda), \beta) := \begin{cases} u & , \text{ if } u_1 \in \text{dom}(\sigma) \text{ and } L(role)(u_1, \lambda) = \{u\} \\ \perp & , \text{ otherwise} \end{cases}$
- $I[\![role(expr_1)]\!](\sigma, \lambda), \beta) := \begin{cases} L(role)(u_1, \lambda) & , \text{ if } u_1 \in \text{dom}(\sigma) \\ \perp & , \text{ otherwise} \end{cases}$

where

$$L(role)(u, \lambda) = \{(u_1, \dots, u_n) \in \lambda(r) \mid u \in \{u_1, \dots, u_n\}\} \downarrow i$$

if

$$\langle r : \langle role_1 : _, _, _, _, _, _ \rangle, \dots \langle role_n : _, _, _, _, _, _ \rangle, \rangle, \quad \underbrace{role}_{\text{---}} = \underbrace{role_i}_{\text{---}}.$$

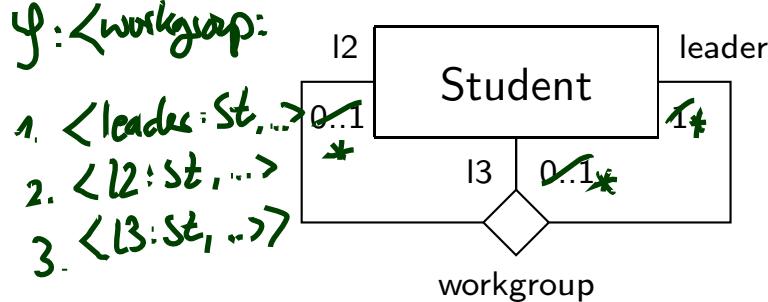
Given a set of n -tuples A ,

$A \downarrow i$ denotes the element-wise projection onto the i -th component.

OCL and Associations Semantics: Example

- $I[\![role(expr_1)]\!]((\sigma, \lambda), \beta) := \begin{cases} u & , \text{ if } u_1 \in \text{dom}(\sigma) \text{ and } L(role)(u_1, \lambda) = \{u\} \\ \perp & , \text{ otherwise} \end{cases}$
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$$L(role)(u, \lambda) = \{(u_1, \dots, u_n) \in \lambda(r) \mid u \in \{u_1, \dots, u_n\}\} \downarrow i$$



$$\lambda(\text{workgroup}) = \{(1_S, 2_S, 3_S), (1_S, 3_S, 5_S), (5_S, 1_S, 1_S)\}$$

$S_S \in \text{dom } \sigma$

$$\text{allInstances}_{\text{Student}} \rightarrow \text{Exists}(s \mid s.\text{I2} = s.\text{I3})$$

$$I[\![s.\text{I2}]\!](\sigma, \lambda, \{s \mapsto S_S\}) = \{(1_S, 3_S, 5_S), (5_S, 1_S, 1_S)\} \downarrow 2$$

$$= \{3_S, 1_S\} = R_1$$

$$I[\![s.\text{I3}]\!](\sigma, \lambda, \beta) = \{(1_S, 3_S, 5_S), (5_S, 1_S, 1_S)\} \downarrow 3 = \{5_S, 1_S\} = R_2$$

$$\begin{aligned} I[\![s.\text{I2}]\!](\sigma, \lambda, \{s \mapsto S_S\}) &= I[\!=\!](R_1, R_2) \\ &= \text{false} \end{aligned}$$

Associations: The Rest

The Rest

Recapitulation: Consider the following association:

$$\langle r : \langle role_1 : C_1, \mu_1, P_1, \xi_1, \nu_1, o_1 \rangle, \dots, \langle role_n : C_n, \mu_n, P_n, \xi_n, \nu_n, o_n \rangle \rangle$$

- **Association name** r and **role names / types** $role_i / C_i$ induce extended system states (σ, λ) .
- **Multiplicity** μ is considered in OCL syntax.
- **Visibility** ξ / **Navigability** ν : well-typedness (in a minute).

Now the rest:

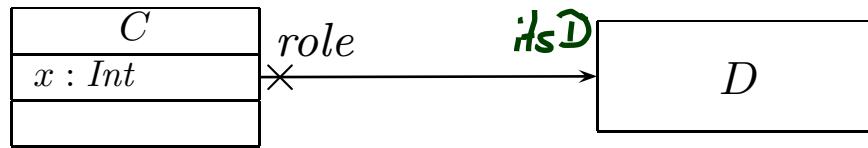
- **Multiplicity** μ : we propose to view them as constraints.
- **Properties** P_i : even more typing.
- **Ownership** o : getting closer to pointers/references.
- **Diamonds**: exercise.

Navigability

Navigability is treated similar to visibility:

Using names of non-navigable association ends ($\nu = \times$) are **forbidden**.

Example: Given



the following OCL expression is **not well-typed** wrt. navigability,

context D inv : $role.x > 0$

The standard says: navigation is...

- '—': ...possible
- '×': ...not possible
- '>': ...efficient

So: In general, UML associations **are different** from pointers / references in general!

But: Pointers / references **can faithfully** be modelled by UML associations.

Multiplicities as Constraints

Recall: Multiplicity is a term of the form $N_1..N_2, \dots, N_{2k-1}..N_{2k}$ where $N_i \leq N_{i+1}$ for $1 \leq i \leq 2k$, $N_1, \dots, N_{2k-1} \in \mathbb{N}$, $N_{2k} \in \mathbb{N} \cup \{\ast\}$.

Define $\mu_{\text{OCL}}^C(role) :=$



context $C \text{ inv} : (N_1 \leq role \rightarrow \text{size}() \leq N_2) \text{ or } \dots \text{ or } (N_{2k-1} \leq role \rightarrow \underbrace{\text{size}()} \leq N_{2k})$
omit if $N_{2k} = \ast$

for each $\langle r : \dots, \langle role : D, \mu, _, _, _, _, _ \rangle, \dots, \langle role' : C, _, _, _, _, _ \rangle, \dots \rangle \in V$ or

$\langle r : \dots, \langle role' : C, _, _, _, _, _ \rangle, \dots, \langle role : D, \mu, _, _, _, _ \rangle, \dots \rangle \in V,$

with $role \neq role'$, if $\mu \neq 0..1$, $\mu \neq 1..1$, and

$\mu_{\text{OCL}}^C(role) := \text{context } C \text{ inv} : \text{not}(\text{ocllsUndefined}(role))$

if $\mu = 1..1$.

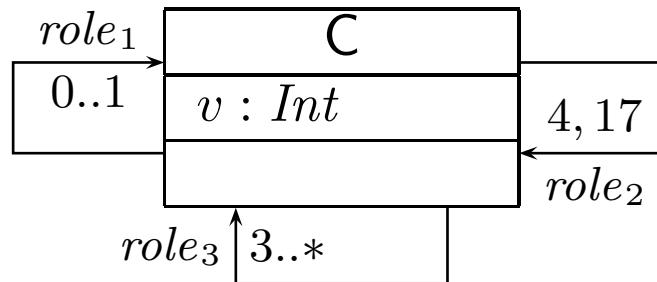
Note: in n -ary associations with $n > 2$, there is redundancy.

Multiplicities as Constraints Example

$\mu_{OCL}^C(role) = \text{context } C \text{ inv :}$

$$(N_1 \leq role \rightarrow \text{size}() \leq N_2) \text{ or } \dots \text{ or } (N_{2k-1} \leq role \rightarrow \text{size}() \leq N_{2k})$$

$\mathcal{CD} :$

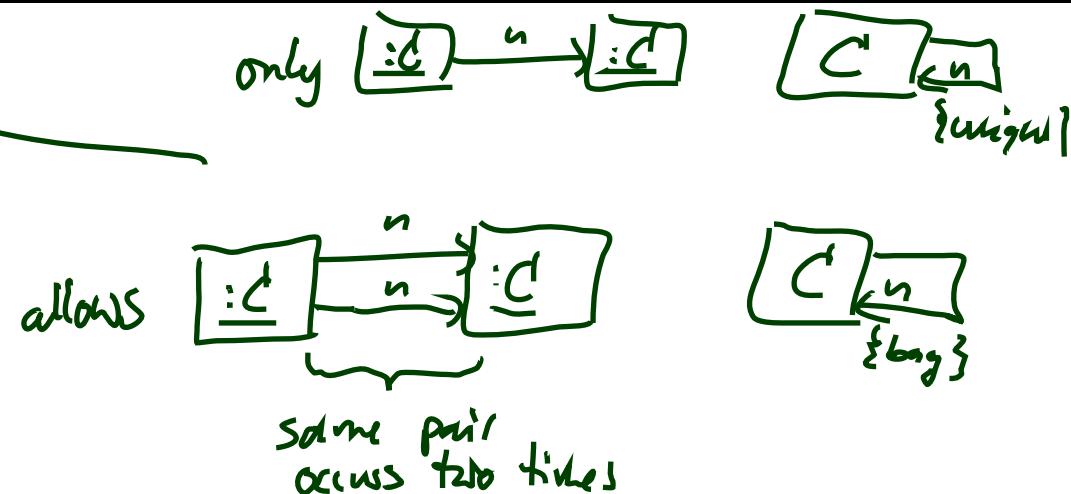


- $\mu_{OCL}^C(role_2) = \text{context } C \text{ inv : } 4 \leq role_2 \rightarrow \text{size} \leq 4 \text{ or } 17 \leq role_2 \rightarrow \text{size} \leq 17$
= {context C inv : $role_2 \rightarrow \text{size}() = 4$ or $role_2 \rightarrow \text{size}() = 17$ }
- $\mu_{OCL}^C(role_3) = \text{context } C \text{ inv : } 3 \leq role_3 \rightarrow \text{size}$

Properties

We don't want to cover association **properties** in detail,
only some observations (assume binary associations):

Property	Intuition	Semantical Effect
unique	one object has at most one r -link to a single other object	current setting
bag	one object may have multiple r -links to a single other object	have $\lambda(r)$ yield multi-sets
ordered, sequence	an r -link is a sequence of object identities (possibly including duplicates)	have $\lambda(r)$ yield sequences



Properties

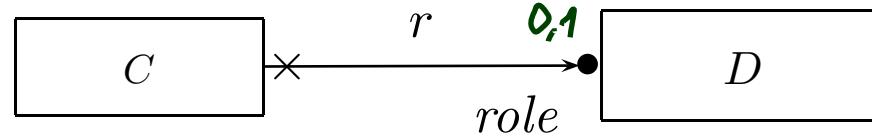
We don't want to cover association **properties** in detail,
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Property	OCL Typing of expression $role(expr)$
unique	$T_D \rightarrow Set(T_C)$
bag	$T_D \rightarrow Bag(T_C)$
ordered, sequence	$T_D \rightarrow Seq(T_C)$

For **subsets**, **redefines**, **union**, etc. see (? , 127).

Ownership



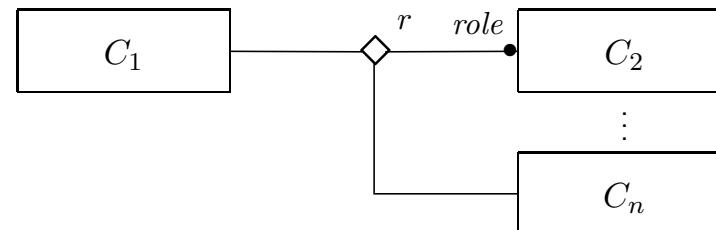
Intuitively it says:

Association r is **not a “thing on its own”** (i.e. provided by λ),
but association end ‘ $role$ ’ is **owned** by C (!).
(That is, it’s stored inside C object and provided by σ).

So: if multiplicity of $role$ is $0..1$ or $1..1$, then the picture above is very close to concepts of pointers/references.

Actually, ownership is seldom seen in UML diagrams. Again: if target platform is clear, one may well live without (cf. (OMG, 2011b, 42) for more details).

Not clear to me:



Back to the Main Track

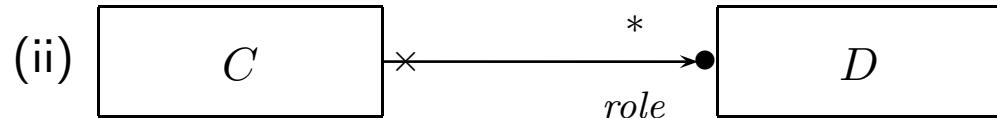
Back to the main track:

Recall: on some earlier slides we said, the extension of the signature is **only** to study associations in “full beauty”.

For the remainder of the course, we should look for something simpler...

Proposal:

- **from now on**, we only use associations of the form



(And we may omit the non-navigability and ownership symbols.)

- Form (i) introduces $\text{role} : C_{0,1}$, and form (ii) introduces $\text{role} : C_*$ in V .
- In both cases, $\text{role} \in \text{atr}(C)$.
- We drop λ and go back to our nice σ with $\sigma(u)(\text{role}) \subseteq \mathcal{D}(D)$.

OCL Constraints in (Class) Diagrams

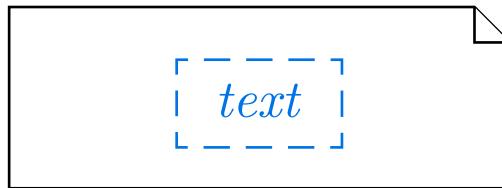
Where Shall We Put OCL Constraints?

Two options:

- (i) Notes.
- (ii) Particular dedicated places.

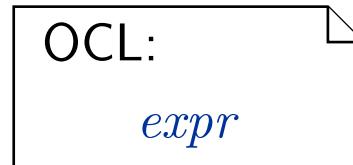
(i) **Notes:**

A UML **note** is a picture of the form

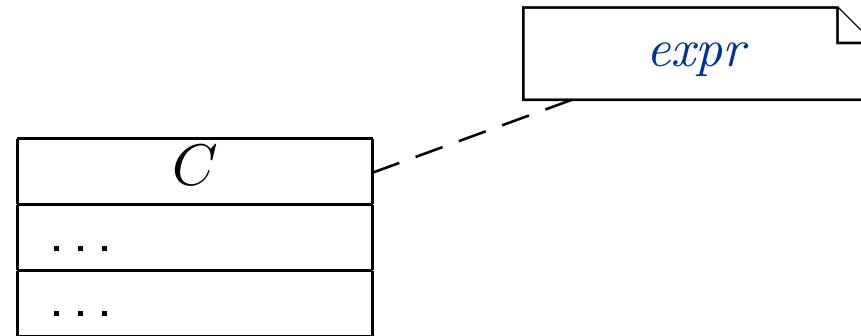


text can principally be **everything**, in particular **comments** and **constraints**.

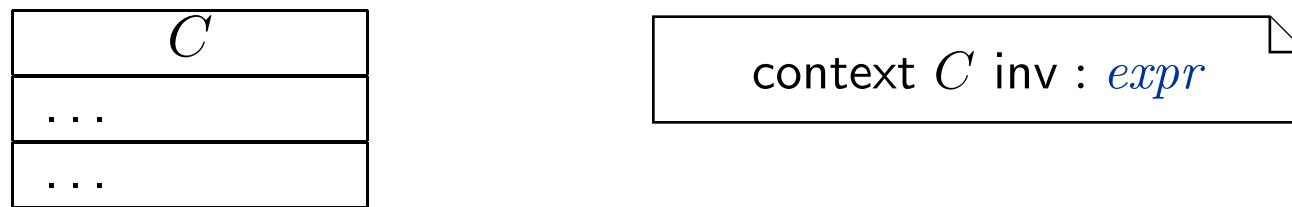
Sometimes, content is **explicitly classified** for clarity:



OCL in Notes: Conventions



stands for

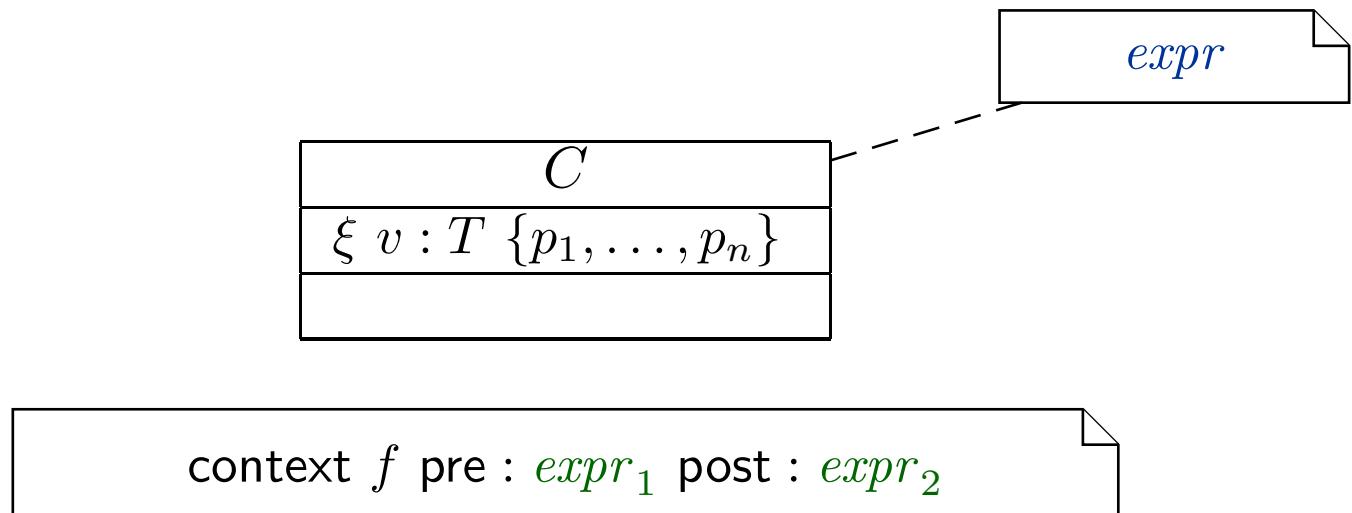


Where Shall We Put OCL Constraints?

- (ii) **Particular dedicated places** in class diagrams: (behavioural features: later)

C
$\xi v : T \{p_1, \dots, p_n\} \{expr\}$
$\xi f(v_1 : T, \dots, v_n : T_n) : T \{p_1, \dots, p_n\} \{pre : expr_1$ $\quad \quad \quad post : expr_2\}$

For simplicity, we view the above as an abbreviation for



Invariants of a Class Diagram

- Let \mathcal{CD} be a class diagram.
- We are (now) able to recognise OCL constraints when we see them, so define

$$Inv(\mathcal{CD})$$

as the set $\{\varphi_1, \dots, \varphi_n\}$ of OCL constraints **occurring** in notes in \mathcal{CD} — after **unfolding** all **graphical** abbreviations (cf. previous slides).

- **As usual:** consider all invariants in all notes in any class diagram — plus implicit multiplicity-induced invariants.

$$Inv(\mathcal{CD}) = \bigcup_{\mathcal{CD} \in \mathcal{CD}} Inv(\mathcal{CD}) \cup$$

$$\left\{ \mu_{OCL}^C(role) \mid \langle r : \dots, \langle role : D, \mu, -, -, -, - \rangle, \dots, \langle role' : C, -, -, -, -, - \rangle, \dots \rangle \in V \text{ or} \right. \\ \left. \langle r : \dots, \langle role' : C, -, -, -, -, - \rangle, \dots, \langle role : D, \mu, -, -, -, - \rangle, \dots \rangle \in V \right\}.$$

- **Analogously:** $Inv(\cdot)$ for any kind of diagram (like **state machine diagrams**).

Semantics of a Class Diagram

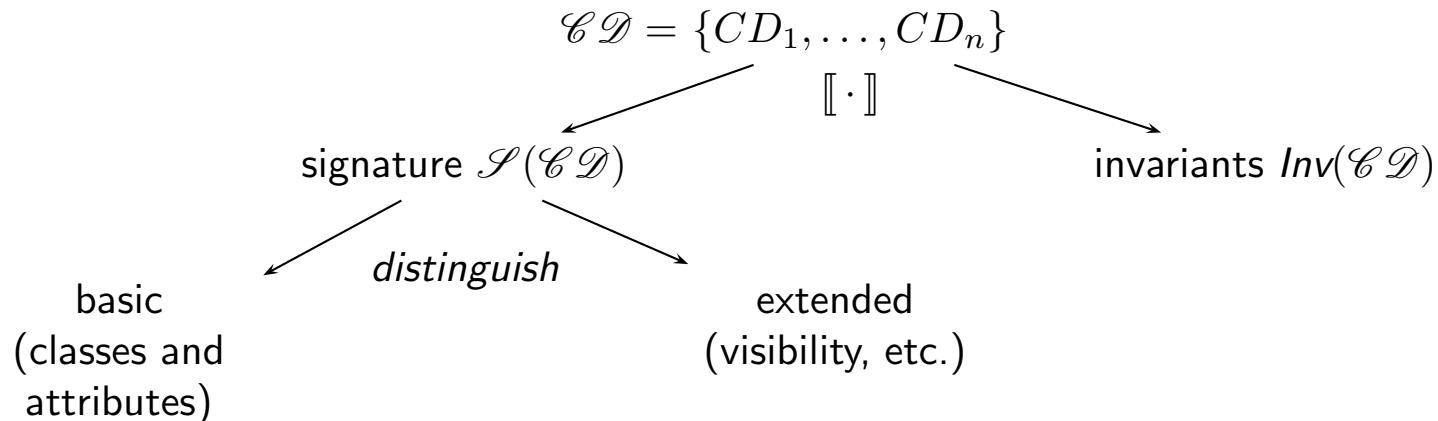
Definition. Let \mathcal{CD} be a set of class diagrams.

We say, the **semantics** of \mathcal{CD} is the signature it induces and the set of OCL constraints occurring in \mathcal{CD} , denoted

$$[\![\mathcal{CD}]\!] := \langle \mathcal{S}(\mathcal{CD}), \text{Inv}(\mathcal{CD}) \rangle.$$

Given a structure \mathcal{D} of \mathcal{S} (and thus of \mathcal{CD}), the class diagrams **describe** the system states $\Sigma_{\mathcal{S}}^{\mathcal{D}}$, of which **some** may satisfy $\text{Inv}(\mathcal{CD})$.

In pictures:



Pragmatics

Recall: a UML **model** is an image or pre-image of a software system.

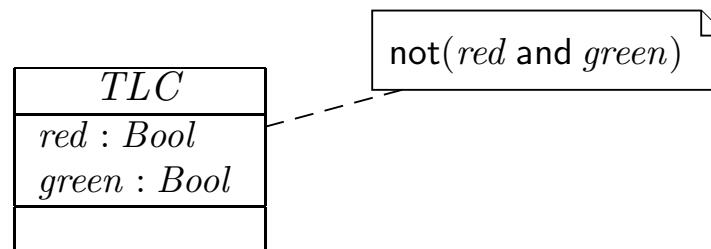
A set of class diagrams \mathcal{CD} describes the **structure** of system states.

Together with the invariants $Inv(\mathcal{CD})$ it can be used to state:

- **Pre-image:** Dear programmer, please provide an implementation which **uses** only system states that satisfy $Inv(\mathcal{CD})$.
- **Post-image:** Dear user/maintainer, in the existing system, only system states which satisfy $Inv(\mathcal{CD})$ are **used**.

(The exact meaning of “**use**” will become clear when we study behaviour — intuitively: the system states that are reachable from the initial system state(s) by calling methods or firing transitions in state-machines.)

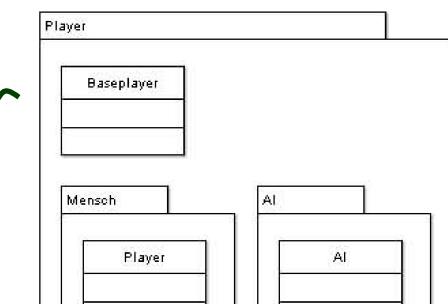
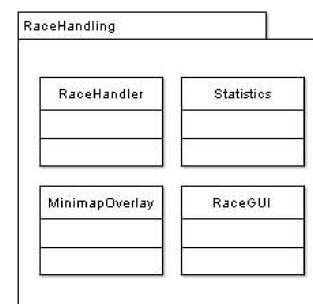
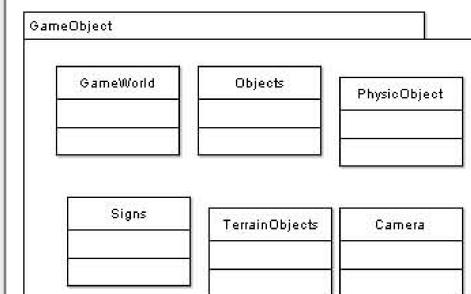
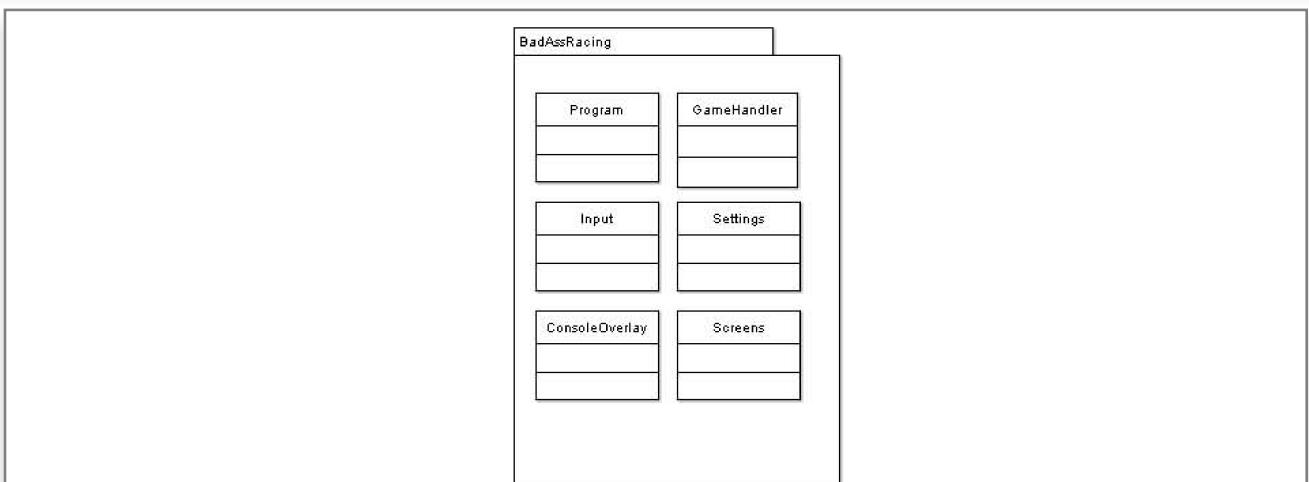
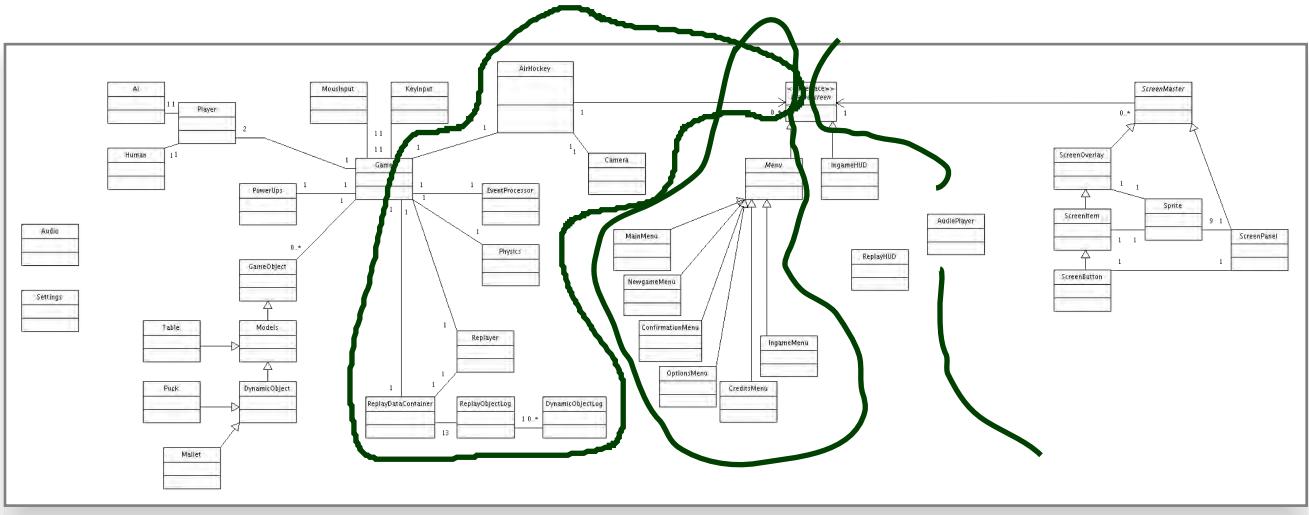
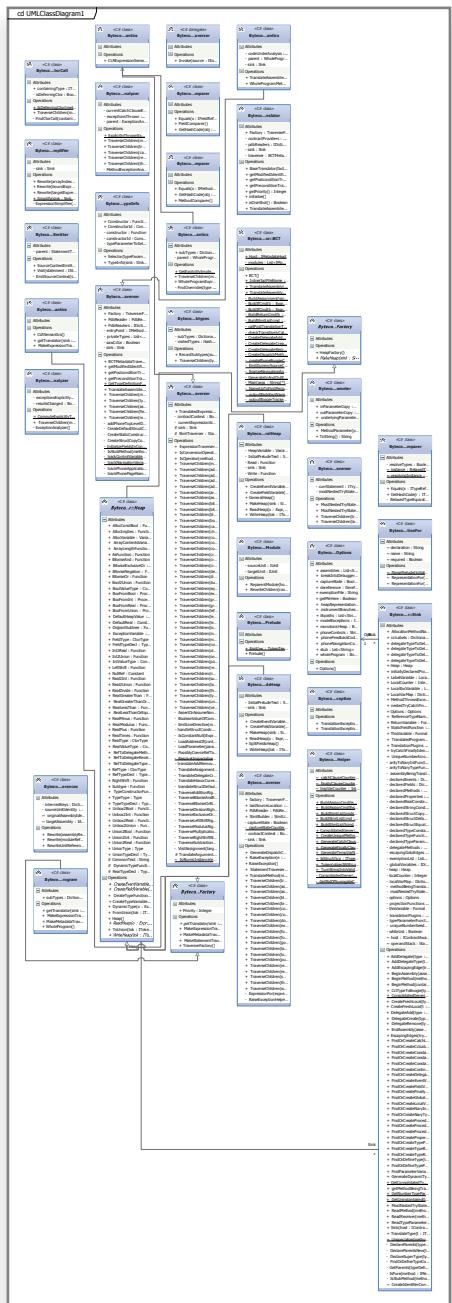
Example: highly abstract model of traffic lights controller.



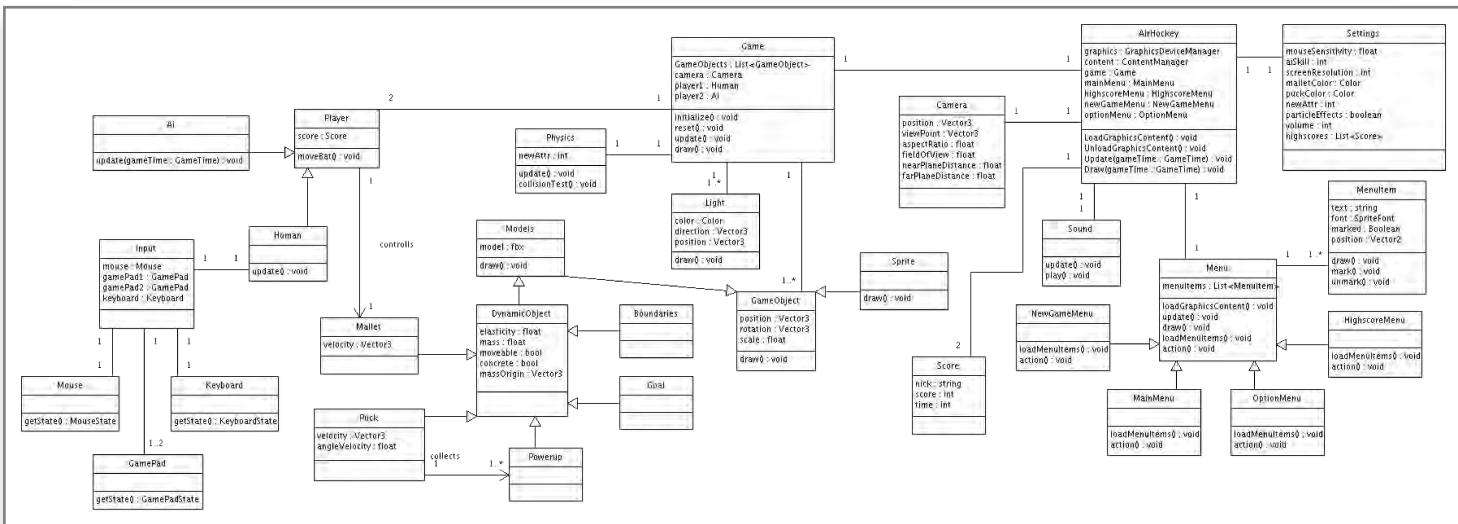
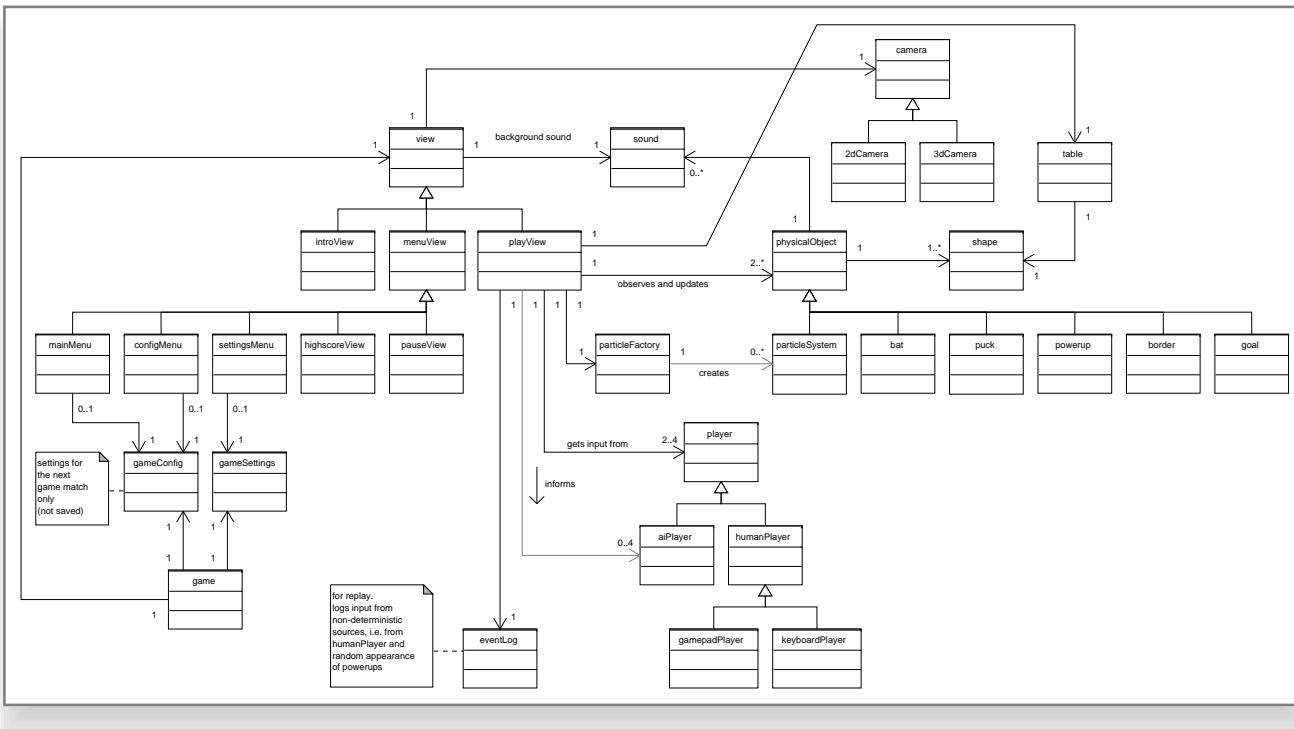
Design Guidelines for (Class) Diagram

(partly following Ambler (2005))

Some Example Class Diagrams



Some More Example Class Diagrams



So: what makes a class diagram a **good class diagram?**

Main and General Modelling Guideline

Be good to your audience.

“Imagine you’re given **your** diagram \mathcal{D} and asked to conduct task \mathcal{T} .

- Can you do \mathcal{T} with \mathcal{D} ?
(semantics sufficiently clear? all necessary information available? ...)
- Does doing \mathcal{T} with \mathcal{D} cost you more nerves/time/money/... than it should?”
(syntactical well-formedness? readability? intention of deviations from standard syntax clear? reasonable selection of information? layout? ...)

In other words:

- the things **most relevant** for task \mathcal{T} , do they **stand out** in \mathcal{D} ?
- the things **less relevant** for task \mathcal{T} , do they **disturb** in \mathcal{D} ?

Main and General Quality Criterion

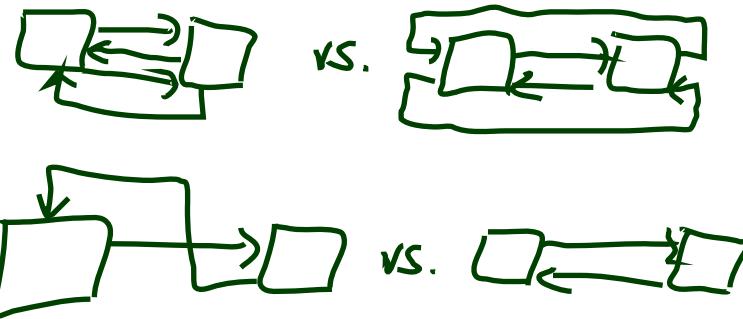
- **Q:** When is a (class) diagram a good diagram?
- **A:** If it serves its purpose/makes its point.

Examples for purposes and points and rules-of-thumb:

- **Analysis/Design**
 - realizable, no contradictions
 - abstract, focused, admitting degrees of freedom for (more detailed) design
 - platform independent – as far as possible but not (artificially) farer
- **Implementation/A**
 - close to target platform
($C_{0,1}$ is easy for Java, C_* comes at a cost — other way round for RDB)
- **Implementation/B**
 - complete, executable
- **Documentation**
 - Right level of abstraction: “if you’ve only one diagram to spend, illustrate the concepts, the architecture, the difficult part”
 - The more detailed the documentation, the higher the probability for regression
“outdated/wrong documentation is worse than none”

General Diagramming Guidelines Ambler (2005)

(Note: “Exceptions prove the rule.”)



- **2.1 Readability**

- 1.–3. Support Readability of Lines

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References

References

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- OMG (2011a). Unified modeling language: Infrastructure, version 2.4.1. Technical Report formal/2011-08-05.
- OMG (2011b). Unified modeling language: Superstructure, version 2.4.1. Technical Report formal/2011-08-06.