

# Software Design, Modelling and Analysis in UML

## Lecture 2: Semantical Model

2016-10-20

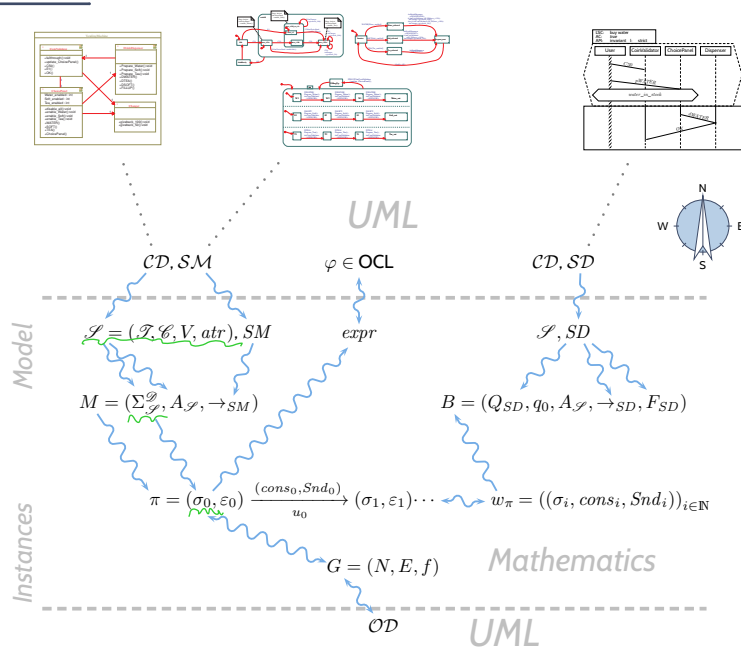
Prof. Dr. Andreas Podelski, Dr. ~~Bernd Westphal~~

Jochen Hoernicke

Albert-Ludwigs-Universität Freiburg, Germany

-2-2016-10-20-math-

### Course Map



-2-2016-10-20-Stephan-

## Content

- **Basic Object System Signature**
  - (basic) types, classes,
  - typed attributes,
  - attribute mapping.
- **Basic Object System Structure**
  - objects / object identities,
  - domains of basic and derived types.
- **System State**
  - concrete and symbolic,
  - dangling references,
- **A Complete Example**

## *Semantical Foundation*

## Basic Object System Signature

**Definition.** A (Basic) Object System **Signature** is a quadruple

$$\mathcal{S} = (\mathcal{T}, \mathcal{C}, V, \text{atr}) \quad C_1, C_2$$

where

- $\mathcal{T}$  is a set of (basic) **types**,
- $\mathcal{C}$  is a finite set of **classes**,
- $V$  is a finite set of **typed attributes**, i.e., each  $v \in V$  has a type
  - $\tau \in \mathcal{T}$ , or
  - $C_{0,1}$  or  $C_*$ , where  $C \in \mathcal{C}$
 (written  $v : \tau$  or  $v : C_{0,1}$  or  $v : C_*$ ),
- $\text{atr} : \mathcal{C} \rightarrow 2^V$  maps each class to its set of attributes.

Classes name  
total functions  
PowerSet of V

**Note:** Inspired by OCL 2.0 standard [OMG \(2006\)](#), Annex A.

-2-2016-10-10 - Semindan-

5/20

## Basic Object System Signature Example

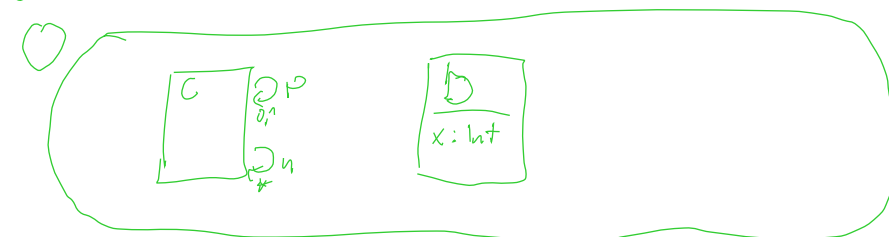
$\mathcal{S} = (\mathcal{T}, \mathcal{C}, V, \text{atr})$  where

- (basic) **types**  $\mathcal{T}$  and **classes**  $\mathcal{C}$  (both finite),
- **typed attributes**  $V$ ,  $\tau$  from  $\mathcal{T}$ , or  $C_{0,1}$  or  $C_*$ , for some  $C \in \mathcal{C}$ ,
- $\text{atr} : \mathcal{C} \rightarrow 2^V$  mapping classes to attributes.

**Example:**

$$\mathcal{S}_0 = (\{\text{Int}\}, \{C, D\}, \{x : \text{Int}, p : C_{0,1}, n : C_*\}, \{C \mapsto \{p, n\}, D \mapsto \{x\}\})$$

set of basic type  
Attributes V  
atr  
set of classe C  
 $\text{atr}(C) = \{p, n\}$   
 $\text{atr}(D) = \{x\}$



-2-2016-10-10 - Semindan-

6/20

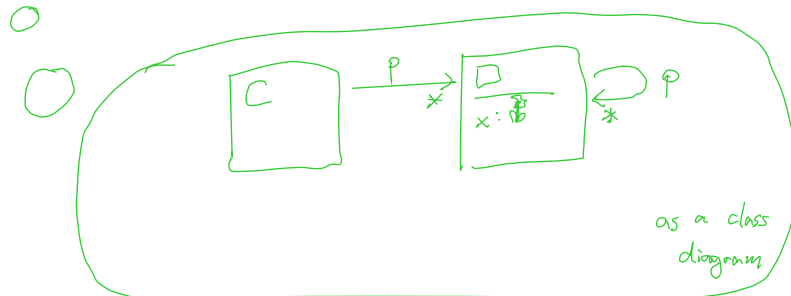
## Basic Object System Signature Another Example

$\mathcal{S} = (\mathcal{T}, \mathcal{C}, V, atr)$  where

- (basic) types  $\mathcal{T}$  and classes  $\mathcal{C}$  (both finite),
- typed attributes  $V, \tau$  from  $\mathcal{T}$ , or  $C_{0,1}$  or  $C_*$ , for some  $C \in \mathcal{C}$ ,
- $atr : \mathcal{C} \rightarrow 2^V$  mapping classes to attributes.

Example:

$\mathcal{S}_1 = (\{ \mathbb{B}, MyType \}, \{ C, D \}, \{ x: \mathbb{B}, p: D_*, q: D_{0,1} \}, \{ C \mapsto \{ p \}, D \mapsto \{ x, p \} \})$



-2-2016-10-10 - Semindam-

7/20

## Basic Object System Structure

**Definition.** A Basic Object System Structure of  $\mathcal{S} = (\mathcal{T}, \mathcal{C}, V, atr)$  is a domain function  $\mathcal{D}$  which assigns to each type a domain, i.e.

- $\tau \in \mathcal{T}$  is mapped to  $\mathcal{D}(\tau)$ ,  $\mathcal{D}(\mathcal{T})$  is a set
- $C \in \mathcal{C}$  is mapped to an infinite set  $\mathcal{D}(C)$  of (object) identities.  
Note: Object identities only have the "=" operation.
- Sets of object identities for different classes are disjoint, i.e.

$$\forall C, D \in \mathcal{C} : C \neq D \rightarrow \mathcal{D}(C) \cap \mathcal{D}(D) = \emptyset.$$

- $C_*$  and  $C_{0,1}$  for  $C \in \mathcal{C}$  are mapped to  $2^{\mathcal{D}(C)}$ .

We use  $\mathcal{D}(\mathcal{C})$  to denote  $\bigcup_{C \in \mathcal{C}} \mathcal{D}(C)$ ; analogously  $\mathcal{D}(\mathcal{C}_*)$ .

$$\mathcal{D}(C_*) := 2^{\mathcal{D}(C)}$$

$$\mathcal{D}(C_{0,1}) := 2^{\mathcal{D}(C)}$$

$$\mathcal{D}(\mathcal{C}_*) = \bigcup_{C \in \mathcal{C}} \mathcal{D}(C_*)$$

**Note:** We identify objects and object identities, because both uniquely determine each other (cf. OCL 2.0 standard).

-2-2016-10-10 - Semindam-

8/20

## Basic Object System Structure Example

**Wanted:** a structure for signature

$$\mathcal{S}_0 = (\{Int\}, \{C, D\}, \{x : Int, p : C_{0,1}, n : C_*\}, \{C \mapsto \{p, n\}, D \mapsto \{x\}\})$$

$\mathcal{D}$  needs to map:

- $\tau \in \mathcal{T}$  to **some**  $\mathcal{D}(\tau)$ .
- $C \in \mathcal{C}$  to **some** set of identities  $\mathcal{D}(C)$  (infinite, disjoint for different classes).
- $C_*$  and  $C_{0,1}$  for  $C \in \mathcal{C}$ : always mapped to  $\mathcal{D}(C_*) = \mathcal{D}(C_{0,1}) = 2^{\mathcal{D}(C)}$ .

$$\begin{aligned} \mathcal{D}(Int) &= \mathbb{Z} \\ \mathcal{D}(C) &= \mathbb{N}^+ \times \{C\} \cong \{1_C, 2_C, 3_C, 4_C, \dots\} \\ \mathcal{D}(D) &= \mathbb{N}^+ \times \{D\} \cong \{1_D, 2_D, 3_D, \dots\} \\ \mathcal{D}(C_{0,1}) = \mathcal{D}(C_*) &= 2^{\mathcal{D}(C)} \\ \mathcal{D}(D_{0,1}) = \mathcal{D}(D_*) &= 2^{\mathcal{D}(D)} \end{aligned} \quad \left| \begin{aligned} \mathcal{D}(Int) &= \{-2, -1, 0, 1, 2\} \\ \mathcal{D}(C) &= \{a, aa, aaa, \dots\} \\ \mathcal{D}(D) &= \{b, bb, bbb, \dots\} \end{aligned} \right.$$

-2-2016-10-10 - Seminar -

9/20

## System State

**Definition.** Let  $\mathcal{D}$  be a structure of  $\mathcal{S} = (\mathcal{T}, \mathcal{C}, V, atr)$ .

A **system state** of  $\mathcal{S}$  wrt.  $\mathcal{D}$  is a **type-consistent** mapping

$$\sigma : \mathcal{D}(\mathcal{C}) \rightarrow (V \rightarrow (\mathcal{D}(\mathcal{T}) \cup \mathcal{D}(\mathcal{C}_*)))$$

$\mathcal{D}(\mathcal{C})$ : Set of all class identities  
 $V$ : partial func.  
 $\mathcal{D}(\mathcal{T}) \cup \mathcal{D}(\mathcal{C}_*)$ : union of all domains  $\mathcal{D}(\tau)$  and all subset of object identities

That is, for each  $u \in \mathcal{D}(C)$ ,  $C \in \mathcal{C}$ , if  $u \in \text{dom}(\sigma)$

- $\text{dom}(\sigma(u)) = atr(C)$
  - $(\sigma(u))(v) \in \mathcal{D}(\tau)$  if  $v : \tau, \tau \in \mathcal{T}$
  - $\sigma(u)(v) \in \mathcal{D}(D_*)$  if  $v : D_{0,1}$  or  $v : D_*$  with  $D \in \mathcal{C}$
- } type-consistent

We call  $u \in \mathcal{D}(\mathcal{C})$  **alive** in  $\sigma$  if and only if  $u \in \text{dom}(\sigma)$ .

We use  $\Sigma_{\mathcal{D}}$  to denote the set of all system states of  $\mathcal{S}$  wrt.  $\mathcal{D}$ .

-2-2016-10-10 - Seminar -

10/20

## System State Example

$$\mathcal{S}_0 = (\{Int\}, \{C, D\}, \{x : Int, p : C_{0,1}, n : C_*\}, \{C \mapsto \{p, n\}, D \mapsto \{x\}\})$$

$$\mathcal{D}(Int) = \mathbb{Z}, \quad \mathcal{D}(C) = \{1_C, 2_C, 3_C, \dots\}, \quad \mathcal{D}(D) = \{1_D, 2_D, 3_D, \dots\}$$

**Wanted:**  $\sigma : \mathcal{D}(\mathcal{C}) \rightarrow (V \rightarrow (\mathcal{D}(\mathcal{T}) \cup \mathcal{D}(\mathcal{C}_*)))$  such that (i)  $\text{dom}(\sigma(u)) = \text{atr}(C)$ , and  
(ii)  $\sigma(u)(v) \in \mathcal{D}(\tau)$  if  $v : \tau, \tau \in \mathcal{T}$ , (iii)  $\sigma(u)(v) \in \mathcal{D}(C_*)$  if  $v : D_*$  with  $D \in \mathcal{C}$ .

$\sigma_0 = \emptyset$  ("empty function")  
alive in  $\sigma_0$ : none

$$\sigma_1 = \{ 1_C \mapsto \{ p \mapsto \emptyset, n \mapsto \{ 1_C, 5_C \} \}, 5_C \mapsto \{ p \mapsto \{ 1_C \}, n \mapsto \emptyset \}, 3_D \mapsto \{ x \mapsto 3 \} \}$$

$$\sigma_2 = \{ 1_D \mapsto \{ x \mapsto 27 \}, 2_D \mapsto \{ x \mapsto 27 \}, 1_{7_D} \mapsto \{ x \mapsto 0 \} \}$$

## System State Example

$$\mathcal{S}_0 = (\{Int\}, \{C, D\}, \{x : Int, p : C_{0,1}, n : C_*\}, \{C \mapsto \{p, n\}, D \mapsto \{x\}\})$$

$$\mathcal{D}(Int) = \mathbb{Z}, \quad \mathcal{D}(C) = \{1_C, 2_C, 3_C, \dots\}, \quad \mathcal{D}(D) = \{1_D, 2_D, 3_D, \dots\}$$

**Wanted:**  $\sigma : \mathcal{D}(\mathcal{C}) \rightarrow (V \rightarrow (\mathcal{D}(\mathcal{T}) \cup \mathcal{D}(\mathcal{C}_*)))$  such that (i)  $\text{dom}(\sigma(u)) = \text{atr}(C)$ , and  
(ii)  $\sigma(u)(v) \in \mathcal{D}(\tau)$  if  $v : \tau, \tau \in \mathcal{T}$ , (iii)  $\sigma(u)(v) \in \mathcal{D}(C_*)$  if  $v : D_*$  with  $D \in \mathcal{C}$ .

### Two options:

- **Concrete, explicit identities:**

$$\sigma = \{ 1_C \mapsto \{ p \mapsto \emptyset, n \mapsto \{ 5_C \} \}, 5_C \mapsto \{ p \mapsto \emptyset, n \mapsto \emptyset \}, 1_D \mapsto \{ x \mapsto 23 \} \}$$

$$1_C \mapsto \emptyset$$

- **Alternative: symbolic system state.**

$$\sigma = \{ c_1 \mapsto \{ p \mapsto \emptyset, n \mapsto \{ c_2 \} \}, c_2 \mapsto \{ p \mapsto \emptyset, n \mapsto \emptyset \}, d \mapsto \{ x \mapsto 23 \} \}$$

assuming  $c_1, c_2 \in \mathcal{D}(C), d \in \mathcal{D}(D), c_1 \neq c_2$ .

## System State: Spot the 10 (?) Mistakes

$$\mathcal{S}_0 = (\{Int\}, \{C, D\}, \{x : Int, p : C_{0,1}, n : C_*\}, \{C \mapsto \{p, n\}, D \mapsto \{x\}\})$$

$$\mathcal{D}(Int) = \mathbb{Z}, \quad \mathcal{D}(C) = \{1_C, 2_C, 3_C, \dots\}, \quad \mathcal{D}(D) = \{1_D, 2_D, 3_D, \dots\}$$

**Wanted:**  $\sigma : \mathcal{D}(\mathcal{C}) \rightarrow (V \rightarrow (\mathcal{D}(\mathcal{T}) \cup \mathcal{D}(\mathcal{C}_*)))$  such that (i)  $\text{dom}(\sigma(u)) = \text{atr}(C)$ , and  
 (ii)  $\sigma(u)(v) \in \mathcal{D}(\tau)$  if  $v : \tau, \tau \in \mathcal{T}$ , (iii)  $\sigma(u)(v) \in \mathcal{D}(C_*)$  if  $v : D_*$  with  $D \in \mathcal{C}$ .

- $\sigma = \{1_C \mapsto \{p \mapsto \emptyset, n \mapsto \{5_C\}\}, 5_C \mapsto \{p \mapsto \emptyset, n \mapsto 1_C\}, 1_D \mapsto \{x \mapsto 2, 3\}\}$ .
- $\sigma = \{1_C \mapsto \{p \mapsto \emptyset, n \mapsto \{5_C\}\}, 5_C \mapsto \{p \mapsto 1_C, n \mapsto \emptyset\}, 1_D \mapsto \{x \mapsto 23\}\}$ .  
*(iii)  $\nexists \notin \mathcal{D}(C_*)$   $\mathcal{D}(Int)$*
- $\sigma = \{1_C \mapsto \{p \mapsto \emptyset, n \mapsto \{1_D\}\}, 5_C \mapsto \{p \mapsto \emptyset, n \mapsto \emptyset\}, 1_D \mapsto \{x \mapsto 22\}\}$ .
- $\sigma = \{1_C \mapsto \{p \mapsto \emptyset, n \mapsto \{5_C\}\}, 5_C \mapsto \{n \mapsto \emptyset\}, 1_D \mapsto \{x \mapsto 1, p \mapsto \{1_C\}\}\}$ .  
*(iii)  $\nexists \notin \mathcal{D}(C_*)$*
- $\sigma = \{1_C \mapsto \{p \mapsto \emptyset, n \mapsto \{5_C\}\}, 5_C \mapsto \{p \mapsto \emptyset, n \mapsto \{9_C\}\}\}$   
*(iii)  $\{9_C\} \in \mathcal{D}(C_*)$*

-2-2016-10-10 - Seminar -

13/20

## Dangling References

**Definition.** Let  $\sigma \in \Sigma_{\mathcal{D}}$  be a system state.

We say attribute  $v \in V_{0,1,*}$ , i.e.  $v : C_{0,1}$  or  $v : C_*$ , in object  $u \in \text{dom}(\sigma)$  has a **dangling reference** if and only if the attribute's value comprises an object which is not alive in  $\sigma$ , i.e. if

$$\sigma(u)(v) \not\subseteq \text{dom}(\sigma).$$

We call  $\sigma$  **closed** if and only if no attribute has a dangling reference in any object alive in  $\sigma$ .

**Example:**

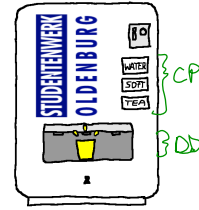
- $\sigma = \{1_C \mapsto \{p \mapsto \emptyset, n \mapsto \{5_C\}\}\}$

-2-2016-10-10 - Seminar -

14/20

## A Complete Example: Vending Machine Structure

$$\mathcal{J} = \left( \begin{array}{l} \{ \text{Int}, \text{Bool} \}, \\ \{ \text{VM}, \text{CP}, \text{DD} \}, \\ \{ \text{cp} : \text{CP}_x, \text{dd} : \text{DD}_{0,1}, \text{win} : \text{Int}, \text{wen} : \text{Bool} \}, \\ \left( \begin{array}{l} \text{VM} \mapsto \{ \text{cp}, \text{dd} \} \\ \text{CP} \mapsto \{ \text{wen} \} \\ \text{DD} \mapsto \{ \text{win} \} \end{array} \right) \end{array} \right)$$



$$\mathcal{D}(\text{Bool}) = \{ \text{true}, \text{false} \} \quad \mathcal{D}(\text{Int}) = \mathbb{Z}$$

$$\mathcal{D}(\text{VM}) = \{ 1_{\text{VM}}, 2_{\text{VM}}, \dots \}$$

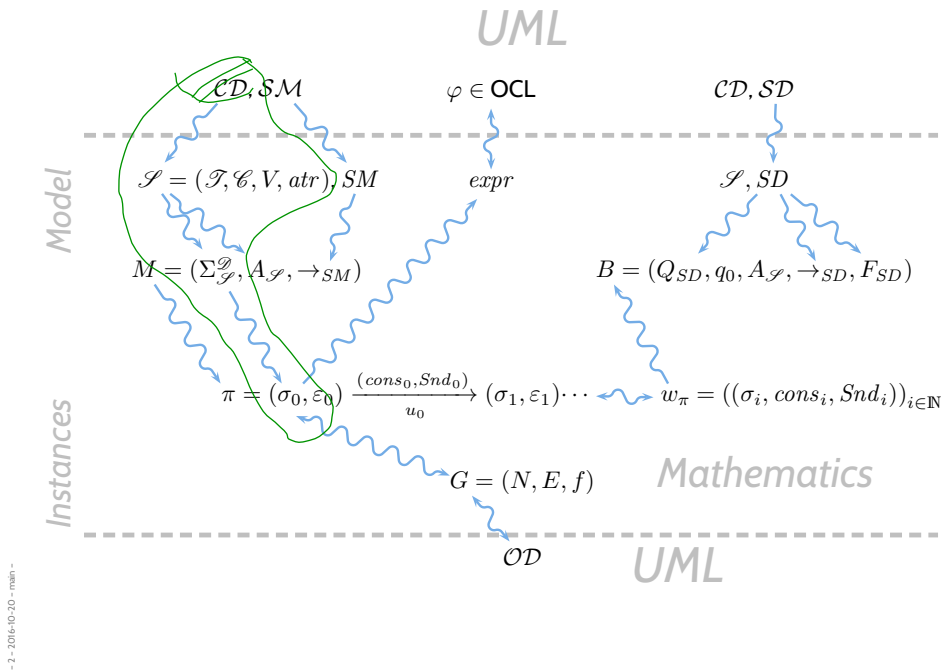
$$\mathcal{D}(\text{DD}) = \dots$$

$$\sigma = \left\{ \begin{array}{l} 1_{\text{VM}} \mapsto \left\{ \begin{array}{l} \text{dd} \mapsto \{ 1_{\text{DD}} \}, \text{cp} \mapsto \{ 4_{2_{\text{CP}}}, 3_{\text{CP}} \} \\ 1_{\text{DD}} = \{ \text{win} \mapsto 10 \}, 4_{2_{\text{CP}}} = \{ \text{wen} \mapsto \text{true} \}, 3_{\text{CP}} = \{ \text{wen} \mapsto \text{false} \} \end{array} \right\} \end{array} \right\}$$

You Are Here.



## Course Map



17/20

## Tell Them What You've Told Them...

- We can directly use **object system signatures** to model the structure of systems.
  - We don't **need** diagrams, they will be more pleasant to read.
- We introduce
  - **basic types** and **classes**,
  - basic type and derived type **attributes**, and
  - assign to each class a set of attributes.
- **Object system structures** provide domains for basic and derived types.
- An **object system signature**  $\mathcal{S}$  and an **object system structure**  $\mathcal{D}$  uniquely define the set  $\Sigma_{\mathcal{S}}$  of **system states**.
- **Outlook:**
  - **Object system signatures** will be used to capture the **abstract syntax** of class diagrams.
  - **OCL expressions** will be evaluated on **system states**.
  - **State machines** will define sequences of **system configurations** (consisting of a **system state** and an **ether**).

-2-2016-10-30-SMv11-

18/20

## *References*

## *References*

OMG (2006). Object Constraint Language, version 2.0. Technical Report formal/06-05-01.

OMG (2011a). Unified modeling language: Infrastructure, version 2.4.1. Technical Report formal/2011-08-05.

OMG (2011b). Unified modeling language: Superstructure, version 2.4.1. Technical Report formal/2011-08-06.