

Software Design, Modelling and Analysis in UML

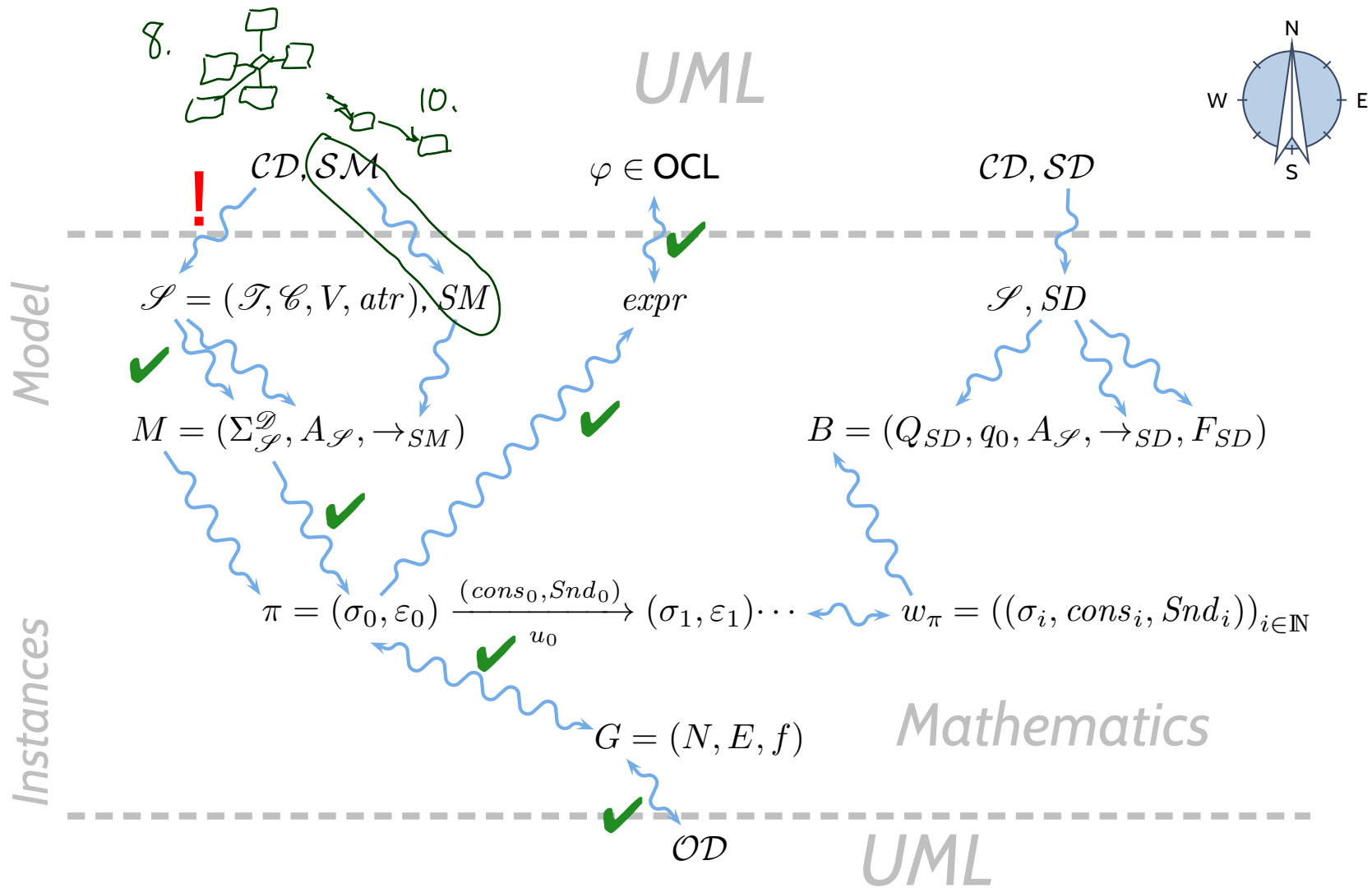
Lecture 8: Class Diagrams III

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Course Map



- **Recall: Associations**
 - Overview & Plan
 - (Temporarily) **Extend Signature**
- From **Class Diagrams** to **Signatures**
 - What if Things are Missing?
- **Association Semantics**
 - **Links** in System States
 - Associations and **OCL**
- **The Rest**
 - **Visibility, Navigability**
 - **Multiplicity, Properties,**
 - **Ownership, “Diamonds”**

```
graph LR; C[C] --> D[D]
```
- **Back to the Main Track**

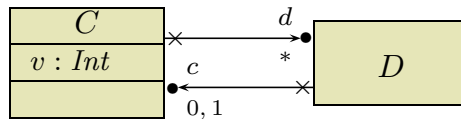
Recall: Plan & Extended Signature

Overview

- **Class diagram:**



- **Alternative presentation:**



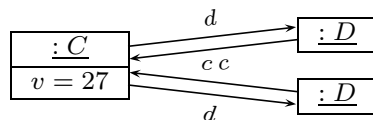
- **Signature:**

$$\mathcal{S} = (\{Int\}, \{C, D\}, \{v : Int, d : D_*, c : C_{0,1}\}, \{C \mapsto \{v, d\}, D \mapsto \{c\}\})$$

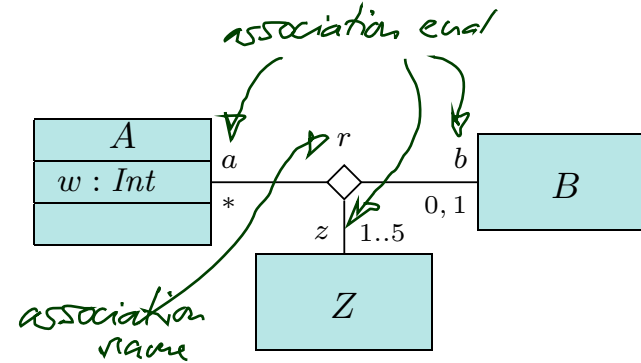
- **Example system state:**

$$\sigma = \{1_C \mapsto \{v \mapsto 27, d \mapsto \{5_D, 7_D\}\}, 5_D \mapsto \{c \mapsto \{1_C\}\}, 7_D \mapsto \{c \mapsto \{1_C\}\}\}$$

- **Object diagram:**



- **Class diagram (with ternary association):**



- **Signature:** extend again to represent

- **association** r with

- **association ends** $a, b,$ and z (each with multiplicity, visibility, etc.)

- **Example system state:** (σ, λ)

$$\sigma = \{1_A \mapsto \{w \mapsto 13\}, 1_B \mapsto \emptyset, 1_Z \mapsto \emptyset\}$$

$$\lambda = \{r \mapsto \{(1_A, 1_B, 1_Z), (1_A, 1_B, 2_Z)\}\}$$



- **Object diagram:** No...

So, What Do We (Have to) Cover?

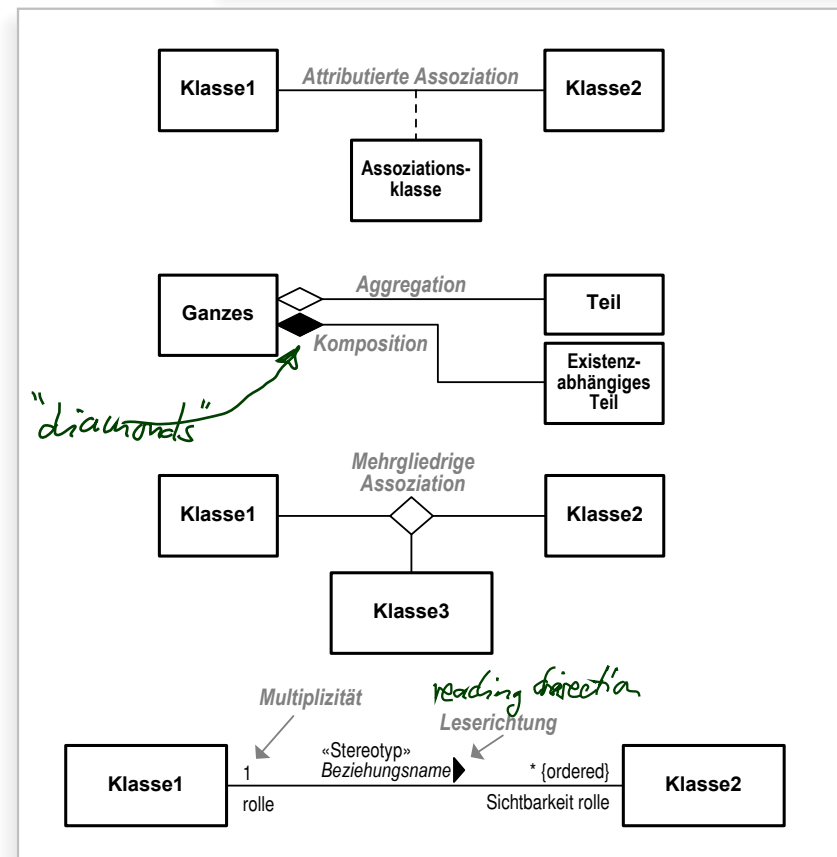
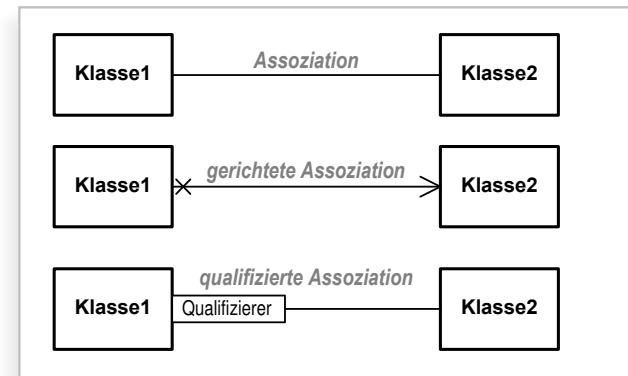
An **association** has

- a **name**,
- a **reading direction**, and
- at least two **ends**.

Each **end** has

- a **role name**,
- a **multiplicity**,
- a set of **properties**, such as **unique**, **ordered**, etc.
- a **qualifier**, (not in lect.)
- a **visibility**,
- a **navigability**,
- an **ownership**,
- and possibly a **diamond**.

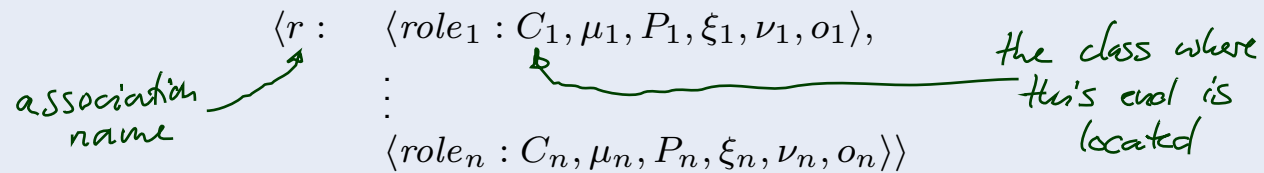
Wanted: places in the signature to represent the information from the picture.



Temporarily (Lecture 7 – 9) Extended Signature

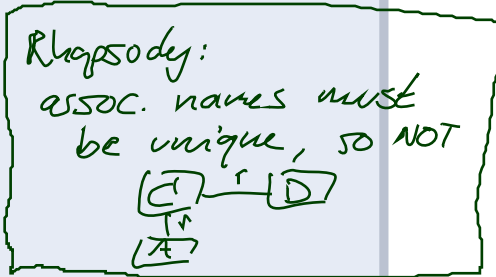
Definition. An (Extended) Object System **Signature** (with Associations) is a quadruple $\mathcal{S} = (\mathcal{T}, \mathcal{C}, V, atr)$ where

- ...
- each element of V is
 - either a **basic type attribute** $\langle v : T, \xi, expr_0, P_v \rangle$ with $T \in \mathcal{T}$
 - or an **association** of the form



(ends with multiplicity μ_i , properties P_i , visibility ξ_i , navigability ν_i , ownership o_i , $1 \leq i \leq n$)

- ...
- $atr : \mathcal{C} \rightarrow 2^{\{v \in V \mid v:T, T \in \mathcal{T}\}}$ maps classes to **basic type (!)** attributes.



In other words:

- only **basic type attributes** “belong” to a class (may appear in $atr(C)$),
- **associations** are not “owned” by a class (not in any $atr(C)$), but “live on their own”.

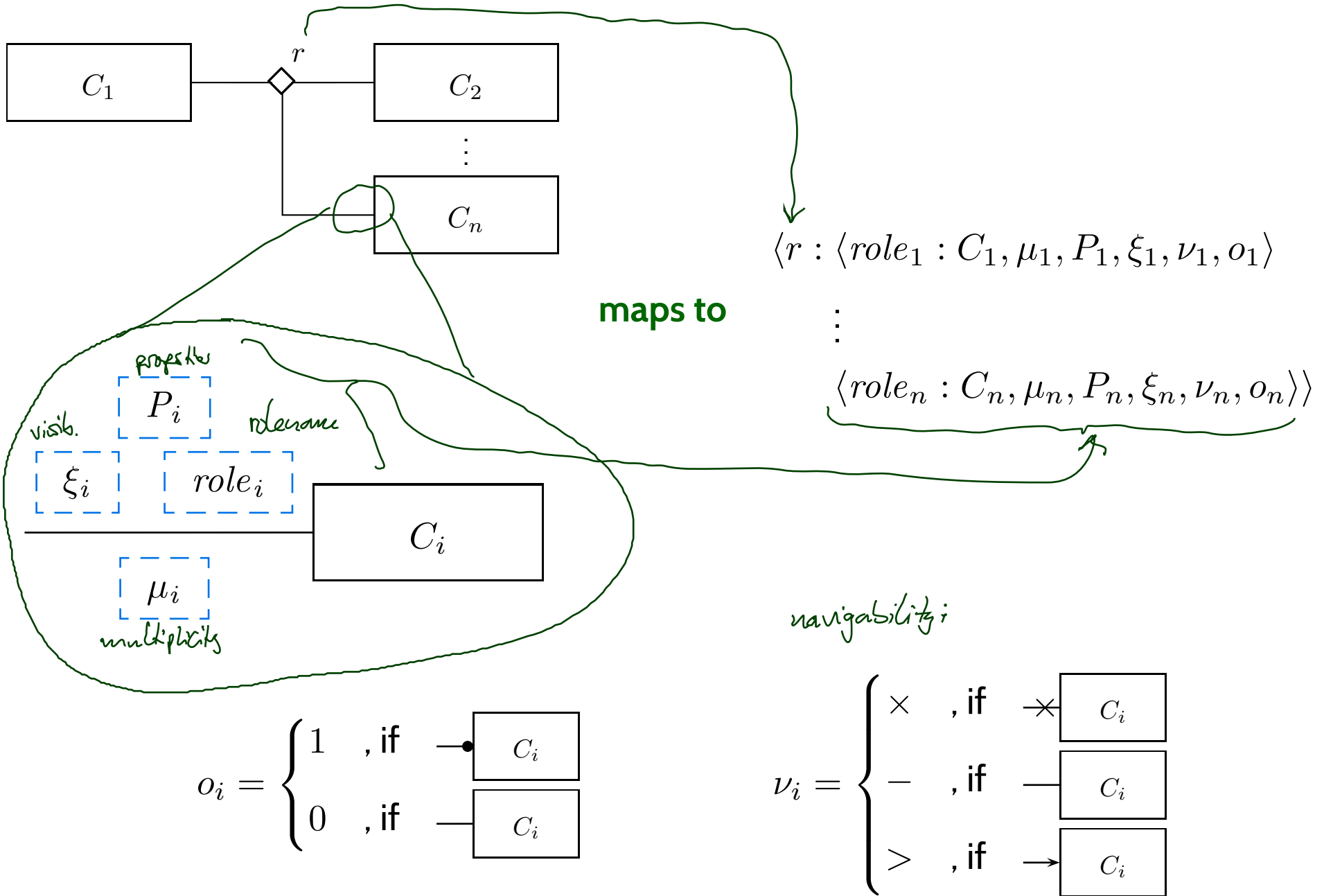
$$M ::= N..M \mid N..* \mid \mu, \mu$$

$$\langle * := 0..* \rangle$$

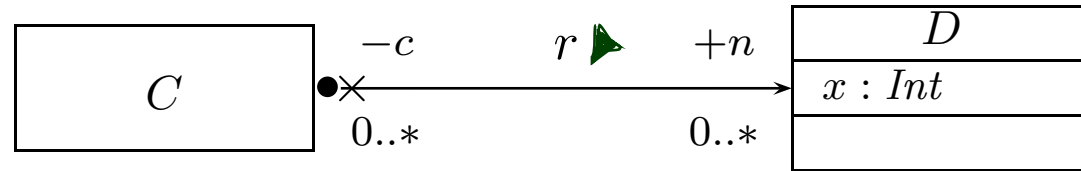
$$N ::= N..N$$

Associations in Class Diagrams

From Association Lines to Extended Signatures



Association Example



Signature:

$$\mathcal{S} = (\{Int\}, \{C, D\}, \{ \langle x: Int, +, \mathbb{K}, \emptyset \rangle, \\
 \langle r: \langle n: D, *, \{unique\}, +, \rangle, 0 \rangle, \\
 \langle c: C, 0..*, \{unique\}, -, x, ? \rangle \} , \\
 \{ C \mapsto \emptyset, D \mapsto \{x\} \})$$

What If Things Are Missing?

Most components of associations or association end may be omitted.
For instance (OMG, 2011b, 17), Section 6.4.2, proposes the following rules:

- **Name:** Use

$$A_ \langle C_1 \rangle _ \cdots _ \langle C_n \rangle$$

if the name is missing.

Example:



- **Reading Direction:** no default.
- **Role Name:** use the class name at that end in lower-case letters

Example:



Other convention: (used e.g. by modelling tool Rhapsody)



What If Things Are Missing?

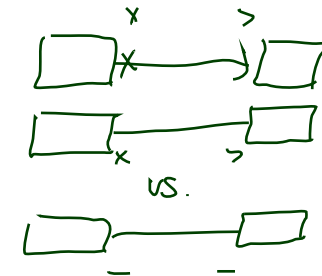
- **Multiplicity:** 1

In my opinion, it's safer to assume 0..1 or * (for 0..*) if there are no fixed, written, agreed conventions (“expect the worst”).

- **Properties:** \emptyset (in course: {unique})

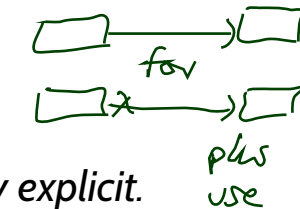
- **Visibility:** public

- **Navigability and Ownership:** not so easy. (OMG, 2011b, 43)



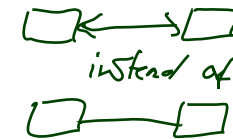
“Various options may be chosen for showing navigation arrows on a diagram.

In practice, it is often convenient to suppress some of the arrows and crosses and just show exceptional situations:



- Show all arrows and \times 's: Navigation and its absence are made completely explicit.

- Suppress all arrows and \times 's: No inference can be drawn about navigation.



This is similar to any situation in which information is suppressed from a view.

- Suppress arrows for associations with navigability in both directions, and show arrows only for associations with one-way navigability.

In this case, the two-way navigability cannot be distinguished from situations where there is no navigation at all; however, the latter case occurs rarely in practice.”

Wait, If Omitting Things...

- ...**is causing so much trouble** (e.g. leading to misunderstanding), why does the standard say “**In practice, it is often convenient...**”?

Is it a good idea to trade **convenience** for **precision/unambiguity**?

It depends.

- Convenience as such is a **legitimate goal**.
- In UML-As-Sketch mode, precision “**doesn't matter**”, so convenience (for writer) can even be a primary goal.
- In UML-As-Blueprint mode, **precision** is the **primary goal**. And misunderstandings are in most cases annoying.

But: (even in UML-As-Blueprint mode)

If all associations in your model have multiplicity *, then it's probably a good idea not to write all these *'s.

So: tell the reader about your convention and leave out the *'s.

Associations: Semantics

Associations in General

Recall: We consider associations of the following form:

$$\langle r : \langle role_1 : C_1, \mu_1, P_1, \xi_1, \nu_1, o_1 \rangle, \dots, \langle role_n : C_n, \mu_n, P_n, \xi_n, \nu_n, o_n \rangle \rangle$$

Only these parts are relevant for extended system states:

$$\langle r : \langle role_1 : C_1, _, P_1, _, _, _ \rangle, \dots, \langle role_n : C_n, _, P_n, _, _, _ \rangle \rangle$$

(recall: we assume $P_1 = P_n = \{\text{unique}\}$).

The UML standard “thinks” of associations as **n-ary relations** which “**live on their own**” in a system state.

That is, **links** (= association instances)

- **do not** belong (in general) to certain objects (in contrast to pointers, e.g.)
- are “first-class citizens” **next to objects**,
- are (in general) **not** directed (in contrast to pointers).

Links in System States

$$\langle r : \langle role_1 : C_1, _, P_1, _, _, _ \rangle, \dots, \langle role_n : C_n, _, P_n, _, _, _ \rangle \rangle$$

Only for the course of lectures ~~7~~8 / 9 we change the definition of system states:

Definition. Let \mathcal{D} be a structure of the (extended) signature with associations $\mathcal{S} = (\mathcal{I}, \mathcal{C}, V, atr)$.

A **system state** of \mathcal{S} wrt. \mathcal{D} is a pair (σ, λ) consisting of

- a type-consistent mapping (as before)

$$\sigma : \mathcal{D}(\mathcal{C}) \rightarrow (atr(\mathcal{C}) \rightarrow \mathcal{D}(\mathcal{I})),$$

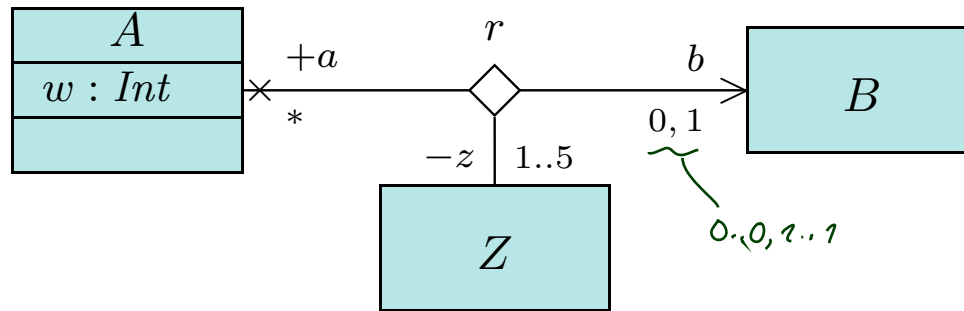
only basic type
attributes here

- a mapping λ which maps each association $\langle r : \langle role_1 : C_1 \rangle, \dots, \langle role_n : C_n \rangle \rangle \in V$ to a **relation**

$$\lambda(r) \subseteq \mathcal{D}(C_1) \times \dots \times \mathcal{D}(C_n)$$

(i.e. a set of type-consistent n -tuples of identities).

Association / Link Example



Signature:

$$\mathcal{S} = (\{Int\}, \{A, Z, B\}, \{w: Int, \langle r: \langle a: A, 0..*, +, \{unique\}, *, 0 \rangle, \langle b: \langle z: Z, 1..5, -, \{unique\}, -, 0 \rangle, \langle z: \langle b: B, 0, 1, +, \{unique\}, >, 0 \rangle \rangle, \{A \mapsto \{w\}, Z \mapsto \emptyset, B \mapsto \emptyset\})$$

System state:

$$\sigma = \{ \begin{array}{l} 1_A \mapsto \{w \mapsto 27\}, \\ 2_A \mapsto \{w \mapsto 13\}, \\ 4_Z \mapsto \emptyset, \\ 3_B \mapsto \emptyset, \\ 7_B \mapsto \emptyset, \\ 8_B \mapsto \emptyset, \\ 3_A \mapsto \emptyset \end{array} \}$$

$$\in \mathcal{D}(A) \times \mathcal{D}(Z) \times \mathcal{D}(B)$$

$$\lambda = \{ r \mapsto \{ \begin{array}{l} (1_A, 4_Z, 3_B), \\ (1_A, 4_Z, 7_B), \\ (1_A, 4_Z, 5_B), \\ (2_A, 4_Z, 3_B) \end{array} \} \}$$

NOT: $(4_Z, 3_B, 2_A)$
NOT: $(1_A, 4_Z)$

a	z	b
1 _A	4 _Z	3 _B
1 _A	4 _Z	7 _B
1 _A	4 _Z	5 _B
2 _A	4 _Z	3 _B

Associations and OCL

OCL and Associations: Syntax

Recall: OCL syntax as introduced in Lecture 3, interesting part:

$$\begin{array}{l|ll} \text{expr} ::= \dots & r_1(\text{expr}_1) & : \tau_C \rightarrow \tau_D & r_1 : D_{0,1} \in \text{atr}(C) \\ & r_2(\text{expr}_1) & : \tau_C \rightarrow \text{Set}(\tau_D) & r_2 : D_* \in \text{atr}(C) \end{array}$$

Now becomes

$$\begin{array}{l|ll} \text{expr} ::= \dots & \text{role}(\text{expr}_1) & : \tau_C \rightarrow \tau_D & \mu = 0..1 \text{ or } \mu = 1..1 \\ & \text{role}(\text{expr}_1) & : \tau_C \rightarrow \text{Set}(\tau_D) & \text{otherwise} \end{array}$$

if there is

$$\begin{array}{l} \langle r : \dots, \langle \text{role} : D, \mu, _ , _ , _ , _ \rangle, \dots, \langle \text{role}' : C, _ , _ , _ , _ \rangle, \dots \rangle \in V \text{ or} \\ \langle r : \dots, \langle \text{role}' : C, _ , _ , _ , _ \rangle, \dots, \langle \text{role} : D, \mu, _ , _ , _ , _ \rangle, \dots \rangle \in V, \quad \text{role} \neq \text{role}' \end{array}$$

Note:

- Association name as such **does not occur** in OCL syntax, role names do.
- expr_1 has to denote an object of a class which “participates” in the association.

OCL and Associations: Semantics

Recall:

Assume $expr_1 : \tau_C$ for some $C \in \mathcal{C}$. Set $u_1 := I[expr_1](\sigma, \beta) \in \mathcal{D}(T_C)$.

- $I[r_1(expr_1)](\sigma, \beta) := \begin{cases} u & , \text{if } u_1 \in \text{dom}(\sigma) \text{ and } \sigma(u_1)(r_1) = \{u\} \\ \perp & , \text{otherwise} \end{cases}$
- $I[r_2(expr_1)](\sigma, \beta) := \begin{cases} \sigma(u_1)(r_2) & , \text{if } u_1 \in \text{dom}(\sigma) \\ \perp & , \text{otherwise} \end{cases}$

Now needed:

$$I[role(expr_1)]((\sigma, \lambda), \beta)$$

- We cannot simply write $\sigma(u)(role)$.

Recall: *role* is (for the moment) not an attribute of object u (not in $atr(C)$).

- What we have is $\lambda(r)$ (with association name r , not with role name *role*!).

$$\langle r : \dots, \langle role : D, \mu, _ , _ , _ , _ \rangle, \dots, \langle role' : C, _ , _ , _ , _ , _ \rangle, \dots \rangle$$

But it yields a set of n -tuples, of which **some** relate u and some instances of D .

- *role* denotes the position of the D 's in the tuples constituting the value of r .

OCL and Associations: Semantics Cont'd

Assume $expr_1 : \tau_C$ for some $C \in \mathcal{C}$. Set $u_1 := I[expr_1](\langle \sigma, \lambda \rangle, \beta) \in \mathcal{D}(T_C)$.

- $I[role(expr_1)](\langle \sigma, \lambda \rangle, \beta) := \begin{cases} u & , \text{if } u_1 \in \text{dom}(\sigma) \text{ and } L(role)(u_1, \lambda) = \{u\} \\ \perp & , \text{otherwise} \end{cases}$
- $I[role(expr_1)](\langle \sigma, \lambda \rangle, \beta) := \begin{cases} L(role)(u_1, \lambda) & , \text{if } u_1 \in \text{dom}(\sigma) \\ \perp & , \text{otherwise} \end{cases}$

where

$$L(\underline{role})(u, \lambda) = \left\{ \{ (u_1, \dots, u_n) \in \lambda(r) \mid u \in \{u_1, \dots, u_n\} \} \right\} \downarrow i$$

if

$$\langle r : \langle role_1 : _ , _ , _ , _ , _ , _ \rangle , \dots \langle role_n : _ , _ , _ , _ , _ , _ \rangle \rangle , \quad \underline{role} = \underline{role}_i$$

project onto
i-th
component

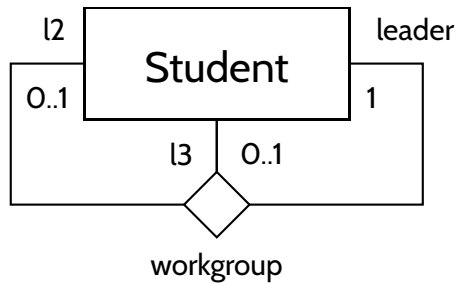
Given a set of n -tuples A ,

$A \downarrow i$ denotes the element-wise projection onto the i -th component.

OCL and Associations Semantics: Example

$$I[\text{role}(\text{expr}_1)]((\sigma, \lambda), \beta) := \begin{cases} u & , \text{if } u_1 \in \text{dom}(\sigma) \text{ and } L(\text{role})(u_1, \lambda) = \{u\} \\ \perp & , \text{otherwise} \end{cases}$$

$$I[\text{role}(\text{expr}_1)]((\sigma, \lambda), \beta) := \begin{cases} L(\text{role})(u_1, \lambda) & , \text{if } u_1 \in \text{dom}(\sigma) \\ \perp & , \text{otherwise} \end{cases} \quad \begin{array}{l} L(\text{role})(u, \lambda) = \{(u_1, \dots, u_n) \\ \in \lambda(r) \mid u \in \{u_1, \dots, u_n\}\} \downarrow i \end{array}$$



$\neq := \text{allInstances}_{\text{Student}} \rightarrow$
 $\text{Exists}(s \mid s.l2 = s.l3)$

1. 2. 3.
leader l₂ l₃

$$\lambda(\text{workgroup}) = \{(1_s, 2_s, 3_s), (1_s, 3_s, 4_s), (5_s, 1_s, 1_s)\}$$

$$I[\neq](\sigma, \lambda, \beta) =: \beta_1$$

$$I[s.l2](\sigma, \lambda, \{s \mapsto 5_s\}) = 1_s$$

$$u_1 = I[s](\sigma, \lambda, \beta_1) = \beta_1(s) = 5_s$$

$$L(l2)(u_1, \lambda) = \{(5_s, 1_s, 1_s)\} \downarrow 2 = \{1_s\}$$

$$I[s.l2](\sigma, \lambda, \{s \mapsto 1_s\}) = \perp$$

$u_1 = 1_s$

$$L(l2)(u_1, \lambda) = \{(5_s, 1_s, 1_s), (1_s, 2_s, 3_s), (1_s, 3_s, 4_s)\} \downarrow 2 = \{2_s, 3_s\}$$

1_s

Associations: The Rest

Recapitulation: Consider the following association:

$$\langle r : \langle role_1 : C_1, \mu_1, P_1, \xi_1, \nu_1, o_1 \rangle, \dots, \langle role_n : C_n, \mu_n, P_n, \xi_n, \nu_n, o_n \rangle \rangle$$

- **Association name** r and **role names / types** $role_i / C_i$ induce extended system states (σ, λ) .
- **Multiplicity** μ is considered in OCL syntax.
- **Visibility** ξ / **Navigability** ν : well-typedness (in a minute).

Now the rest:

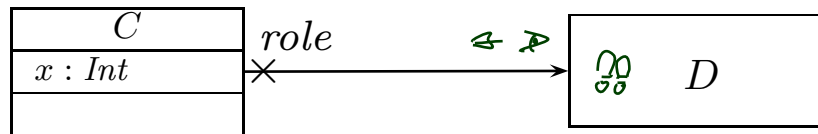
- **Multiplicity** μ : we propose to view them as constraints.
- **Properties** P_i : even more typing.
- **Ownership** o : getting closer to pointers/references.
- **Diamonds**: exercise.

Navigability

Navigability is treated similar to visibility:

Using names of non-navigable association ends ($\nu = \times$) are **forbidden**.

Example: Given



the following OCL expression is **not well-typed** wrt. navigability,

context D inv : role.x > 0

The standard says: navigation is...

- '—': ...possible
- '×': ...not possible
- '>': ...efficient



So: In general, UML associations **are different** from pointers / references in general!

But: Pointers / references **can faithfully** be modelled by UML associations.

Multiplicities as Constraints

Recall: Multiplicity is a term of the form $N_1..N_2, \dots, N_{2k-1}..N_{2k}$

where $N_i \leq N_{i+1}$ for $1 \leq i \leq 2k$, $N_1, \dots, N_{2k-1} \in \mathbb{N}$, $N_{2k} \in \mathbb{N} \cup \{*\}$.

Define $\mu_{\text{OCL}}^C(\text{role}) :=$

context C inv : $(N_1 \leq \text{role} \rightarrow \text{size}() \leq N_2)$ or ... or $(N_{2k-1} \leq \text{role} \rightarrow \text{size}() \leq \underbrace{N_{2k}}_{\text{omit if } N_{2k} = *})$

for each $\langle r : \dots, \langle \text{role} : D, \mu, _ , _ , _ , _ \rangle, \dots, \langle \text{role}' : C, _ , _ , _ , _ \rangle, \dots \rangle \in V$ or

$\langle r : \dots, \langle \text{role}' : C, _ , _ , _ , _ \rangle, \dots, \langle \text{role} : D, \mu, _ , _ , _ , _ \rangle, \dots \rangle \in V,$

with $\text{role} \neq \text{role}'$, if $\mu \neq 0..1$, $\mu \neq 1..1$, and

$\mu_{\text{OCL}}^C(\text{role}) := \text{context } C \text{ inv} : \text{not}(\text{ocllsUndefined}(\text{role}))$

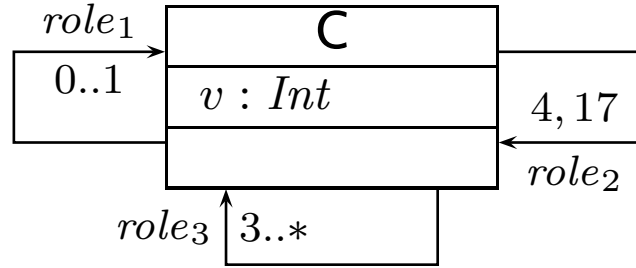
if $\mu = 1..1$.

Note: in n -ary associations with $n > 2$, there is redundancy.

Multiplicities as Constraints Example

$$\mu_{\text{OCL}}^C(\text{role}) = \text{context } C \text{ inv :} \\ (N_1 \leq \text{role} \rightarrow \text{size}() \leq N_2) \text{ or } \dots \text{ or } (N_{2k-1} \leq \text{role} \rightarrow \text{size}() \leq N_{2k})$$

\mathcal{CD} :



- $\{\text{context } C \text{ inv : } 4 \leq \text{role}_2 \rightarrow \text{size}() \leq 4 \text{ or } 17 \leq \text{role}_2 \rightarrow \text{size}() \leq 17\}$
 $= \{\text{context } C \text{ inv : } \text{role}_2 \rightarrow \text{size}() = 4 \text{ or } \text{role}_2 \rightarrow \text{size}() = 17\}$
- $\cup \{\text{context } C \text{ inv : } 3 \leq \text{role}_3 \rightarrow \text{size}()\}$

Properties

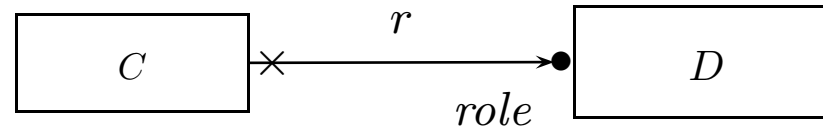
We don't want to cover association **properties** in detail, only some observations (assume binary associations):

Property	Intuition	Semantical Effect
unique	one object has at most one r -link to a single other object	current setting
bag	one object may have multiple r -links to a single other object	have $\lambda(r)$ yield multi-sets
ordered, sequence	an r -link is a sequence of object identities (possibly including duplicates)	have $\lambda(r)$ yield sequences

Property	OCL Typing of expression $role(expr)$
unique	$T_D \rightarrow Set(T_C)$
bag	$T_D \rightarrow Bag(T_C)$
ordered, sequence	$T_D \rightarrow Seq(T_C)$

For **subsets**, **redefines**, **union**, etc. see (? , 127).

Ownership



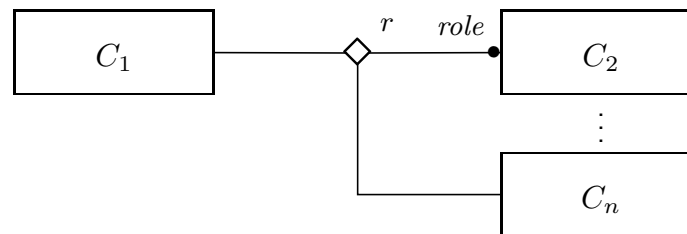
Intuitively it says:

Association r is **not a “thing on its own”** (i.e. provided by λ), but association end ‘ $role$ ’ is **owned** by C (!). (That is, it’s stored inside C object and provided by σ).

So: if multiplicity of $role$ is $0..1$ or $1..1$, then the picture above is very close to concepts of pointers/references.

Actually, ownership is seldom seen in UML diagrams. Again: if target platform is clear, one may well live without (cf. [\(OMG, 2011b, 42\)](#) for more details).

Not clear to me:



Back to the Main Track

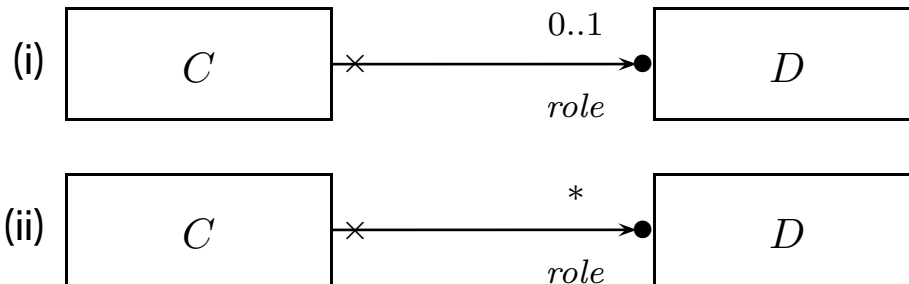
Back to the main track:

Recall: on some earlier slides we said, the extension of the signature is **only** to study associations in “full beauty”.

For the remainder of the course, we should look for something simpler...

Proposal:

- **from now on**, we only use associations of the form



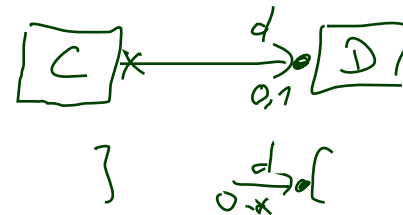
(And we may omit the non-navigability and ownership symbols.)

- Form (i) introduces $role : C_{0,1}$, and form (ii) introduces $role : C_*$ in V .
- In both cases, $role \in atr(C)$.
- We drop λ and go back to our nice σ with $\sigma(u)(role) \subseteq \mathcal{D}(D)$.

Tell Them What You've Told Them...

- From class diagrams with (general) **associations**, we obtain **extended signatures**. ✓
- Links (instances of associations) “live on their own” in the λ in extended system states (σ, λ) . ✓
- OCL considers **role names**, the **semantics** is (more or less) **straightforward**. ✓
- **The Rest:**
 - **navigability** is treated like visibility, ✓
 - view **multiplicities** as shorthand for **constraints**, ⚠
 - properties, ownership, “diamonds”: exist ✓
- **Back to the main track:**

For simplicity, let's restrict the following discussion to $C_{0,1}$ and C_* as before (now viewed as abbreviations for particular associations).



References

References

OMG (2011a). Unified modeling language: Infrastructure, version 2.4.1. Technical Report formal/2011-08-05.

OMG (2011b). Unified modeling language: Superstructure, version 2.4.1. Technical Report formal/2011-08-06.