

Software Design, Modelling and Analysis in UML

Lecture 10: Constructive Behaviour, State Machines Overview

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Contents & Goals

Last Lecture:

- Completed discussion of modelling **structure**.

This Lecture:

- Educational Objectives:** Capabilities for following tasks/questions.
 - Discuss the style of this class diagram.
 - What's the difference between reflective and constructive descriptions of behaviour?
 - What's the purpose of a behavioural model?
 - What does this State Machine mean? What happens if I inject this event?
 - Can you please model the following behaviour.
- Content:**
 - Purposes of Behavioural Models
 - Constructive vs. Reflective
 - UML Core State Machines (first half)

Modelling Behaviour

Stocktaking...

Have: Means to model the **structure** of the system.

- Class diagrams graphically, concisely describe sets of system states.
- OCL expressions logically state constraints/invariants on system states.

Want: Means to model **behaviour** of the system.

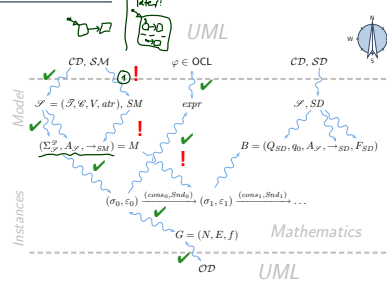
- Means to describe how system states evolve over **time**, that is, to describe sets of **sequences**

$$\sigma_0, \sigma_1, \dots \in \Sigma^\omega$$

of system states.

not refl.-free, discrete time

Course Map



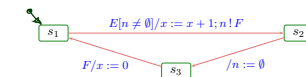
Constructive UML

UML provides two visual formalisms for constructive description of behaviours:

- Activity Diagrams**
- State-Machine Diagrams**

We (exemplary) focus on State-Machines because

- somehow "practice proven" (in different flavours),
 - prevalent in embedded systems community,
 - indicated useful by [Dobing and Parsons, 2006] survey, and
 - Activity Diagram's intuition changed (between UML 1.x and 2.x) from transition-system-like to petri-net-like...
- Example state machine:



UML State Machines

"we can take this transition if there is an E ready for us and..."

"state can F to object if checked by G in"

action

low corresponding transitions

Brief History:

- Rooted in Moore/Mealy machines, Transition Systems
- [Harel, 1987]: Statecharts as a concise notation, introduces in particular hierarchical states.
- Manifest in tool **Statemate** [Harel et al., 1990] (simulation, code-generation); nowadays also in **Matlab/Simulink**, etc.
- From UML 1.x on: **State Machines** (not the official name, but understood: UML-Statecharts)
- Late 1990's: tool **Rhapsody** with code-generation for state machines.

Note: there is a common core, but each dialect interprets some constructs subtly different [Crane and Dingel, 2007]. (*Would be too easy otherwise...*)

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Roadmap: Chronologically

- (i) What do we (have to) cover?
UML State Machine Diagrams **Syntax**.
- (ii) Def.: Signature with **signals**.
- (iii) Def.: **Core state machine**.
- (iv) Map UML State Machine Diagrams to core state machines.

Semantics:
The Basic Causality Model

- (v) Def.: **Ether** (aka. event pool)
- (vi) Def.: **System configuration**.
- (vii) Def.: **Event**.
- (viii) Def.: **Transformer**.
- (ix) Def.: **Transition system**, computation.
- (x) Transition relation induced by core state machine.
- (xi) Def.: **step, run-to-completion step**.
- (xii) Later: Hierarchical state machines.

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UML State-Machines: What do we have to cover?

Handwritten annotations in German: *initial state*, *state*, *transitions*, *actions*, *guards*, *initial state (class. def.)*, *initial state (class. def.)*, *initial state (class. def.)*, *initial state (class. def.)*.

UML State-Machines: What do we have to cover?

Handwritten annotations in German: **Proven approach:**
Start out simple, consider the essence, namely
• basic/leaf states
• transitions,
then extend to cover the complicated rest.

Signature With Signals

Definition. A tuple $\mathcal{S} = (\mathcal{S}, \mathcal{E}, V, \text{atr}, \delta)$, δ a set of signals, is called signature (with signals) if and only if $(\mathcal{S}, \mathcal{E} \cup \{\delta, \tau, \text{atr}$ is a signature (as before).

Handwritten note: WE USE THIS

Note: Thus conceptually, a signal is a class and can have attributes of plain type and associations.
Alternative (maybe even better) definition:
 $\mathcal{E}(\mathcal{S}) = \{ \langle C, \mathcal{S}, a, \tau \rangle \in \mathcal{E} \mid \text{signal} \in \mathcal{S} \}$

Handwritten examples:
• $\langle \text{Signal}, \mathcal{S}, \tau, \text{atr} \rangle$
• $\langle \text{Signal}, \mathcal{S}, \tau, \text{atr} \rangle$

Core State Machine

Definition. A core state machine over signature $\mathcal{S} = (\mathcal{S}, \mathcal{E}, V, \text{atr}, \delta)$ is a tuple $M = (S, s_0, \rightarrow)$ where S is a non-empty, finite set of (basic) states, $s_0 \in S$ is an initial state, and $\rightarrow \subseteq S \times (\mathcal{E} \cup \{\tau\}) \times \text{Act}_{\mathcal{S}} \times S$ is a labelled transition relation.

Handwritten notes: disjoint union: - should not already be in \mathcal{E} (otherwise renamed first), source state, signals in \mathcal{S} , dest. state, trigger, guard, action.

We assume a set $\text{Expr}_{\mathcal{S}}$ of boolean expressions over \mathcal{S} (for instance OCL, may be something else) and a set $\text{Act}_{\mathcal{S}}$ of actions.

From UML to Core State Machines: By Example

UML state machine diagram SM:

Handwritten notes: *initial state*, *trigger*, *guard*, *action*, *annot := [[(event) ['.' (event)]* ['!' (guard)]*]* ['*' (action)]]*, *(default: true, assumed to be in Expr_S)*, *(default: stop, assumed to be in Act_S)*.

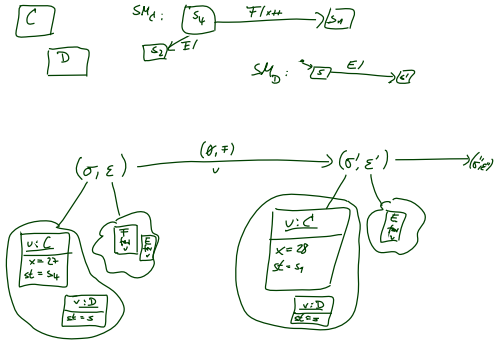
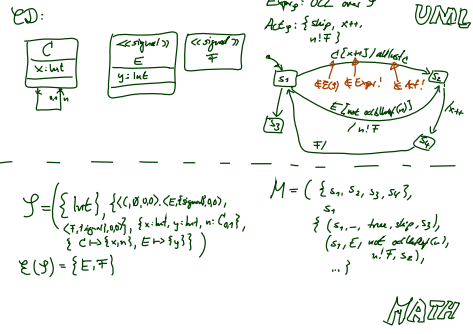
maps to $M(SM) = ((s_1, s_2), s_1, (s_1, \text{event}, \text{guard}, \text{action}, s_2))$

Annotations and Defaults in the Standard

Reconsider the syntax of transition annotations:
 $\text{annot} ::= [[(\text{event}) ['.' (\text{event})]* ['!' (\text{guard})]*] ['*' (\text{action})]]$
and let's play a bit with the defaults:
(empty annot.): $\rightsquigarrow (s_1, -, \text{true}, \text{stop}, s_2)$
 $/$ $\rightsquigarrow (s_1, -, \text{true}, \text{stop}, s_2)$
 $E /$ $\rightsquigarrow (s_1, E, \text{true}, \text{stop}, s_2)$
 $/ \text{act}$ $\rightsquigarrow (s_1, -, \text{true}, \text{act}, s_2)$
 E / act $\rightsquigarrow (s_1, E, \text{true}, \text{act}, s_2)$
 $E [e]$ $\rightsquigarrow (s_1, E, e, \text{act}, s_2)$

In the standard, the syntax is even more elaborate: (we don't discuss those)
• $E(v)$ — when consuming E in object u , attribute v of u is assigned the corresponding attribute of $E \in \mathcal{E}$
• $E(v : \tau)$ — similar, but v is a local variable, scope is the transition

Handwritten notes: we view as an address for $E \in \mathcal{E}$, $F \in \mathcal{E}$, $F \in \mathcal{E}$.



State-Machines belong to Classes

- In the following, we assume that a UML models consists of a set \mathcal{C} of class diagrams and a set \mathcal{M} of state chart diagrams (each comprising one state machines SM).
 - Furthermore, we assume each that each state machine $SM \in \mathcal{M}$ is associated with a class $C_{SM} \in \mathcal{C} \setminus \{ \emptyset \}$.
 - For simplicity, we even assume a bijection, i.e. we assume that each class $C \in \mathcal{C} \setminus \{ \emptyset \}$ has a state machine SM_C and that its class C_{SM_C} is C . If not explicitly given, then this one:
 $SM_C := (\{ s_0 \}, s_0, \text{true}, \text{skip}, s_0)$
- We'll see later that, semantically, this choice does no harm.
- Intuition 1:** SM_C describes the behaviour of the instances of class C .
 - Intuition 2:** Each instance of class C executes SM_C .
- Note:** we don't consider multiple state machines per class. Because later (when we have AND-states) we'll see that this case can be viewed as a single state machine with as many AND-states.

References

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