

Hybrid Systems

Approximation of reachable state sets

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Alongkrit Chutinan and Bruce H. Krogh:

Computing Polyhedral Approximations to Flow Pipes for Dynamic Systems

In Proceedings of the 37rd IEEE Conference on Decision and Control, 1998

Olaf Stursberg and Bruce H. Krogh:

Efficient Representation and Computation of Reachable Sets for Hybrid Systems

Hybrid Systems: Computation and Control, LNCS 2623, pp. 482-497, 2003

We had a look at state set approximations by

- convex polyhedra,

and at the basic operations

- testing for membership,
- intersection, and
- union

on these.

Thus we can

- approximate state sets and
- compute with them.

How is all this used in the reachability analysis procedure?

General reachability procedure

Input: Set **Init** of initial states.

Algorithm:

```
 $R^{\text{new}} := \text{Init};$   
 $R := \emptyset;$   
while ( $R^{\text{new}} \neq \emptyset$ ) {  
     $R := R \cup R^{\text{new}};$   
     $R^{\text{new}} := \text{Reach}(R^{\text{new}} \setminus R);$   
}
```

Output: Set **R** of reachable states.

What is "Reach"?

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For **hybrid systems**, independently of the exact definition of “Reach”, it will involve the following computations:

Given a state set R , compute

- the set of states reachable from R by a **flow** (i.e., time transition),
and
- the set of states reachable from R by a **jump** (i.e., discrete transition).

Computing the jump successors, i.e., the flow pipe, of a set can be done with the operations we already introduced.

The harder part is computing the flow successors. So let's start with that...

Approximating a flow pipe

Consider a dynamical system with **state equation**

$$\dot{x} = f(x(t)).$$

We assume f to be **Lipschitz continuous** so that for every initial state x_0 there is a unique solution $x(t, x_0)$ to the state equation.

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We describe a solution which approximates the flow pipe by a sequence of **convex polytopes**.

Definition (Convex polytope)

Let $POLY(C, d)$ denote the convex polytope defined by the pair $(C, d) \in \mathbb{R}^{m \times n} \times \mathbb{R}^m$ according to

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- Given a finite set of points Γ , the **convex hull $CH(\Gamma)$** of Γ is the smallest convex set that contains Γ .

Problem statement for polyhedral approximation of flow pipes

Given

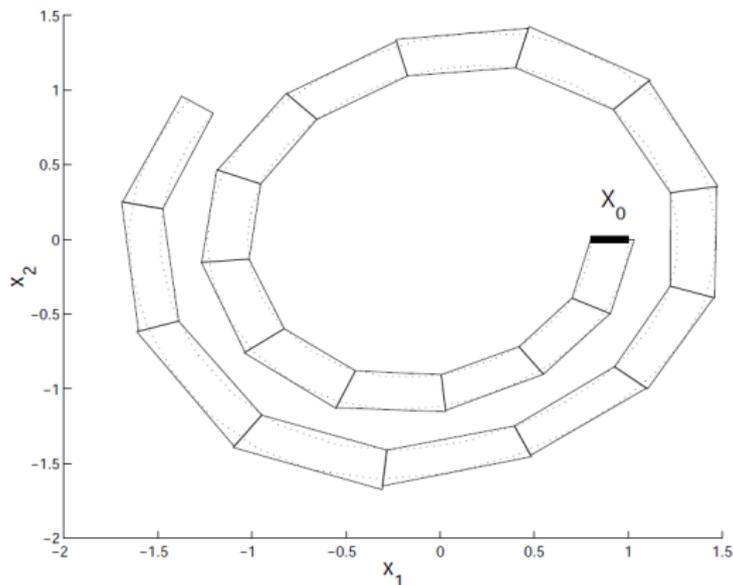
- a set X_0 of initial states which is a polytope, and
- a final time t_f ,

compute a polyhedral approximation $\hat{\mathcal{R}}_{[0,t_f]}(X_0)$ to the flow pipe $\mathcal{R}_{[0,t_f]}(X_0)$ such that

$$\mathcal{R}_{[0,t_f]}(X_0) \subseteq \hat{\mathcal{R}}_{[0,t_f]}(X_0).$$

Flow pipe segmentation

Since a single convex polyhedra would strongly overapproximate the flow pipe, we compute a **sequence of convex polyhedra**, each approximating a **flow pipe segment**.



Segmented flow pipe approximation

Let the time interval $[0, t_f]$ be divided into $0 < N \in \mathbb{N}$ time segments

$$[0, t_1], [t_1, t_2], \dots, [t_{N-1}, t_f]$$

with $t_i = i \cdot \frac{t_f}{N}$.

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We generate an approximation $\hat{\mathcal{R}}_{[t_1, t_2]}(X_0)$ for each flow pipe segment:

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$$\mathcal{R}_{[t_1, t_2]}(X_0) \subseteq \hat{\mathcal{R}}_{[t_1, t_2]}(X_0).$$

The complete **flow pipe approximation** is the union of the approximation of all N pipe segments:

$$\mathcal{R}_{[0, t_f]}(X_0) \subseteq \hat{\mathcal{R}}_{[0, t_f]}(X_0) = \bigcup_{k=1, \dots, N} \hat{\mathcal{R}}_{[t_{k-1}, t_k]}(X_0)$$

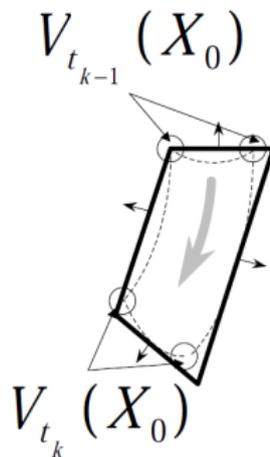
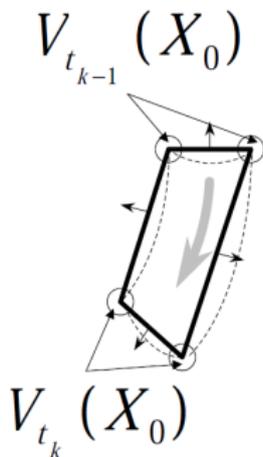
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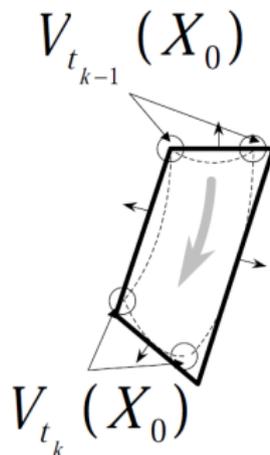
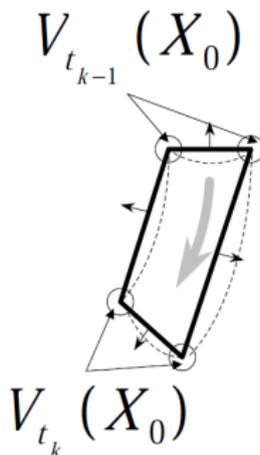
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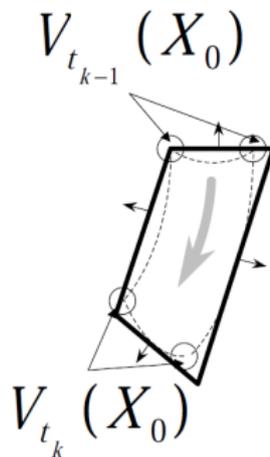
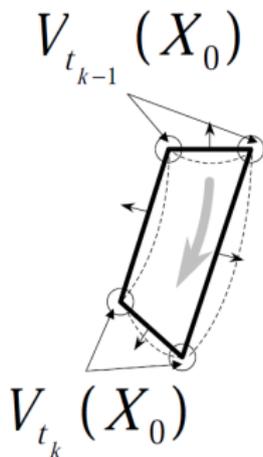
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The approximation of the flow pipe for the time segment $[t_{k-1}, t_k]$ ($k \in \{1, \dots, N\}$) consists of the following steps:

- **Evolve vertices:** Compute the set of points reachable from the vertices of X_0 in time t_{i-1} and in time t_i .
- **Determine hull:** Compute the convex hull of those points.
- **Bloat hull:** Enlarge the hull until it contains all points of the flow pipe segment.



1. Evolve vertices

To gain some geometrical information about the flow pipe segment, we begin with taking sample points at times t_{k-1} and t_k from the trajectories emanating from the vertices of X_0 .

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In particular, we compute the sets $V_{t_{k-1}}(X_0)$ and $V_{t_k}(X_0)$ where

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Each point in the above sets can be obtained

- by analytic solution of the state equation and computing the value, or
- by simulation.

2. Determine hull

We use the evolved vertices in $V_{t_{k-1}}(X_0)$ and $V_{t_k}(X_0)$ to form a **convex hull** which serves as an **initial approximation** to the flow pipe segment $\mathcal{R}_{[t_{k-1}, t_k]}(X_0)$, denoted by

$$\Phi_{[t_{k-1}, t_k]}(X_0) = CH(V_{t_{k-1}}(X_0) \cup V_{t_k}(X_0)).$$

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Note that $\Phi_{[t_{k-1}, t_k]}(X_0)$ may not contain the whole flow pipe segment $\mathcal{R}_{[t_{k-1}, t_k]}(X_0)$.

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Let (C_Φ, d_Φ) be the matrix-vector pair defining the convex hull, i.e.,

$$\Phi_{[t_{k-1}, t_k]}(X_0) = POLY(C_\Phi, d_\Phi).$$

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- Given: $POLY(C_{\Phi}, d_{\Phi})$.
- We want: $\mathcal{R}_{[t_{k-1}, t_k]}(X_0) \subseteq POLY(C_{\Phi}, \boxed{d})$.

3. Bloat hull

- We compute d as the solution to the following optimization problem:

$$\begin{aligned} \min_d \quad & \text{volume}[POLY(C_\Phi, d)] & (1) \\ \text{s.t.} \quad & \mathcal{R}_{[t_{k-1}, t_k]}(X_0) \subseteq POLY(C_\Phi, d). \end{aligned}$$

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- The i th component d_i^* of the optimum d^* can be found by solving

$$\max_x c_i^T x \quad \text{s.t. } x \in \mathcal{R}_{[t_{k-1}, t_k]}(X_0). \quad (2)$$

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- Solution (x_0^*, t^*) to 3 \rightarrow

Solution $x(t^*, x_0^*)$ to 2 \rightarrow

Solution $d_i^* = c_i^T x(t^*, x_0^*)$ to 1.

Example

- Van der Pol equation:

$$\dot{x}_1 = x_2$$

$$\dot{x}_2 = -0.2(x_1^2 - 1)x_2 - x_1.$$

- Initial set: $X_0 = \{(x_1, x_2) \mid 0.8 \leq x_1 \leq 1 \wedge x_2 = 0\}$.
- Time: $t_f = 10$.
- Segments: 20

