Real-Time Systems

Lecture 01: Introduction

2013-04-16

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Contents & Goals

Last Lecture:

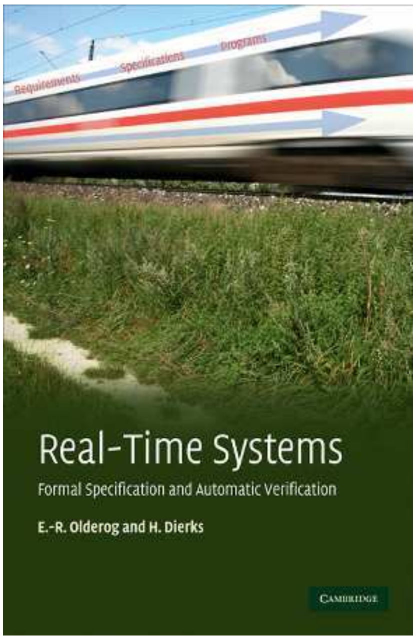
- ./.

This Lecture:

- **Educational Objectives:**
  - Be able to decide whether you want to stay with us or not.
    (IOW: an advertisement for the lecture.)
  - Agree on formalia.
- **Content:**
  - Overview: content (and non-content) of the lecture.
  - Definition reactive, real-time, hybrid system.
  - Outlook on methodology for precise development of (provably) correct real-time systems.
  - Formalia: dates/times, exercises, exam admission.
  - Literature
  - A formal model of real-time behaviour.

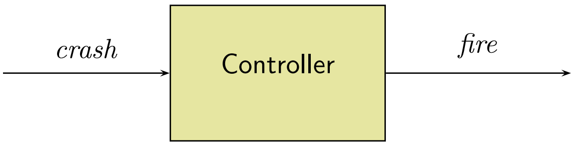
Introduction

Subject of the Lecture



What is a Real-Time System?

Classical example: **Airbag Controller**



Requirement: "When a crash is detected, fire the airbag."

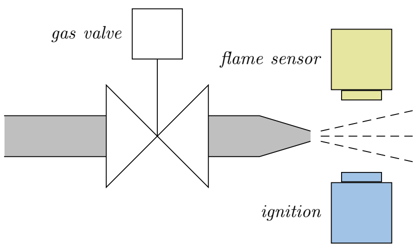
- When firing **too early**: airbag ineffective.
- When firing **too late**: additional threat.

Say, 300ms (plus/minus small ε) after a crash is the right™ time to fire.

Then the **precise requirement** is
"When a crash is detected at time t, fire the airbag at $t + 300ms \pm \varepsilon$."

What is a Real-Time System?

- Other example: **Gas Burner**



- **Leakage** is practically unavoidable:
  - for ignition, first open valve
  - then ignite the available gas
  - ignition may fail...
- **Leakage** is **safety critical**:
  Igniting large amounts of leaked gas may lead to a dangerous explosion.

No, Really, What is a Real-Time System?

- The examples have in common that

  it matters, **when in time**

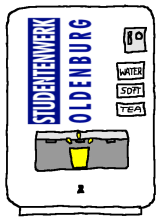
  the output for a given input (sequence) takes place.

  For instance,
  - "fire" 300ms after "crash",
  - within any interval of at least 60s, leakage (= have the gas valve open without a flame) amounts to at most 5% of the time.

  before / after / in between

  Note: **quantitative** (here) vs. **qualitative** notions of time (untimed).

- Often: There is a physical environment, which has a notion of time, and which evolves while our controller is computing.



- (Half-)**Contrast:** vending machine for soft-drinks:
  - If the customer is really thirsty, she'll wait.
  - Neither the usage of a really fast or a really slow contemporary controller causes a violation of (timing) requirements.
- (Real) **Contrast:** transformational systems, such as computing π.

Other Definitions *[Douglass, 1999]*

- "A **real-time** system is one that has **performance deadlines** on its computations and actions."
- Distinguish:
  - "**Hard** deadlines: performance requirements that **absolutely must** be met each and every event or time mark."

    "(Late/early data can be bad data.)"
  - "**Soft** deadlines: for instance about **average** response times."

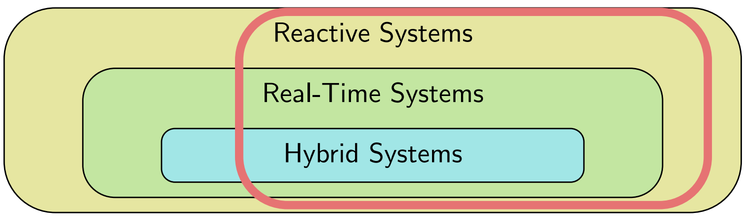
    "(Late data is still good.)"
- Design Goal:

  A **timely system**, i.e. one meeting its performance requirements.
- Note: **performance** can in general be any unit of quantities:
  - (discrete) number of steps or processor instructions,
  - (discrete or continuous) number of seconds,
  - etc.

Definitions: Reactive vs. Real-Time vs. Hybrid Systems

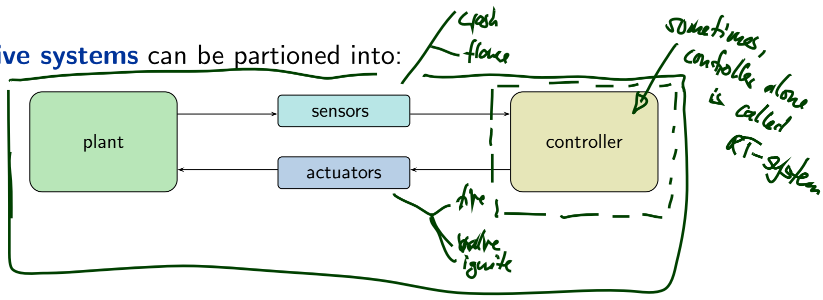
- **Reactive Systems** interact with their environment by reacting to inputs from the environment with certain outputs.
- A **Real-Time System** is a **reactive system** which, for certain inputs, has to compute the corresponding outputs within given time bounds.
- A **Hybrid System** is a **real-time system** consisting of continuous and discrete components. The continuous components are time-dependent (!) physical variables ranging over a continous value set.

- A system is called **Safety Critical** if and only if a malfunction can cause loss of goods, money, or even life.



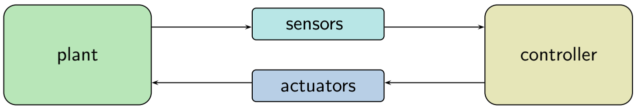
The Problem: Constructing Safety-critical RT Systems

- **Reactive systems** can be partioned into:



The Problem: Constructing Safety-critical RT Systems

- **Reactive systems** can be partioned into:



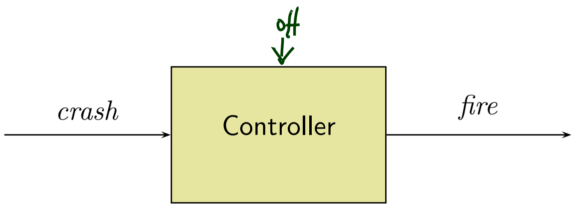
  - "In constructing a **real-time system** the aim is to control a physically existing environment, the **plant**, in such a way that the controlled plant satisfies all desired (timing) requirements."

- The design of **safety critical** (**reactive**) **systems** requires a high degree of precision:
  We want — at best — to be sure that a design meets its requirements.

- **Real-time systems** are often **safety-critical**.

- The lecture presents approaches for the precise development of **real-time systems** based on formal, mathematical methods.

Constructing Safety-critical RT Systems: Examples



"When a crash is detected at time t, fire the airbag at $t + 300ms \pm \varepsilon$."

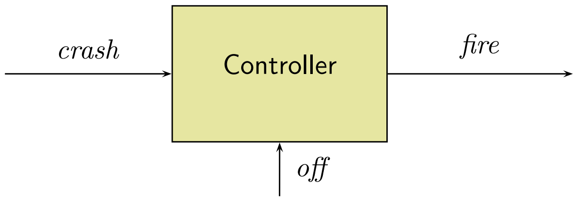
- A controller program is easy:

```
while (true) do
  poll_sensors();
  if (crash) tmr.start(300ms);
  if (tmr.elapsed()) fire := 1;
  update_actuators();
od
```

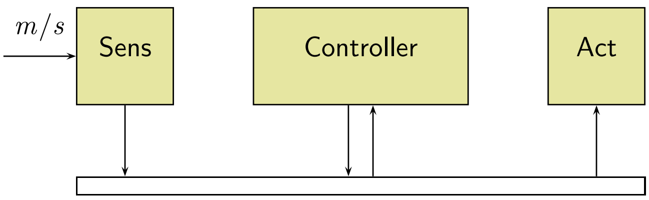
- And likely to be believed to be correct.

Constructing Safety-critical RT Systems: Examples

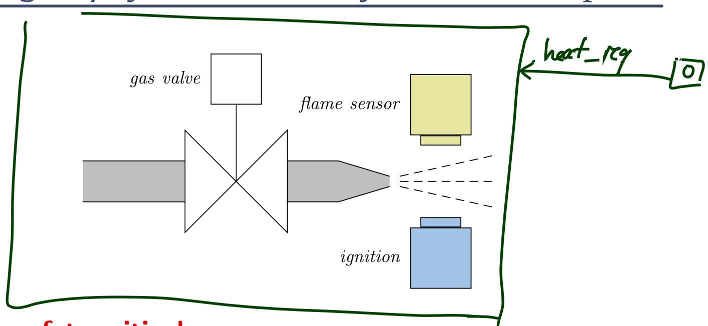
- More complicated: **additional features**.



- More complicated: **distributed implementation**.



Constructing Safety-critical RT Systems: Examples



- Leakage is **safety critical**:
  Igniting large amounts of leaked gas may lead to a dangerous explosion.

- Controller program for ignition is easy:

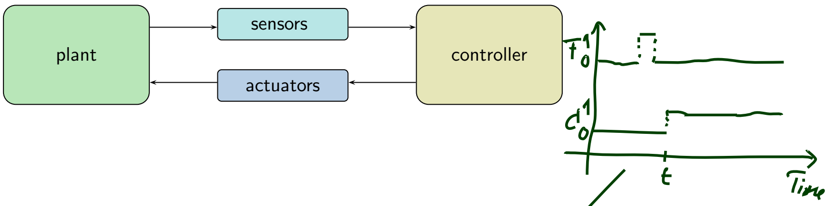
```
while (!flame) do
  open_valve();
  wait(t);
  ignite();
od
```

wait(tsafe);

if (!flame) {
close_valve();
wait(tsafe);
}

close_valve();
wait(tsafe);

- Is it **correct**? (Here: Is it avoiding dangerous explosions?)



To

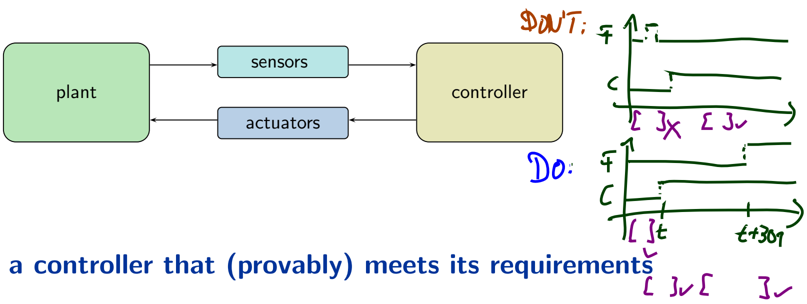
design a controller that (provably) meets its requirements

we need

- a formal model of behaviour in (quantitative) time,
-
-
-

Then we can devise **a methodology** to get **from requirements** to a **(correct) implementation** — here: following [Olderog and Dierks, 2008].

Prerequisites for Precise Development of Real-Time Systems



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-
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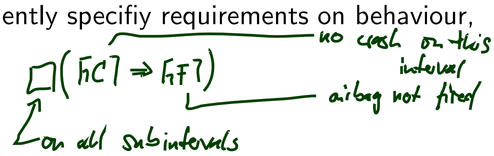

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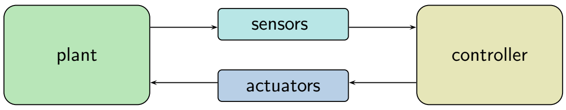


To

design a controller that (provably) meets its requirements

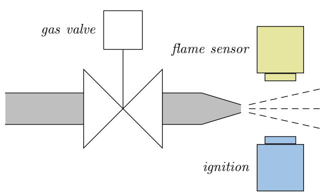
we need

- a formal model of behaviour in (quantitative) time,
- a language to concisely, conveniently specifiy requirements on behaviour,
- a language to specify behaviour of controllers,
- a notion of "meet" and a methodology to verify (or prove) "meeting".

Then we can devise **a methodology** to get **from requirements** to a **(correct) implementation** — here: following [Olderog and Dierks, 2008].

Sketch of the Methodology: Gas Burner Example

- **Requirements**
  - At most 5% of any at least 60s long interval amounts to leakage.
- **Reflective Design**
  - Time intervals with leakage last at most 1s.
  - After each leak, wait 30s before opening valve again.
- **Constructive Design**
  - PLC Automaton
    (open valve for 0.5s;
    ignite;
    if no flame after 0.1s close valve)
- **Implementation**
  - IEC 61131-3 program



Content Overview

Content

Introduction

- **First-order Logic**
- **Duration Calculus** (DC)
- Semantical Correctness Proofs with DC
- DC Decidability
- DC Implementables
- **PLC-Automata**

- **Timed Automata** (TA), Uppaal
- Networks of Timed Automata
- Region/Zone-Abstraction
- Extended Timed Automata
- Undecidability Results

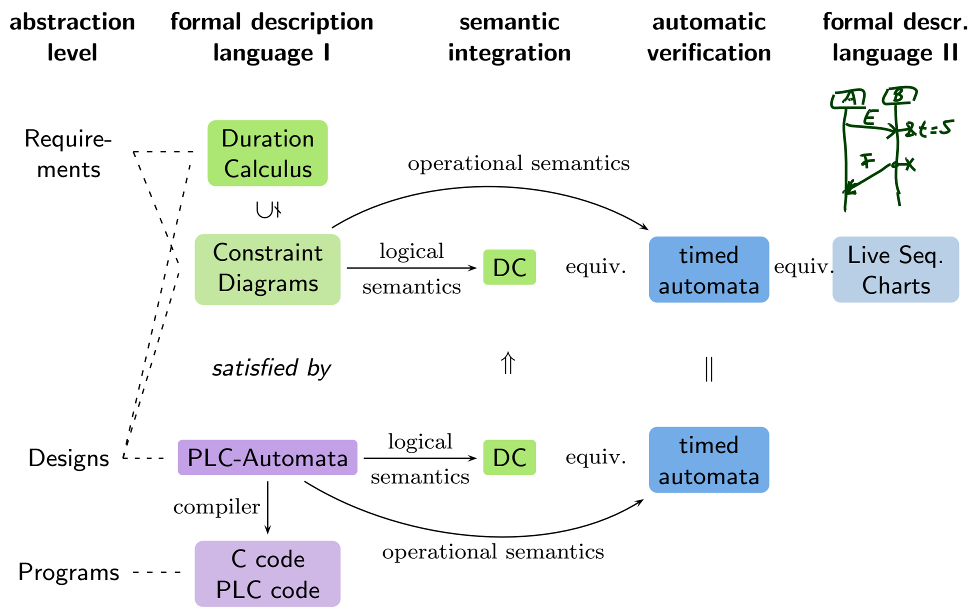
$obs : \mathsf{Time} \to \mathscr{D}(obs)$

$\langle obs_0, \nu_0 \rangle, t_0 \xrightarrow{\lambda_0} \langle obs_1, \nu_1 \rangle, t_0 \xrightarrow{\lambda_1} \langle obs_1, \nu_2 \rangle, t_1, \ldots$

- **Automatic Verification**...
- ...whether TA satisfies DC formula, observer-based

Recap

Tying It All Together



Maybe-Content

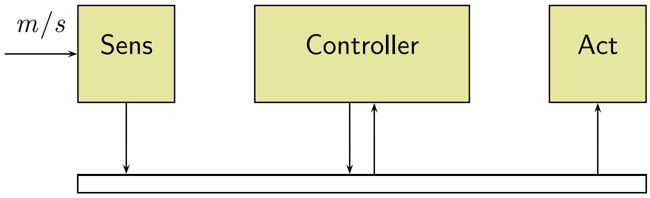
- **Worst Case Execution Time**
  - Recall over-simplified airbag controller:

```
while (true) do
  poll_sensors();
  if (crash) tmr.start(300ms);
  if (tmr.elapsed()) fire := 1;
  update_actuators();
od
```

  - The execution of `poll_sensors()` and `update_actuators()` also takes time! (And we have to consider it!)
  - **Maybe in lecture**:
    How to determine the WCET of, for instance, C code.
    (A science of its own.)

Non-Content

- **Scheduling**
  - Recall over-simplified airbag controller:



  - **Not in lecture**: Specialised methods to determine...
    - ...whether the bus provides sufficient bandwidth.
    - ...whether the Real-Time OS controlling CPU 'Controller' schedules the airbag control code in time.
    - ...how to distribute tasks over multiple CPUs.
    - etc.

  (Also a science of its own.)

Formalia

Formalia: Event

- **Lecturer:** Dr. Bernd Westphal
- **Support:** Dennis Gauss
- **Homepage:**

  http://swt.informatik.uni-freiburg.de/teaching/SS2013/rtsys

- **Questions:**
  - **"online"**:
    (i) ask immediately or in the break
  - **"offline"**:
    (i) try to solve yourself
    (ii) discuss with colleagues
    (iii) contact lecturer by mail (cf. homepage) or just drop by:
    Building 52, Room 00-020

Formalia: Dates/Times, Break

- **Schedule:**

Wednesday,	week N:	10–12 **lecture**	(exercises M **online**)
Tuesday,	week $N+1$:	14–16 **lecture**	
Wednesday,	week $N+1$:	10–12 **lecture**	
Monday,	week $N+2$:	14:00	(exercises M **early turn-in**)
Tuesday,	week $N+2$:	14–16 **tutorial**	(exercises M **late turn-in**)
Wednesday,	week $N+2$:	10–12 **lecture**	(exercises $M+1$ **online**)

  With a prefix of lectures, with public holidays; see homepage for details.

- **Location:**
  - Tuesday, Wednesday: here

- **Break:**
  - Unless a majority objects **now**,
    we'll have a **10 min. break** in the middle of each event from now on.

Formalia: Lectures

- **Course language:** **English**
  (slides/writing, presentation, questions/discussions)
- **Presentation:**
  half slides/half on-screen **hand-writing** — for reasons
- **Script/Media:**
  - slides without annotations on **homepage**,
    trying to put them there **before** the lecture
  - slides with annotations on **homepage**, 2-up for printing,
    typically soon **after** the lecture
  - recording on eLectures portal with max. 1 week delay
    (link on **homepage** – eLectures is updated first, look there!)
- **Interaction:**
  absence often moaned but **it takes two**,
  so please ask/comment immediately

Formalia: Exercises and Tutorials

- **Schedule/Submission:**
  - Recall: exercises **online** on Wednesday before (or soon after) lecture,
    regular **turn in** on corresponding tutorial day until **14:00 local time**
  - should work in groups of **max. 3**, clearly give **names** on submission
  - please submit **electronically** by Mail to **me** (cf. homepage),
    some LaTeX styles on homepage; paper submissions are tolerated
- **Didactical aim:**
  - deal more extensively with notions from lecture (easy)
  - explore corner cases or alternatives (medium)
  - evaluate/appreciate approaches (difficult)
  - additional **difficulty**: imprecise/unclear tasks — by intention
- **True aim:** **most complicated** rating system **ever**, namely two ratings
  - Good-will ("reasonable solution with knowledge **before** tutorial")
  - Evil/Exam ("reasonable solution with knowledge **after** tutorial")

  10% **bonus** for **early** submission.

Formalia: Exam

- **Exam Admission:**

  50% of the maximum possible non-bonus **good-will points** in total are **sufficient** for admission to exam

- **Exam Form:** (oral or written) not yet decided

Formalia: Evaluation

Speaking of **grading and examination**...

- **Mid-term Evaluation:**
  We will have a **mid-term evaluation**[a], but we're **always** interested in comments/hints/proposals concerning form or content.

[a] that is, students are asked to evaluate lecture, lecturer, and tutor...

Formalia: Questions?

References

References

[Douglass, 1999] Douglass, B. P. (1999). *Doing Hard Time*. Addison-Wesley.

[Olderog and Dierks, 2008] Olderog, E.-R. and Dierks, H. (2008). *Real-Time Systems - Formal Specification and Automatic Verification*. Cambridge University Press.