Formal Methods for C

Seminar – Summer Semester 2014

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Plan

- **(15 min.)** Topic lottery: prepare lottery ticket with *P.S.*
  - name
  - first preference topic
  - second preference topic

▷ Introduction to C (1)

- **(15 min.)** The VM (Marius)

Overview

Goals

Educational Objectives: Capabilities for following tasks/questions.

- Rough overview over concepts of the C programming language.
  - IOW: know, what to look for in books/manuals.
    (We try to stick with the names from of ISO/IEC 9899:1999)
  - IOW: don't be badly surprised from the examples.
- The concept of pointers.
- Basic work-flow, tool-usage (headers, sources, compiler, linker).
- Orthogonal: Rough overview over common sources of errors.
- Formal methods and C.
- **Not:** reference manual, each and every feature.
- **Not:** programming course.

Everybody Say "Hello" to C

```
1   #include <stdio.h>
2
3   int g( int x ) {
4     return x/2;
5   }
6
7   int f() {
8     return g(1);
9   }
10
11  int main() {
12    printf( "Hello World.\n" );
13    return f();
14  }
```
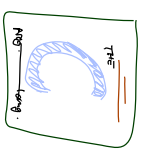
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42 Years of C

1972: Created by **Dennis Ritchie** († 2011) for Unix system programming.

1978: Brian W. Kernighan & Dennis Ritchie: "The C Programming Language" – "K&R C".

1989: ANSI X3.159-1989 – C89, **C90** (still most widely used (?))

1999: ISO/IEC 9899:1999 – C99, (use --std=c99 for gcc(1))

2011: ISO/IEC 9899:2011 – C11

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1972: Created by **Dennis Ritchie** († 2011) for Unix system programming.
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1989: ANSI X3.159-1989 – C89, **C90**.
1999: ISO/IEC 9899:1999 – C99, (use –std=c99 for gcc4))
2011: ISO/IEC 9899:2011 – C11

Content

- Compilers for virtually every platform (CPU + operating system) available. Virtually every CPU vendor offers an own C compiler, in particular in the embedded domain (MSP 430, ARM, intel...).
- Still No. 1 programming language for embedded systems software, hardware drivers, performance critical applications, ...
- Preferred by many embedded programmers for "lack of surprises"; (without optimisation) direct correspondence between C code and assembler.
- Resources widely controllable by programmer, downside: programmer needs to "know what one's doing".

Comments

Comments (6.5.9)

- one line comment, until end of line: `// ...`
- generic comment, no nesting: `/* ... */`
- corner cases:

```
1   "a//b"              // four-character string literal
2   #include "//e"      // undefined behavior
3   /* */               // comment, not syntax error
4   f = g/**/h;         // equivalent to f = g / h;
5   //\
6   i();                // part of a two-line comment
7   \
8   j();                // part of a two-line comment
9   #define glue(x,y) x##y   // syntax error, not comment
10  glue(/,/) k();      // equivalent to !();
11  /*/*/ l();          // equivalent to !();
12  m = n//**/o         // equivalent to m = n + p;
13      + p;
```

Declarations and Scopes

Variables

Basic Types (6.2.5), Constants (6.4.4)

```
1  char c = 'a', d = 93;
2
3  int x = 027; // octal!
4  long int y = 3L, z;
5  short int w = 0xBEEF;
```

"[...] **char** is large enough to store any member of the basic exec. character set."

"A 'plain' **int** object has the natural size suggested by the architecture of the execution environment (large enough [for values] **INT_MIN** to **INT_MAX**)."

char, short int, int, long int, long long int also as **unsigned**.

```
1  unsigned int x = 27U;
```

```
1  float f = 1;
2  double g = 314159e-5;
3  long double h;
```

Bool (6.2.5)

- Only introduced in C99.
- "An object declared as type **_Bool** is large enough to store the values 0 and 1."
- **<stdbool.h>** (→ later) defines **bool, true, false** as macros (→ later).

- Before C99, and **still very common**:
- Use scalar type (including pointers).
- **0**: **false**
- **everything else**: **true**
- values of boolean expressions: 0, 1

```
1  int y = 27;
2  int x = 13 && (y > 0); // value of x becomes 1
```

Derived Types (6.2.5), (6.7)

- **array types**:

```
1  int a[10];
2  char b[2][3];
```

- **structured types**:

```
1  typedef struct {
2    int n;
3    double d;
4  } S; // declaration of type 'S'
5  S x; // declaration of variable 'x' of type 'S'
6
7  struct { int a; double b; } y; // declaration of ...?
8  typedef struct { S[3] c; double d; } T;
```

Note: x and y are of different type!

- **union types**: not here
- **function types**, **pointer types**: later

Declaration Syntax

- ...takes a little bit getting used to:

```
1  int c, *p, a[3], *q[2], (*f)(int);
```

- is a shorthand notation for:

```
1  int c;          // integer
2  int *p;         // pointer to integer
3  int a[3];       // array of 3 integers
4  int *q[2];      // array of 2 pointers to integer
5  int (*f)(int);  // function pointer...
```

- And what's declared here, what is its type?

```
1  int (*g)( int (*)( int*[2]) ); // ...?
```

Expressions

Expressions (6.5)

- "An **expression** is a sequence of operators and operands that specifies computation of a value, or that designates an object or a function, or that generates side effects, or that performs a combination thereof."

- basically like Java:
- postfix operators: a++, p->g->y
                                             (double)a
- unary operators: ++a, sizeof(int), &a, *p
- cast operators:
- multiplicative, additive
- relational, equality: a < b, a == b, a != b
- logical operators: (a < b) && (c > 0)
- conditional operator: a < b ? a : b
- assignment operator (are expressions):   a = b, a += b, a = b = 0
- comma operator: a = b, c = d

Statements

Bitwise Operators (6.5.3.3, 6.5.7, 6.5.10–12)

- Often used in hardware level programming: Communicate with "the hardware" via memory-mapped registers – single bits or groups of bits have particular, platform dependent meaning.

- Bitwise And, Or, Xor (6.5.10–12):

$$0101_2 \;\&\; 1100_2 == \ldots$$
$$0101_2 \;\&\&\; 1100_2 == \ldots$$

- Useful idioms (assuming 4-bit type):
  - Set the 3rd bit: $a \mathrel{|=} 0100_2$
  - Clear the 2nd bit: $a \mathrel{\&=} 1101_2$
  - Test whether 2nd bit set: $a \;\&\; 0010_2$

- Shift (6.5.7): $a << 2$, $a >> 2$ (unsigned (!) filled up with 0 at left and right)

- Bitwise complement (6.5.3.3): $\sim a$

- **Be careful with signed types** (bit-operations at best on unsigned):

$$\sim(\texttt{int}(1)) == \texttt{0xFFFFFFFE} \; \textit{// == -2}$$

Lvalues (6.3.2.1)

```
1  int x, a[3], *p;
2
3  x = 27;
4  a[1] = 0;
5  p = &x;
6  *p = 13;
7
8  &x = ...; // no
9  a = ...; // no, only as initializer
```

- "An **lvalue** is an expression with an object type or an incomplete type other than void."

- "The name "lvalue" [comes from] E1 = E2, in which the **left** operand E1 is required to be a (modifiable) lvalue. What is sometimes called "rvalue" is in this International Standard described as the "value of an expression". An obvious example of an lvalue is an identifier of an object."

Boolean Logic (6.3.2.1, 6.5.13–6.5.15)

- "When any scalar value is converted to _Bool, the result is 0 [**false**] if the value compares equal to 0; otherwise, the result is 1." (6.3.2.1)

- (a != 0) and 'a' are equivalent (if a is of scalar type), so are a == 0 and !a
  null check

- for pointers (later): p == NULL and !p are equivalent

Statements (6.8)

- "A **statement** specifies an action to be performed. Except as indicated, statements are executed in sequence."

- also basically like Java:
  - Selection statements (6.8.4):   if, else, switch
  - Iteration statements (6.8.5):   while, do ... while, for
  - Jump statements (6.8.6):   goto, continue, break, return

Functions

Function Definitions (6.9.1)

```
1  int max( int a, int b )
2  {
3    return a > b ? a : b;
4  }
```

(handwritten: decl. + def.)

- no nesting, no member functions
- all in file (global) scope (but: module scope possible (later))
- call-by-value semantics (call-by-reference: later)

Function declaration (vs. **definition**):

```
1  int max( int a, int b ); // param. names just "decoration"
2  int max( int, int );
```

- "Zero or many declarations, exactly one definition"

Scopes

Scopes of Identifiers (6.2.1)

```
1   int a; // file scope (F)
2
3   struct { int a; } s; // name-space
4
5   int f( int a ) // block scope, block (A)
6   { /* <--          uses a:(A)
7     if ( a ) { /* <--          (A) */
8                           (B) */
9       a = 0; // uses a:(A)
10      int a = 27; // block scope, block (B)
11
12      s.a = a; // uses a:(B)
13    }
14    return a; // uses a:(A)
15  }
16
17  int main() { return f(a); /* uses a:(F) */ }
```

- "The same identifier can denote different entities at diff. points in the program."
- "For each different entity that an identifier designates, the identifier is **visible** (i.e. can be used) only within a region of program text called its **scope**."

Scopes of Identifiers (6.2.1)

- "**Different entities** designated by the **same identifier** either have different scopes, or are in different name spaces. **There are four kinds of scopes:** function, file, block, and function prototype."
- "A label name is the only kind of identifier that has function scope."
- "Every other identifier has scope determined by the placement of its declaration (in a declarator or type specifier)."
- **Declare before use:** each identifier must be declared before (i.e. earlier in the source file) its first use in, e.g., an expression. (**Unlike Java!**)

Scopes of Identifiers (6.2.1)

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- **Declare before use:** each identifier must be declared before (i.e. earlier in the source file) its first use in, e.g., an expression. (**Unlike Java!**)
- "**Hidden**" identifiers are **not accessible:**

```
1  int a; /* a (F) */
2
3  void f( int a ) { /* a (A) */
4    a = 0; // uses a:(A).
5    // a:(F) not accessible here. "hidden" by a:(A)
6  }
```

Pointers

Variables in the System's Memory

```
1  char c = 127;
```

the compiler chose to
store values of 'c' at
memory cell with
address 0x1001

Assigning Variables = Update Memory

```
1  char c = 127;
2  c = c + 1;
```

"$c = c + 1$" means: the new
value of c is the old value
plus 1; in assembler:
read 0x1001, R; inc R;
write R, 0x1001

A Pointer to 'c' (16-bit Architecture)

```
1  char c = 127;
2  c = c + 1;
3  char* p = &c;
```

the compiler chose to store
values of variable 'p' at
memory cells (!) with
addresses 0x1002 and 0x1003

Dereferencing Pointers

```
1  char c = 127;
2  c = c + 1;
3  char* p = &c;
4  *p = *p + 3;
```

dereference

Assigning Variables = Update Memory

```
1  char c = 127;
2  c = c + 1;
```

"$c = c + 1$" means: the new
value of c is the old value
plus 1; in assembler:
read 0x1001, R; inc R;
write R, 0x1001

Dereferencing Pointers

```
1  char c = 127;
2  c = c + 1;
3  char* p = &c;
4  *p = *p + 3;
```

*p, rhs: get the value of
p (0x1001) and read the
value **at that address** (at
0x1001, yields 0x80)

add 3 to the value just
obtained (yields 0x83)

*p, lhs: get the value of
p (0x1001) again, write the
addition result (0x83) **to
that address** (to 0x1001)

Dereferencing Pointers

```
1  char c = 127;
2  c = c + 1;
3  char* p = &c;
4  *p = *p + 3;
5  char* q = p;
```

*p, rhs: get the value of *p* (0x1001) and read the value **at that address** (at 0x1001, yields 0x80)

*p, lhs: get the value of *p* again (0x1001), write the addition result (0x83) **to that address** (to 0x1001)

add 3 to the value just obtained (yields 0x83)

0x1000	0x1001	...
	0x83	
	0x01	

Assigning Pointers

```
1  char c = 127;
2  c = c + 1;
3  char* p = &c;
4  *p = *p + 3;
5  char* q = p;
```

assume the compiler choses to store values of variable **'q'** at memory cells (!) with addresses 0x1004 and 0x1005

0x1000	0x1001	...	0x1004	0x1005
	0x80		??	??
	0x10		??	??
	0x01			

Assigning Pointers

```
1  char c = 127;
2  c = c + 1;
3  char* p = &c;
4  *p = *p + 3;
5  char* q = p;
```

0x1000	0x1001	0x1002	0x1003	0x1004	0x1005
	0x83	0x10	0x01	0x10	0x01

```
P = *P + 3;
 ↳ 0x1002   0x1005
   0x83      0x10
   0x80      0x83
*P = 0;
```

Pointers to Pointers

```
1  char c = 127;
2  c = c + 1;
3  char* p = &c;
4  *p = *p + 3;
5  char* q = p;
6  char** r = &q;
```

assume the compiler choses to store values of variable **'r'** at memory cells (!) with addresses 0x1006 and 0x1007

0x1000	0x1001	0x1002	0x1003	0x1004	0x1005	0x1006	0x1007
	0x80	0x10	0x01	0x10	0x01	??	??

Pointers to Pointers

```
1  char c = 127;
2  c = c + 1;
3  char* p = &c;
4  *p = *p + 3;
5  char* q = p;
6  char** r = &q;
```

0x1000	0x1001	0x1002	0x1003	0x1004	0x1005	0x1006	0x1007
	0x83	0x10	0x01	0x10	0x01	0x10	0x04

Using Pointers to Pointers

```
1  char c = 127;
2  c = c + 1;
3  char* p = &c;
4  *p = *p + 3;
5  char* q = p;
6  char** r = &q;
7  *r = p;
8  **r = 5;
```

0x1000	0x1001	0x1002	0x1003	0x1004	0x1005	0x1006	0x1007
	0x83	0x10	0x01	0x10	0x01	0x10	0x04

Using Pointers to Pointers

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1  char c = 127;
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7  *r = p;
8  **r = 5;
```

Using Pointers to Pointers

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1  char c = 127;
2  c = c + 1;
3  char* p = &c;
4  *p = *p + 3;
5  char* q = p;
6  char** r = &q;
7  *r = p;
8  **r = 5;
```

Pointers vs. Arrays

Arrays

```
1  char a[5] = { 'H', 'e', 'l', 'l', 'o' };
```

Arrays

reserve some space
for 5 **chars**...

...and let **a** point
to that space

```
1  char a[5] = { 'H', 'e', 'l', 'l', 'o' };
```
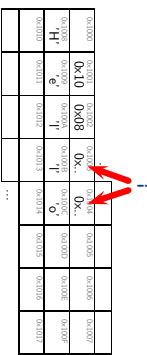
Arrays

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```
1  char a[5] = { 'H', 'e', 'l', 'l', 'o' };
```

Arrays

```
1  char a[5] = { 'H', 'e', 'l', 'l', 'o' };
2  int i;
3  for (i = 0; i < 5; ++i)
4    a[i] = 'x';
```
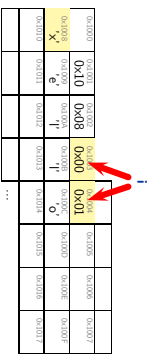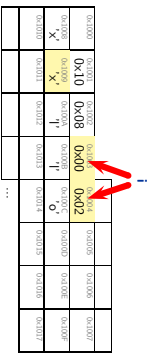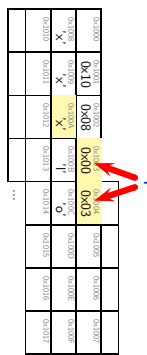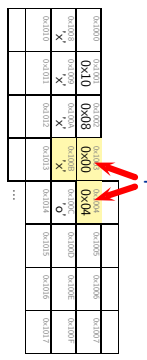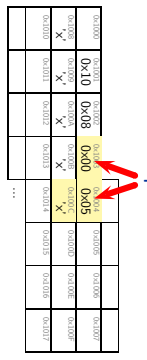
Arrays

```
1  char a[5] = { 'H', 'e', 'l', 'l', 'o' };
2  int i;
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```

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Arrays vs. Pointers

```
1  char a[5] = { 'H', 'e', 'l', 'l', 'o' };
2  char* p = a; // not &a !
3  for (int i = 0; i < 5; ++i, ++p)
4    *p = 'o';
```

Arrays vs. Pointers

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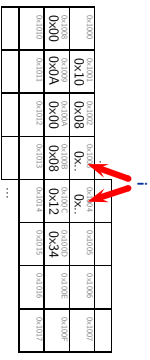

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```

p

Integer Arrays

reserve some space
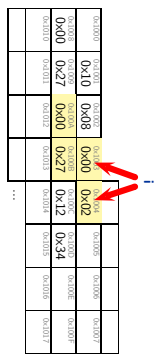for 3 **int**s…

...and let a point
to that space

a

```
1  int a[3] = { 10, 010, 0x1234 };
```

Integer Arrays

```
1  int a[3] = { 10, 010, 0x1234 };
2  int i;
3  for (i = 0; i < 3; ++i)
4    a[i] = 0x27;
```
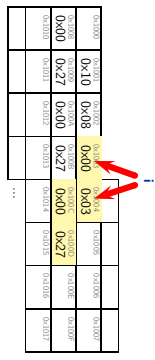
i

Integer Arrays

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2  int i;
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```

i

Integer Arrays vs. Pointers

```
1  int a[3] = { 10, 010, 0x1234 };
2  int* p = a;
3  for (int i = 0; i < 3; ++p)
4    *p = 0x3421;
```

Integer Arrays vs. Pointers

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1  int a[3] = { 10, 010, 0x1234 };
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4    *p = 0x3421;
```

Pointers to 'void', Pointer Arithmetic

Pointer to 'void'

```
1  int[3] a = { 10, 010, 0x1234 };
2  int* p = a;
3  void* q = a;
4  for ( int i = 0; i < 3; ++i ) {
5      p++;
6      q++;
7  }
```

p
q

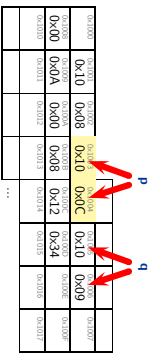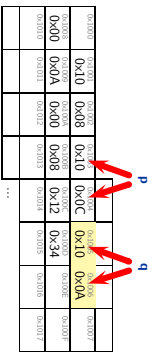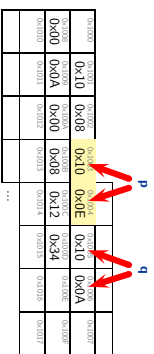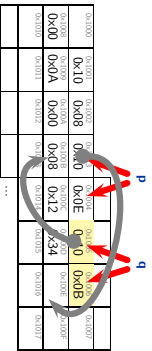## Pointer to 'void'

```
1  int[3] a = { 10, 010, 0x1234 };
2  int* p = a;
3  void* q = a;
4  for ( int i = 0; i < 3; ++i ) {
5      p++;
6      q++;
7  }
```

p
q

Pointer to 'void'

```
1  int[3] a = { 10, 010, 0x1234 };
2  int* p = a;
3  void* q = a;
4  for ( int i = 0; i < 3; ++i ) {
5      p++;
6      q++;
7  }
```

p
q

Pointer to 'void'

```
1  int[3] a = { 10, 010, 0x1234 };
2  int* p = a;
3  void* q = a;
4  for ( int i = 0; i < 3; ++i ) {
5      p++;
6      q++;
7  }
```

p
q

Pointer to 'void'

```
1  int[3] a = { 10, 010, 0x1234 };
2  int* p = a;
3  void* q = a;
4  for ( int i = 0; i < 3; ++i ) {
5      p++;
6      q++;
7  }
```

p
q

Pointer to 'void'

```
1  int[3] a = { 10, 010, 0x1234 };
2  int* p = a;
3  void* q = a;
4  for ( int i = 0; i < 3; ++i ) {
5      p++;
6      q++;
7  }
```

p
q

Pointer to 'void'

```
1  int[3] a = { 10, 010, 0x1234 };
2  int* p = a;
3  void* q = a;
4  for ( int i = 0; i < 3; ++i ) {
5    p++;
6    q++;
7  }
```

Pointers: Observation

- A variable of pointer type just **stores an address**.
- So do variables of **array type**.
- Pointers can point to a certain type, or to **void**.
- "A pointer to void shall have the same representation and alignment requirements as a pointer to a character type." (6.2.5.26)
- The effect of "incrementing" a pointer depends on the type pointed to.

```
1  int a[2];
2  int* p = a;
3  ++p; // points to a[1]
4  void* q = a;
5  q += sizeof(int); // points to a[1]
6  ++q; // may point into the middle
```

Pointer Arithmetic

```
1  int[3] a = { 10, 010, 0x1234 }, i = 0;
2  int* p = a; // not &a !
3
4  if (a[0]    == *p)     i++;
5  if (a[1]    == *(p+1)) i++;
6  if (a[2]    == *(p+2)) i++;
7
8  if (&(a[2]) - p == 2)  i++;
9
10 void* q = a;
11
12 if (a[2]    == *((int*)(q + (2 * sizeof(int))))) i++;
13
14
15 // i == 5
```

void as such does not have values, we need to **cast** 'q' here... note: **void***
can be casted to everything

Pointers for Call By Reference

Call By Reference with Pointers

```
1  void f( int x, int y ) {
2    x++; y++;
3  }
4  void g( int* p, int* q ) {
5    (*p)++; (*q)++;
6  }
7  int a = 2, b = 5;
8  f( a, b );
9  g( &a, &b );
```
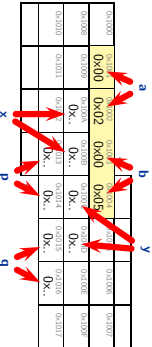
Call By Reference with Pointers

```
1  void f( int x, int y ) {
2    x++; y++;
3  }
4  void g( int* p, int* q ) {
5    (*p)++; (*q)++;
6  }
7  int a = 2, b = 5;
8  f( a, b );
9  g( &a, &b );
```

Call By Reference with Pointers

```
1  void f( int x, int y ) {
2    x++, y++;
3  }
4  void g( int* p, int* q ) {
5    (*p)++, (*q)++;
6  }
7  int a = 2, b = 5;
8  f( a, b );
9  g( &a, &b );
```

Call By Reference with Pointers

```
1  void f( int x, int y ) {
2    x++, y++;
3  }
4  void g( int* p, int* q ) {
5    (*p)++, (*q)++;
6  }
7  int a = 2, b = 5;
8  f( a, b );
9  g( &a, &b );
```

Call By Reference with Pointers

```
1  void f( int x, int y ) {
2    x++, y++;
3  }
4  void g( int* p, int* q ) {
5    (*p)++, (*q)++;
6  }
7  int a = 2, b = 5;
8  f( a, b );
9  g( &a, &b );
```

Call By Reference with Pointers

```
1  void f( int x, int y ) {
2    x++, y++;
3  }
4  void g( int* p, int* q ) {
5    (*p)++, (*q)++;
6  }
7  int a = 2, b = 5;
8  f( a, b );
9  g( &a, &b );
```

Call By Reference with Pointers

```
1  void f( int x, int y ) {
2    x++, y++;
3  }
4  void g( int* p, int* q ) {
5    (*p)++, (*q)++;
6  }
7  int a = 2, b = 5;
8  f( a, b );
9  g( &a, &b );
```

Call By Reference with Pointers

```
1  void f( int x, int y ) {
2    x++, y++;
3  }
4  void g( int* p, int* q ) {
5    (*p)++, (*q)++;
6  }
7  int a = 2, b = 5;
8  f( a, b );
9  g( &a, &b );
```
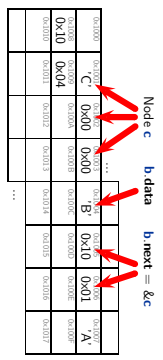
A Linked List

```
1  typedef struct Node {
2    char data;
3    struct Node* next;
4  } Node;
5
6  Node c = { 'C', 0 };
7  Node b = { 'B', &c };
8  Node a = { 'A', &b };
```


Node c   b.data   b.next = &c

data: 'A'
next: 0x004

data: 'B'
next: 0x001

data: 'C'
next: 0x0000

Dynamic Storage & Storage Duration

Dynamic Storage Allocation

Dynamic Storage Allocation

```
1  typedef struct Node {
2    char data; struct Node* next; } Node;
3
4  Node *head = 0, *hlp;
5
6  void insert( char d ) {
7    hlp = (Node*)malloc( sizeof(Node) );
8    hlp->data = d;
9    hlp->next = head;
10   head = hlp;
11 }
12
13 insert( 'C' );
14 insert( 'B' );
15 insert( 'A' );
```

allocate some space for a
Node, return its address;
may fail ("out of memory"),
malloc(3) yields 0 then

A Linked List

```
1  typedef struct Node {
2    char data;
3    struct Node* next;
4  } Node;
5
6  Node c = { 'C', 0 };
7  Node b = { 'B', &c };
8  Node a = { 'A', &b };
```
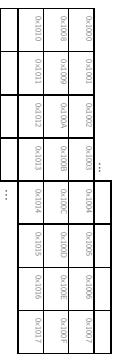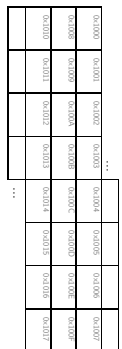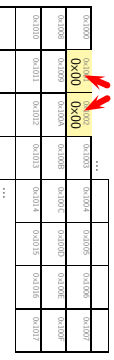

Dynamic Storage Allocation

```
1  typedef struct Node {
2    char data; struct Node* next; } Node;
3
4  Node *head = 0, *hlp;
5
6  void insert( char d ) {
7    hlp = (Node*)malloc( sizeof(Node) );
8    hlp->data = d;
9    hlp->next = head;
10   head = hlp;
11 }
12
13 insert( 'C' );
14 insert( 'B' );
15 insert( 'A' );
```

head

Dynamic Storage Allocation

```c
1  typedef struct Node {
2    char data; struct Node* next; } Node;
3
4  Node *head = 0, *hlp;
5
6  void insert( char d ) {
7    hlp = (Node*)malloc( sizeof(Node) );
8    hlp->data = d;
9    hlp->next = head;
10   head = hlp;
11 }
12
13 insert( 'C' );
14 insert( 'B' );
15 insert( 'A' );
```

Dynamic Storage Allocation

```c
1  typedef struct Node {
2    char data; struct Node* next; } Node;
3
4  Node *head = 0, *hlp;
5
6  void insert( char d ) {
7    hlp = (Node*)malloc( sizeof(Node) );
8    hlp->data = d;
9    hlp->next = head;
10   head = hlp;
11 }
12
13 insert( 'C' );
14 insert( 'B' );
15 insert( 'A' );
```

Dynamic Storage Allocation

```c
1  typedef struct Node {
2    char data; struct Node* next; } Node;
3
4  Node *head = 0, *hlp;
5
6  void insert( char d ) {
7    hlp = (Node*)malloc( sizeof(Node) );
8    hlp->data = d;
9    hlp->next = head;
10   head = hlp;
11 }
12
13 insert( 'C' );
14 insert( 'B' );
15 insert( 'A' );
```

Dynamic Storage Allocation

```c
1  typedef struct Node {
2    char data; struct Node* next; } Node;
3
4  Node *head = 0, *hlp;
5
6  void insert( char d ) {
7    hlp = (Node*)malloc( sizeof(Node) );
8    hlp->data = d;
9    hlp->next = head;
10   head = hlp;
11 }
12
13 insert( 'C' );
14 insert( 'B' );
15 insert( 'A' );
```

Dynamic Storage Allocation

```c
1  typedef struct Node {
2    char data; struct Node* next; } Node;
3
4  Node *head = 0, *hlp;
5
6  void insert( char d ) {
7    hlp = (Node*)malloc( sizeof(Node) );
8    hlp->data = d;
9    hlp->next = head;
10   head = hlp;
11 }
12
13 insert( 'C' );
14 insert( 'B' );
15 insert( 'A' );
```

Dynamic Storage Allocation

```c
1  typedef struct Node {
2    char data; struct Node* next; } Node;
3
4  Node *head = 0, *hlp;
5
6  void insert( char d ) {
7    hlp = (Node*)malloc( sizeof(Node) );
8    hlp->data = d;
9    hlp->next = head;
10   head = hlp;
11 }
12
13 insert( 'C' );
14 insert( 'B' );
15 insert( 'A' );
```

data: 'C'
next: 0x0000

Dynamic Storage Allocation

```
1  typedef struct Node {
2    char data; struct Node* next; } Node;
3
4  Node *head = 0, *hlp;
5
6  void insert( char d ) {
7    hlp = (Node*)malloc( sizeof(Node) );
8    hlp->data = d;
9    hlp->next = head;
10   head = hlp;
11 }
12
13 insert( 'C' );
14 insert( 'B' );
15 insert( 'A' );
```

data: 'C'
next: 0x0000

Dynamic Storage Allocation

```
1  typedef struct Node {
2    char data; struct Node* next; } Node;
3
4  Node *head = 0, *hlp;
5
6  void insert( char d ) {
7    hlp = (Node*)malloc( sizeof(Node) );
8    hlp->data = d;
9    hlp->next = head;
10   head = hlp;
11 }
12
13 insert( 'C' );
14 insert( 'B' );
15 insert( 'A' );
```

data: 'C'
next: 0x0000

Dynamic Storage Allocation

```
1  typedef struct Node {
2    char data; struct Node* next; } Node;
3
4  Node *head = 0, *hlp;
5
6  void insert( char d ) {
7    hlp = (Node*)malloc( sizeof(Node) );
8    hlp->data = d;
9    hlp->next = head;
10   head = hlp;
11 }
12
13 insert( 'C' );
14 insert( 'B' );
15 insert( 'A' );
```

data: 'C'
next: 0x0000

Dynamic Storage Allocation

```
1  typedef struct Node {
2    char data; struct Node* next; } Node;
3
4  Node *head = 0, *hlp;
5
6  void insert( char d ) {
7    hlp = (Node*)malloc( sizeof(Node) );
8    hlp->data = d;
9    hlp->next = head;
10   head = hlp;
11 }
12
13 insert( 'C' );
14 insert( 'B' );
15 insert( 'A' );
```

data: 'C'
next: 0x0000

Dynamic Storage Allocation

```
1  typedef struct Node {
2    char data; struct Node* next; } Node;
3
4  Node *head = 0, *hlp;
5
6  void insert( char d ) {
7    hlp = (Node*)malloc( sizeof(Node) );
8    hlp->data = d;
9    hlp->next = head;
10   head = hlp;
11 }
12
13 insert( 'C' );
14 insert( 'B' );
15 insert( 'A' );
```

data: 'B'
next: 0x1013

data: 'C'
next: 0x0000

Dynamic Storage Allocation

```
1  typedef struct Node {
2    char data; struct Node* next; } Node;
3
4  Node *head = 0, *hlp;
5
6  void insert( char d ) {
7    hlp = (Node*)malloc( sizeof(Node) );
8    hlp->data = d;
9    hlp->next = head;
10   head = hlp;
11 }
12
13 insert( 'C' );
14 insert( 'B' );
15 insert( 'A' );
```

data: 'C'
next: 0x0000

Dynamic Storage Allocation

```c
typedef struct Node {
   char data; struct Node* next; } Node;

Node *head = 0, *hlp;

void insert( char d ) {
   hlp = (Node*)malloc( sizeof(Node) );
   hlp->data = d;
   hlp->next = head;
   head = hlp;
}

insert( 'C' );
insert( 'B' );
insert( 'A' );
```

Dynamic Storage Management

Dynamic Storage Allocation:

- void* **malloc**(**size_t** size);
  - "[...] allocates **size** bytes and returns a pointer to the allocated memory. **The memory is not initialized.** [...]"
  - "On error, [this function] returns NULL."

- void **free**(void* ptr)
  - "[...] frees the memory space pointed to by **ptr**, which **must** have been returned by a previous call to malloc(). [...]"
  - "Otherwise, or if **free(ptr)** has already been called before, **undefined behavior** occurs."
  - "If **ptr** is **NULL**, no operation is performed."

- **No garbage collection!**
  Management of dynamic storage is **responsibility of the programmer**.
  Unaccessible, not free'd memory is called **memory leak**.

Dynamic Storage Management Example

```
1  void remove() {
2    if ( hlp = head ) {
3      head = hlp->next;
4      free(hlp);
5    }
6  }
7  insert( 'C' ); insert( 'B' ); insert( 'A' );
8  remove();
9  insert( 'X' );
```

Dynamic Storage Management Example

```
1 void remove() {
2   if (hlp = head) {
3     head = hlp->next;
4     free(hlp);
5   }
6 }
7 insert('C'); insert('B'); insert('A');
8 remove();
9 insert('X');
```

data: 'B'
next: 0x013

data: 'C'
next: 0x0000

Dynamic Storage Management Example

```
1 void remove() {
2   if (hlp = head) {
3     head = hlp->next;
4     free(hlp);
5   }
6 }
7 insert('C'); insert('B'); insert('A');
8 remove();
9 insert('X');
```

data: 'B'
next: 0x013

data: 'C'
next: 0x0000

Dynamic Storage Management Example

```
1 void remove() {
2   if (hlp = head) {
3     head = hlp->next;
4     free(hlp);
5   }
6 }
7 insert('C'); insert('B'); insert('A');
8 remove();
9 insert('X');
```

data: 'X'
next: 0x008

data: 'B'
next: 0x013

data: 'C'
next: 0x0000

Dynamic Storage Management Example

```
1 void remove() {
2   if (hlp = head) {
3     head = hlp->next;
4     free(hlp);
5   }
6 }
7 insert('C'); insert('B'); insert('A');
8 remove();
9 insert('X');
```

data: 'B'
next: 0x013

data: 'C'
next: 0x0000

Dynamic Linked List Iteration

```
1 Node* find(char d) {
2   hlp = head;
3   while (hlp) {
4     if (hlp->data == d)
5       break;
6     hlp = hlp->next;
7   }
8   return hlp;
9 }
10 insert('C'); insert('B'); insert('A');
11 find('B'); // yields 0x008
12 find('O'); // yields 0x0000, aka. NULL
```

Pointers to Struct/Union — '.' vs. '->'

```
1  typedef struct {
2    int x;
3    int y;
4  } coordinate;
5
6  coordinate pos = { 13, 27 };
7
8  coordinate* p = &pos;
9
10 int tmp;
11
12 tmp = (*p).x;
13 (*p).x = (*p).y;
14 (*p).y = tmp;
15
16 tmp = p->x;
17 p->x = p->y;
18 p->y = tmp;
```

Storage Duration of Objects

Storage Duration of Objects (6.2.4)

- **"static"** – e.g. variables in program scope:
  - live from program start to end
  - if not explicitly initialized, set to 0 (6.7.8)
- **"automatic"** – non-static variables in local scope:
  - live from block entry to exit
  - not automatically initialised: "initial value [...] is indeterminate"
- **"allocated"** – dynamic objects:
  - live from **malloc** to **free**
  - not automatically initialised

"If an object is referred to outside of its lifetime, **the behavior is undefined.**
The value of a pointer becomes indeterminate when the object it points to
reaches the end of its lifetime."

Example: Anatomy of a Linux Program in Memory

local variables
live here

malloc()/free()
work here

uninitialised
global variables,
set to 0, here

initialised global
variables here

program code
lives here

Storage Duration "Automatic" (Simplified)

```
1  void h() { int y; y++; }
2  void g() { int x = 5; x++; }
3  int* f() { int c = 3; g(); h(); return &c; }
4
5  int a = 27, b, *p;
6  p = f();
7  b = *p;
```

Storage Duration "Automatic" (Simplified)

```
1  void h() { int y; y++; }
2  void g() { int x = 5; x++; }
3  int* f() { int c = 3; g(); h(); return &c; }
4
5  int a = 27, b, *p;
6  p = f();
7  b = *p;
```


stack pointer – stack ends at
0x1012 in this case; stack grows
downwards (to smaller addr.)

Storage Duration "Automatic" (Simplified)

```
1  void h() { int y; y++; }
2  void g() { int x = 5; x++; }
3  int* f() { int c = 3; g(); h(); return &c; }
4
5  int a = 27, b, *p;
6  p = f();
7  b = *p;
```

Storage Duration "Automatic" (Simplified)

```
1  void h() { int y; y++; }
2  void g() { int x = 5; x++; }
3  int* f() { int c = 3; g(); h(); return &c; }
4
5  int a = 27, b, *p;
6  p = f();
7  b = *p;
```

c   p   b   a

Storage Duration "Automatic" (Simplified)

```
1  void h() { int y; y++; }
2  void g() { int x = 5; x++; }
3  int* f() { int c = 3; g(); h(); return &c; }
4
5  int a = 27, b, *p;
6  p = f();
7  b = *p;
```

x no longer alive!

p   b   a

Storage Duration "Automatic" (Simplified)

```
1  void h() { int y; y++; }
2  void g() { int x = 5; x++; }
3  int* f() { int c = 3; g(); h(); return &c; }
4
5  int a = 27, b, *p;
6  p = f();
7  b = *p;
```

p   b   a

x

Storage Duration "Automatic" (Simplified)

```
1  void h() { int y; y++; }
2  void g() { int x = 5; x++; }
3  int* f() { int c = 3; g(); h(); return &c; }
4
5  int a = 27, b, *p;
6  p = f();
7  b = *p;
```

(now) y – not explicitly
initialised, thus initial
value is indeterminate

p   b   a

Storage Duration "Automatic" (Simplified)

```
1  void h() { int y; y++; }
2  void g() { int x = 5; x++; }
3  int* f() { int c = 3; g(); h(); return &c; }
4
5  int a = 27, b, *p;
6  p = f();
7  b = *p;
```

p   b   a

Storage Duration "Automatic" (Simplified)

```
1  void h() { int y; y++; }
2  void g() { int x = 5; x++; }
3  int* f() { int c = 3; g(); h(); return &c; }
4
5  int a = 27, b, *p;
6  p = f();
7  b = *p;
```

p   b   a

Storage Duration "Automatic" (Simplified)

```
1  void h() { int y; y++; }
2  void g() { int x = 5; x++; }
3  int* f() { int c = 3; g(); h(); return &c; }
4
5  int a = 27, b, *p;
6  p = f();
7  b = *p;
```

Storage Duration "Automatic" (Simplified)

```
1  void h() { int y; y++; }
2  void g() { int x = 5; x++; }
3  int* f() { int c = 3; g(); h(); return &c; }
4
5  int a = 27, b, *p;
6  p = f();
7  b = *p;
```

y no longer alive!

Storage Duration "Automatic" (Simplified)

```
1  void h() { int y; y++; }
2  void g() { int x = 5; x++; }
3  int* f() { int c = 3; g(); h(); return &c; }
4
5  int a = 27, b, *p;
6  p = f();
7  b = *p;
```

Storage Duration "Automatic" (Simplified)

```
1  void h() { int y; y++; }
2  void g() { int x = 5; x++; }
3  int* f() { int c = 3; g(); h(); return &c; }
4
5  int a = 27, b, *p;
6  p = f();
7  b = *p;
```

Storage Duration "Automatic" (Simplified)

```
1  void h() { int y; y++; }
2  void g() { int x = 5; x++; }
3  int* f() { int c = 3; g(); h(); return &c; }
4
5  int a = 27, b, *p;
6  p = f();
7  b = *p;
```

Storage Duration "Automatic" (Simplified)

```
1  void h() { int y; y++; }
2  void g() { int x = 5; x++; }
3  int* f() { int c = 3; g(); h(); return &c; }
4
5  int a = 27, b, *p;
6  p = f();
7  b = *p;
```

Storage Duration "Automatic" (Simplified)

```
1  void h() { int y; y++; }
2  void g() { int x = 5; x++; }
3  int* f() { int c = 3; g(); h(); return &c; }
4
5  int a = 27, b, *p;
6  p = f();
7  b = *p;
```

p refers to a non-alive object, the behavior is undefined (everything **may happen**, from 'crash' to 'ignore').

Storage Classes and Qualifiers

Storage Class Specifiers (6.7.1)

Storage Class Specifiers (6.7.1)

```
1  typedef char letter;
2
3  extern int x;
4  extern int f();
5
6  static int x; // two uses! (-> later)
7  static int f();
8
9
10 auto x; // "historic"
11
12 register y; // "historic"
```

Storage Class Specifiers: extern (6.7.1)

```
1  // not _defined_ here, "imported" ...
2  //
3  extern int x;
4  extern void f();
5
6  // declared _and_ defined here, "exported" ...
7  //
8  int y;
9
10 int g() {
11   x = y = 27;
12   f();
13 }
```

- → modules, linking (later)
- usually only extern in headers (later)

Storage Class Specifiers: static (6.7.1)

```
1  // declared _and_ defined here,
2  // _not_ "exported" ...
3
4  static int x;
5  static void g();
6
7  int f() {
8    static int a = 0;
9    a++;
10   printf( "%s\n", a );
11 }
12
13
14 f(); f(); f(); // yields 1, 2, 3
```

Excursion: Memory Mapped I/O

- **Intuition:** some memory addresses are wired to hardware
- **writing** to the address causes a pin to change logical value
- **reading** the address gives logical value of a pin

Qualifiers (6.7.3)

```
1   int x;
2
3   const int y;
4
5   volatile int z;
6
7   int* restrict p; // aliasing
8
9
10  const volatile int a;
```

Qualifiers (6.7.3)

```
1   int x;
2
3   const int y;
4
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```

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```
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2
3   const int y;
4
5   volatile int z;
6
7   int* restrict p; // aliasing
8
9
10  const volatile int a;
```

- **restrict:**
- "[... lengthy formal definition ...]"
- "[...] If these requirements are not met, then the behavior is **undefined.**"
- → use **extremely carefully** (i.e. if in doubt, not at all)

Strings & Input/Output

Excursion: Memory Mapped I/O

- **Intuition**: some memory addresses are wired to hardware
- **writing** to the address causes a pin to change logical value
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Strings

Excursion: Memory Mapped I/O

- **Intuition**: some memory addresses are wired to hardware
- **writing** to the address causes a pin to change logical value
- **reading** the address gives logical value of a pin

- The compiler does not know, "**memory is memory**".

Strings are 0-Terminated **char** Arrays

```
1  char* msg = "Hello";
2  char* str = msg;
```

msg

(char) 0

Qualifiers: $volatile$ (6.7.3)

```
1  volatile char* out = 0x1006;
2  volatile char*  in = 0x1007;
3
4  out = 0x01; // switch lamp on
5
6  if (in & 0x01) { /* .... */ }
7
8  if ((in & 0x01) && (in & 0x01)) { /* .... */ }
```

Strings are 0-Terminated *char* Arrays

```
1  char* msg = "Hello";
2  char* str = msg;
```

msg

(char) 0

str

String Manipulation (Annex B)

include<string.h>

provides among others:

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  "[...] calculates length of string s, excluding **the terminating null byte** ('\0')."

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- int **strcmp**(const char* s1, const char* s2)
  "[...] compares the two strings s1 and s2.
  It returns an integer less than, equal to, or greater than zero if s1 is found,
  respectively, to be less than, to match, or be greater than s2."

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- char* **strcpy**(char* s1, const char* s2)
  "The strcpy() function copies the string pointed to by s2,
  including **the terminating null byte** ('\0'), to the buffer pointed to by s1."

String Manipulation (Annex B)

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- char* **strcpy**(char* s1, const char* s2)
  "The strcpy() function copies the string pointed to by s2,
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- char* **strncpy**(char* s1, const char* s2, size_t n)

None of these functions allocates memory!

String Manipulation (Annex B)

include<string.h>

provides among others:

- **size_t strlen(const char* s)**
"[...] calculates length of string s, excluding **the terminating null byte** ('\0')."

- int **strcmp**(const char* s1, const char* s2)
"[...] compares the two strings s1 and s2.
It returns an integer less than, equal to, or greater than zero if s1 is found, respectively, to be less than, to match, or be greater than s2."

- char* **strcpy**(char* s1, const char* s2)
"The strcpy() function copies the string pointed to by s2, including **the terminating null byte** ('\0'), to the buffer pointed to by s1."

- char* **strncpy**(char* s1, const char* s2, size_t n)

None of these functions allocates memory!

Allocate and copy: (not C99, but POSIX)

- char* **strdup**(const char* s)

Input/Ouput

Printing

```
1  #include <stdio.h>
2
3  printf( "%s %i %f\n", "Hello", 27, 3.14 );
```

Tools & Modules

Hello, Again

```
1  #include <stdio.h>
2
3  int g( int x ) { return x/2; }
4
5  int f() { return g(1); }
6
7  int main() {
8    printf( "Hello World.\n" );
9    return f();
10 }
```

- % gcc helloworld.c
- % ls
- a.out helloworld.c
- % ./a.out
- Hello World.
- %

Zoom In: Preprocessing, Compiling, Linking

```
1  #include <stdio.h>
2
3  int g( int x ) { return x/2; }
4
5  int f() { return g(1); }
6
7  int main() {
8    printf( "Hello World.\n" );
9    return f();
10 }
```

- % gcc -E helloworld.c > helloworld.i
- % gcc -c -o helloworld.i
- % ld -o helloworld [...] helloworld.o [...]
- % ./helloworld
- Hello World.
- %

Modules

```
1  #include <stdio.h>
2
3  int g( int x ) {
4    return x/2;
5  }
6
7  int f() {
8    return g(1);
9  }
10
11 int main() {
12   printf( "Hello World \n" );
13   return f();
14 }
```

Modules

Split into:

- .h (header): declarations
- .c: definitions, use headers to "import" declarations

g.h
```
1  #ifndef G_H
2  #define G_H
3
4  extern int
5  g( int x );
6  #endif
```

g.c
```
1  #include "g.h"
2
3  int g( int x ) {
4    return x/2;
5  }
```

f.h
```
1  #ifndef F_H
2  #define F_H
3
4  extern int
5  f();
6  #endif
```

f.c
```
1  #include "g.h"
2  #include "f.h"
3
4  int f() {
5    return g(1);
6  }
```

helloworld.c
```
1  #include <stdio.h>
2  #include "f.h"
3
4  int main() {
5    printf( "Hello World \n" );
6    return f();
7  }
```

Modules

Split into:

- .h (header): declarations
- .c: definitions, use headers to "import" declarations

g.h
```
1  #ifndef G_H
2  #define G_H
3
4  extern int
5  g( int x );
6  #endif
```

g.c
```
1  #include "g.h"
2
3  int g( int x ) {
4    return x/2;
5  }
```

f.h
```
1  #ifndef F_H
2  #define F_H
3
4  extern int
5  f();
6  #endif
```

f.c
```
1  #include "g.h"
2  #include "f.h"
3
4  int f() {
5    return g(1);
6  }
```

helloworld.c
```
1  #include <stdio.h>
2  #include "f.h"
3
4  int main() {
5    printf( "Hello World \n" );
6    return f();
7  }
```

Modules At Work

preprocess & compile:
- % gcc -c g.c f.c \
    helloworld.c
- % ls *.o
  f.o g.o helloworld.o

link:
- % gcc g.o f.o helloworld.o

execute:
- % ./a.out
  Hello World.

g.h
```
1  #ifndef G_H
2  #define G_H
3
4  extern int
5  g( int x );
6  #endif
```

g.c
```
1  #include "g.h"
2
3  int g( int x ) {
4    return x/2;
5  }
```

f.h
```
1  #ifndef F_H
2  #define F_H
3
4  extern int
5  f();
6  #endif
```

f.c
```
1  #include "g.h"
2  #include "f.h"
3
4  int f() {
5    return g(1);
6  }
```

helloworld.c
```
1  #include <stdio.h>
2  #include "f.h"
3
4  int main() {
5    printf( "Hello World \n" );
6    return f();
7  }
```

Modules At Work

preprocess & compile:
- % gcc -c g.c f.c \
    helloworld.c
- % ls *.o

link:
- % gcc g.o f.o helloworld.o

execute:
- % ./a.out
  Hello World.

g.h
```
1  #ifndef G_H
2  #define G_H
3
4  extern int
5  g( int x );
6  #endif
```

g.c
```
1  #include "g.h"
2
3  int g( int x ) {
4    return x/2;
5  }
```

f.h
```
1  #ifndef F_H
2  #define F_H
3
4  extern int
5  f();
6  #endif
```

f.c
```
1  #include "g.h"
2  #include "f.h"
3
4  int f() {
5    return g(1);
6  }
```

helloworld.c
```
1  #include <stdio.h>
2  #include "f.h"
3
4  int main() {
5    printf( "Hi!\n" );
6    return f();
7  }
```

Modules At Work

preprocess & compile:
- % gcc -c g.c f.c \
    helloworld.c
- % ls *.o

link:
- % gcc g.o f.o helloworld.o

execute:
- % ./a.out
  Hello World.

fix and re-build:
- % gcc -c helloworld.c
- % gcc g.o f.o helloworld.o
- % ./a.out
  Hi!

g.h
```
1  #ifndef G_H
2  #define G_H
3
4  extern int
5  g( int x );
6  #endif
```

g.c
```
1  #include "g.h"
2
3  int g( int x ) {
4    return x/2;
5  }
```

f.h
```
1  #ifndef F_H
2  #define F_H
3
4  extern int
5  f();
6  #endif
```

f.c
```
1  #include "g.h"
2  #include "f.h"
3
4  int f() {
5    return g(1);
6  }
```

helloworld.c
```
1  #include <stdio.h>
2  #include "f.h"
3
4  int main() {
5    printf( "Hi!\n" );
6    return f();
7  }
```

Preprocessing

helloworld.c
```
#include <stdio.h>
#include "..."
int main() {
  printf( "Hello-World\n" );
  return f();
}
```

- % gcc -E helloworld.c
  -o helloworld.i

helloworld.i
```
1  # 1 "helloworld.c"
2  # 1 "<command-line>"
3  # 1 "helloworld.c"
4  # 1 "/usr/include/stdio.h" 1 3 4
5  # 28 "/usr/include/stdio.h" 3 4
...
```

Preprocessing Directives (6.10)

```
1   #include <stdio.h>
2   #include "battery.h"
3
4   #define PI 3.1415
5
6   #ifdef DEBUG
7     fprintf( stderr, "honk\n" );
8   #ifdef
9   #endif
10
11  #if __GNUC__ - >= 3
12  # define -pure    -attribute- -((pure))
13  #else
14  # define -pure    /* no pure */
15  #endif
16
17  extern int f() -pure;
```

Linking

provides: int g(int)
needs: -/-   g.o

provides: int f()
needs: int g(int)   f.o

provides: int main()
int f(int)
int printf(const char*,...)
needs: helloworld.o

provides: int printf(const char*,...)
needs: ...   libc.a

Linking

provides: int g(int)
needs: -/-   g.o

provides: int f()
needs: int g(int)   f.o

provides: int main()
int f(int)
int printf(const char*,...)
needs: helloworld.o

provides: int printf(const char*,...)
needs: ...   libc.a

Linking

provides: int g(int)
needs: -/-   g.o

provides: int f()
needs: int g(int)   f.o

provides: int main()
int f(int)
int printf(const char*,...)
needs: helloworld.o

provides: int printf(const char*,...)
needs: ...   libc.a

Linking

provides: int g(int)
needs: -/-   g.o

provides: int f()
needs: int g(int)   f.o

provides: int main()
int f(int)
int printf(const char*,...)
needs: helloworld.o

provides: int printf(const char*,...)
needs: ...   libc.a

Linking

provides: int g(int)
needs: ./. → **g.o**

provides: int f()
needs: int g(int) → **f.o**

provides: int g(int)
needs: int printf(const char*,...) → **helloworld.o**

provides: int main()
needs: int f(int)
int printf(const char*,...)

provides: int printf(const char*,...)
needs: ... → **libc.a**

a.out

Compiler

gcc [OPTION]... infile...

-E – preprocess only

-c – compile only, don't link

-o outfile – write output to **outfile**

  Example: gcc -c main.c — produces main.o

  Example: gcc -c -o x.o main.c — produces x.o

-g – add debug information

-W, **-Wall**, ... – enable warnings

-I dir – add **dir** to **include path** for searching headers

-L dir – add **dir** to **library path** for searching libraries

-D macro[=defn] – define **macro** (to **defn**)

  Example: gcc -DDEBUG -DMAGICNUMBER=27

-l library link against liblibrary.{a,so}, order matters

  Example: gcc a.o b.o main.o -lxy

→ cf. man gcc

---

gdb(1), ddd(1), nm(1), make(1)

- **Command Line Debugger:**

  gdb a.out [core]

- **GUI Debugger:**

  ddd a.out [core]

  (works best with debugging information compiled in (gcc -g))

- **Inspect Object Files:**

  nm a.o

- **Build Utility:**

  make

See battery controller exercise for an example.

Core Dumps

- **Recall:** Anatomy of a Linux Program in Memory
- **Core dump:** (basically) this memory written to a file.

Core Dumps

- **Recall:** Anatomy of a Linux Program in Memory
- **Core dump:** (basically) this memory written to a file.

```
1  int main() {
2    int *p;
3    *p = 27;
4    return 0;
5  }
```

```
1  % gcc -g core.c
2  % limit coredumpsize
3  coredumpsize 0 kbytes
4  % limit coredumpsize 1g
5  % ./a.out
6  Segmentation fault (core dumped)
7  % ls -lh core
8  -rw------- 1 user user 232K Feb 29 11:11 core
9  % gdb a.out core
10 GNU gdb (GDB) 7.4.1-debian
11 [...]
12 Core was generated by `./a.out'.
13 Program terminated with signal 11, Segmentation fault.
14 #0 0x000000000040046d in main () at core.c:3
15 3    *p = 27;
16 (gdb) p p
17 $1 = (int *) 0x0
18 (gdb) p *p
```

Formal Methods for C

Correctness and Requirements

Correctness

- Correctness is defined **with respect to** a specification.
- A program (function, ...) is **correct** (wrt. specification φ) **if and only if** it satisfies φ.
- Definition of "satisfies": **in a minute.**

Examples:

- φ_1: the return value is 10 divided by parameter (if parameter not 0)
- φ_2: the value of variable x is "always" strictly greater than 3
- φ_3: the value of i increases in each loop iteration
- ...

Common Patterns

- **State Invariants:**
  "at **this** program point, the value of p must not be NULL"
  "at **all** program points, the value of p must not be NULL"
  (cf. **sequence points** (Annex C))

- **Data Invariants:**
  "the value of n must be the length of s"

- **(Function) Pre/Post Conditions:**
  Pre-Condition: the parameter must not be 0
  Post-Condition: the return value is 10 divided by the parameter

- **Loop Invariants:**
  "the value of i is between 0 and array length minus 1"

Poor Man's Requirements Specification
aka. How to Formalize Requirements in C?

Diagnostics (7.2)

```
1  #include <assert.h>
2  void assert( /* scalar */ expression );
```

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```
1  #include <assert.h>
2  void assert( /* scalar */ expression );
```

- "The assert macro puts diagnostic tests into programs; [...]
  When it is executed, if expression (which shall have a scalar type) is false (that is, compares equal to 0), the assert macro

- writes information about the particular call that failed [...] on the standard error stream in an implementation-defined format.

- It then calls the **abort** function."

Diagnostics (7.2)

```
1  #include <assert.h>
2  void assert( /* scalar */ expression );
```

- "The assert macro puts diagnostic tests into programs; [...]
  When it is executed, if expression (which shall have a scalar type) is false (that is, compares equal to 0), the assert macro writes information about the particular call that failed [...] on the standard error stream in an implementation-defined format.
  It then calls the abort function."

- Pitfall:

- If macro NDEBUG is defined when including <assert.h>, expression is not evaluated (thus should be side-effect free).

---

abort (7.20.4.1)

```
1  #include <stdlib.h>
2
3  void abort();
```

- "The abort function causes abnormal program termination to occur, unless [...]

- [...] An implementation-defined form of the status unsuccessful termination is returned to the host environment by means of the function call raise(SIGABRT)."

(→ Core Dumps)

Common Patterns with assert

- **State Invariants:**
  "at **this** program point, the value of p must not be NULL"
  "at **all** program points, the value of p must not be NULL"
  (cf. **sequence points** (Annex C))

- **Data Invariants:**
  "the value of n must be the length of s"

- **(Function) Pre/Post Conditions:**
  Pre-Condition: the parameter must not be 0
  Post-Condition: the return value is 10 divided by the parameter

- **Loop Invariants:**
  "the value of i is between 0 and array length minus 1"

State Invariants with <assert.h>

```
1  void f() {
2    int* p = (int*)malloc(sizeof(int));
3
4    if (!p)
5      return;
6
7    assert(p); // assume p is valid from here
8
9  } // ...
10
11 void g() {
12   Node* p = find( 'a' );
13
14   assert(p); // we inserted 'a' before
15 } // ...
16 }
```

Data Invariants with <assert.h>

```
1  typedef struct {
2    char* s;
3    int n;
4  } str;
5
6  str* construct( char* s ) {
7    str* x = (str*)malloc( sizeof(str) );
8    // ...
9    assert( (x->s == NULL && x->n == -1)
10          || (x->n = strlen( x->s ) ) );
11 }
```

Pre/Post Conditions with <assert.h>

```
1  int f( int x ) {
2    assert( x != 0 ); // pre-condition
3
4    int r = 10/x;
5
6    assert( r == 10/x ); // post-condition
7
8    return r;
9  }
```

Loop Invariants with <*assert.h*>

```
1   void f( int a[], int n ) {
2     int i = 0;
3
4     // holds before the loop
5     assert( 0 <= i && i <= n );
6     assert( i < 1 || a[i-1] == 0 );
7
8     while (i < n) {
9       // holds before each iteration
10      assert( 0 <= i && i <= n );
11      assert( i < 1 || a[i-1] == 0 );
12
13      a[i++] = 0;
14    }
15    // holds after exiting the loop
16    assert( 0 <= i && i <= n );
17    assert( i < 1 || a[i-1] == 0 );
18
19
20    return;
   }
```

Old Variables, Ghost Variables

```
1   void xorSwap( unsigned int* a, unsigned int* b ) {
2   #ifndef NDEBUG
3     unsigned int *old_a = a, *old_b = b;
4   #endif
5     assert( a && b ); assert( a != b ); // pre-condition
6
7     *a = *a + *b;
8     *b = *a - *b;
9     *a = *a - *b;
10
11    assert( *a == *old_b && *b == *old_a );  // post-con-
12    assert( a == old_a && b == old_b );      // dition
13  }
```

Outlook

- Some verification tools simply verify for each assert statement:
  When executed, expression is not false.

- Some verification tools support sophisticated requirements specification languages like ACSL with explicit support for
  - pre/post conditions
  - ghost variables, old values
  - data invariants
  - loop invariants
  - . . .

Dependable Verification (Jackson)

Dependability

- **"The program has been verified."** tells us

Dependability

- **"The program has been verified."** tells us **not very much.**

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- **Which specifications** have been considered?
- Under **which assumptions** was the verification conducted?
  - Platform assumptions: finite words (size?), mathematical integers, . . .
  - Environment assumptions, input values, . . .

  Assumptions are often implicit, **"in the tool"**!

Dependability

- **"The program has been verified."** tells us **not very much.**
- One wants to know (and should state):
- **Which specifications** have been considered?
- Under **which assumptions** was the verification conducted?
  - Platform assumptions: finite words (size?), mathematical integers, . . .
  - Environment assumptions, input values, . . .

  Assumptions are often implicit, **"in the tool"**!
- And **what does verification mean** after all?
- In some contexts: **testing.**
- In some contexts: **review.**
- In some contexts: **model-checking** procedure.
- ("We verified the program!" – "What did the tool say?" – "Verification failed.")
- In some contexts: **model-checking tool claims correctness.**

Common Errors

Distinguish

Most **generic errors** boil down to:

- specified but **unwanted behaviour,**
  e.g. under/overflows
- **initialisation issues**
  e.g. automatic block scope objects
- **unspecified behaviour** (J.1)
  e.g. order of evaluation in some cases
- **undefined behaviour** (J.2)
- **implementation defined behaviour** (J.3)

Conformance (4)

- "A program that is
  - correct in all other aspects,
  - operating on correct data,
  - containing **unspecified behavior**

  shall be a correct program and act in accordance with 5.1.2.3. (Program Execution)
- A conforming program is one that is acceptable to a conforming implementation.
- Strictly conforming programs are intended to be maximally portable among conforming implementations.
- An implementation [of C, a compiler] shall be accompanied by a document that defines all implementation-defined and locale-specific characteristics and all extensions.

Over- and Underflows

Over- and Underflows, Casting

- Not specific to C...

```
1   void f( short a, int b ) {
2     a = b; // typing ok, but ...
3   }
4
5   short a; // provisioning, implicit cast
6   if (++a < 0) { /* no */ }
7
8   if (++i > MAX_INT) {
9     /* no */ }
10
11
12  int e = 0;
13
14  void set_error() { e++; }
15  void clear_error() { e = 0; }
16
17  void g() { if (e) { /* ... */ } }
```

Initialisation (6.7.8)

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- "If an object that has automatic storage duration is not initialized explicitly, its value is indeterminate."

```
1   void f() {
2     int a;
3
4     printf( "%i\n", a ); // surprise ...
5   }
```

Unspecified Behaviour (J.1)

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Each implementation (of a compiler) documents how the choice is made.

For example

- whether two string literals result in distinct arrays (6.4.5)
- the order in which the function designator, arguments, and subexpressions within the arguments are evaluated in a function call (6.5.2.2)
- the layout of storage for function parameters (6.9.1)
- the result of rounding when the value is out of range (7.12.9.5, ...)
- the order and contiguity of storage allocated by successive calls to malloc (7.20.3)
- etc. pp.

```
1  char a[] = "hello", b[] = "hello"; // a == b?
2
3  i = 0; f( ++i, ++i, ++i ); // f(1,2,3)?
4
5  int g() { int a, b; } // &a > &b ?
6
7  int* p = malloc(sizeof(int));
8  int* q = malloc(sizeof(int)); // q > p?
```

Undefined Behaviour (J.2)

Undefined Behaviour (3.4.3)

"Behaviour, upon use of a non-portable or erroneous program construct or of erroneous data, for which this International Standard imposes no requirements."

"**Possible undefined behaviour ranges from**

- ignoring the situation completely with **unpredictable results,**
- to behaving during **translation or program execution** in a documented manner characteristic of the environment (with or without the issuance of a diagnostic message),
- to terminating a **translation or execution** (with the issuance of a diagnostic message)."

"An example of undefined behaviour is the behaviour on **integer overflow.**"

Undefined Behaviour (J.2)

More examples:

- an identifier [...] contains an invalid multibyte character (5.2.1.2)
- an object is referred to outside of its lifetime (6.2.4)

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Null-Pointer

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1  int main() {
2    int *p;
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- the program removes the definition of a macro whose name begins with an underscore and either an uppercase letter or another underscore (7.1.3)

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- "An integer constant expression with the value 0, or such an expression cast to type 'void*, is called a **null pointer constant**. [..]"
- "The macro **NULL** is defined in <**stddef.h**> (and other headers) as a null pointer constant; see 7.17."
- "Among the invalid values for dereferencing a pointer by the unary * operator are a null pointer, [..]" (6.5.3.2)

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Segmentation Violation

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1  int main() {
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- Modern operating systems provide **memory protection**.
- Accessing memory which the process is not allowed to access is observed by the operating system.
- Typically an instance of "accessing an object outside its lifetime".
- **But:** other way round does not hold, accessing an object outside its lifetime does not imply a segmentation violation.

Implementation-Defined Behaviour (J.3)

"A conforming implementation is required to document its choice of behavior in each of the areas listed in this subclause. The following are implementation-defined:"

- J.3.2 Environment, e.g.
  The set of signals, their semantics, and their default handling (7.14).

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- **But:** other way round does not hold, accessing an object outside its lifetime does not imply a segmentation violation.

- Some platforms (e.g. SPARC): unaligned memory access, i.e. outside word boundaries, not supported by hardware ("bus error").
  Operating system notifies process, default handler: terminate, dump core.

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- J.3.6 Floating Point, e.g.
  The accuracy of the floating-point operations [...] (5.2.4.2.2).
- J.3.7 Arrays and Pointers, e.g.
  The result of converting a pointer to an integer or vice versa (6.3.2.3).
- etc. pp.

Locale and Common Extensions (J.4, J.5)

- J.4 Locale-specific behaviour
- J.5 Common extensions
  "The following extensions are widely used in many systems, but are not portable to all implementations."

[ISO, 1999] ISO (1999). Programming languages – C. Technical Report
ISO/IEC 9899:1999, ISO. Second edition, 1999-12-01.