Real-Time Systems

Lecture 16: The Universality Problem for TBA

2014-07-29

Dr. Bernd Westphal

Albert-Ludwigs-Universität Freiburg, Germany

Contents & Goals

Last Lecture:

- Extended Timed Automata Cont'd
- A Fragment of TCTL
- Testable DC Formulae

This Lecture:

- **Educational Objectives:** Capabilities for following tasks/questions.
  - Are all DC formulae testable?
  - What's a TBA and what's the difference to (extended) TA?
  - What's undecidable for timed (Büchi) automata? Idea of the proof?
- **Content:**
  - An untestable DC formula.
  - Timed Büchi Automata and timed regular languages [Alur and Dill, 1994].
  - The Universality Problem is undecidable for TBA [Alur and Dill, 1994]
  - Why this is unfortunate.
  - Timed regular languages are not everything.

Untestable DC Formulae

Recall: Testability

Definition 6.1. A DC formula F is called **testable** if an observer (or test automaton (or monitor)) $\mathcal{A}_F$ exists such that for all networks $\mathcal{N} = \mathcal{C}(\mathcal{A}_1, \dots, \mathcal{A}_n)$ it holds that

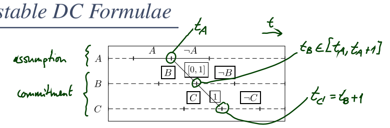
$$\mathcal{N} \models F \quad \text{iff} \quad \mathcal{C}(\mathcal{A}'_1, \dots, \mathcal{A}'_n, \mathcal{A}_F) \models \forall\Box\neg(\mathcal{A}_F.q_{bad})$$

Otherwise it's called **untestable**.

Proposition 6.3. There exist untestable DC formulae.

Theorem 6.4. DC implementables are testable.

Untestable DC Formulae



"Whenever we observe a change from A to $\neg A$ at time t_A,
the system has to produce a change from B to $\neg B$ at some time $t_B \in [t_A, t_A + 1]$
and a change from C to $\neg C$ at time $t_B + 1$.

Sketch of Proof: Assume there is $\mathcal{A}_F$ such that, for all networks $\mathcal{N}$, we have

$$\mathcal{N} \models F \quad \text{iff} \quad \mathcal{C}(\mathcal{A}'_1, \dots, \mathcal{A}'_n, \mathcal{A}_F) \models \forall\Box\neg(\mathcal{A}_F.q_{bad})$$

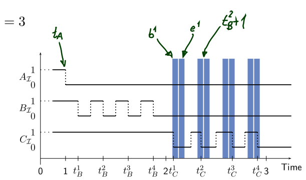
Assume the number of clocks in $\mathcal{A}_F$ is $n \in \mathbb{N}_0$.

Untestable DC Formulae Cont'd

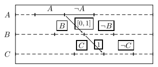
Consider the following time points:

- $t_A := 1$
- $t_B^i := t_A + \frac{2i-1}{2(n+1)}$ for $i = 1, \dots, n+1$
- $t_C^i \in \left]\underbrace{t_B^i + 1 - \frac{1}{4(n+1)}}_{=b^i}, \underbrace{t_B^i + 1 + \frac{1}{4(n+1)}}_{=e^i}\right[$ for $i = 1, \dots, n+1$
  with $t_C^i - t_B^i \neq 1$ for $1 \leq i \leq n+1$.

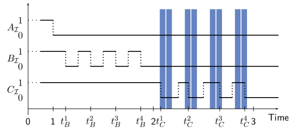
Example: $n = 3$



Untestable DC Formulae Cont'd



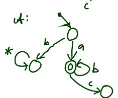
Example: $n = 3$



- The shown interpretation $\mathcal{I}$ satisfies **assumption** of property.
- It has $n + 1$ candidates to satisfy **commitment**.
- By choice of t_C^i, the commitment is not satisfied; so F not satisfied.
- Because $\mathcal{A}_F$ is a test automaton for F, is has a computation path to q_{bad}.
- Because $n = 3$, $\mathcal{A}_F$ can not save all $n + 1$ time points t_B^i.
- Thus there is $1 \leq i_0 \leq n$ such that all clocks of $\mathcal{A}_F$ have a valuation which is not in $2 - t_B^{i_0} + (-\frac{1}{4(n+1)}, \frac{1}{4(n+1)})$

$\Sigma = \{a, b, c\}$

$\mathcal{A}$:



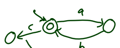
$\mathcal{A} = \{Q, q_0, \to, F\}$, $F \subseteq Q$

$L = ab^*$

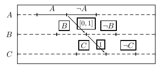
$L(\mathcal{A}) = \{\sigma_0\sigma_1 \ldots \in \Sigma^* \mid q_0 \xrightarrow{\sigma_0} q_1 \cdots \xrightarrow{\sigma_n} q_n,\ q_n \in F\}$

baab $\in^?$ $L(\mathcal{A})$
abb $\in^?$ $L(\mathcal{A})$
aba $\in^?$ $L(\mathcal{A})$
ϵ $\in^?$ $L(\mathcal{A})$
abc $\in^?$ $L(\mathcal{A})$

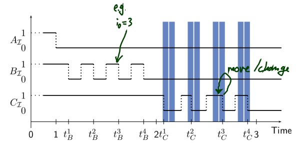
$\mathcal{B}$: $(ab)^*$



Untestable DC Formulae Cont'd



Example: $n = 3$



- Because $\mathcal{A}_F$ is a test automaton for F, is has a computation path to q_{bad}.
- Thus there is $1 \leq i_0 \leq n$ such that all clocks of $\mathcal{A}_F$ have a valuation which is not in $2 - t_B^{i_0} + (-\frac{1}{4(n+1)}, \frac{1}{4(n+1)})$
- Modify the computation to $\mathcal{I}'$ such that $t_C^{i_0} := t_B^{i_0} + 1$.
- Then $\mathcal{I}' \models F$, but $\mathcal{A}_F$ reaches q_{bad} via the same path.
- That is: $\mathcal{A}_F$ claims $\mathcal{I}' \not\models F$.
- Thus $\mathcal{A}_F$ is not a test automaton. **Contradiction**.

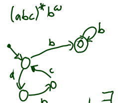
$\Sigma = \{a, b, c\}$

Σ^ω

$\mathcal{A} = (Q, q_{ini}, \to, F)$, $F \subseteq Q$

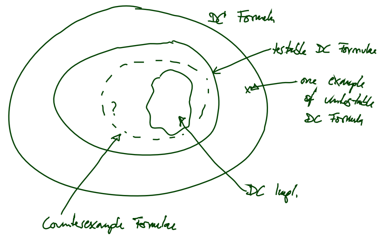
$(abc)^* b^\omega$

$\mathcal{A}$



$L(\mathcal{A}) = \{\sigma_0\sigma_1 \ldots \in \Sigma^\omega \mid \exists\, q_0 \xrightarrow{\sigma_0} q_1 \xrightarrow{\sigma_1} q_2 \ldots$ and $q_i \in F$ for infinitely many $i\}$

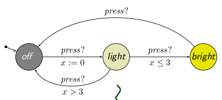
$\mathcal{B}$: $(abc)^\omega$



Timed Büchi Automata

[Alur and Dill, 1994]

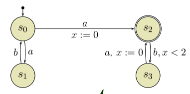
... vs. Timed Automata



produces

$\xi = \langle off, 0\rangle, 0 \xrightarrow{1} \langle off, 1\rangle, 1$
$\xrightarrow{press?} \langle light, 0\rangle, 1 \xrightarrow{3} \langle light, 3\rangle, 4$
$\xrightarrow{press?} \langle bright, 3\rangle, 4 \xrightarrow{\ } \ldots$

ξ is a **computation path** and **run** of $\mathcal{A}$.



accept

New: Given a **timed word**

$(a, 1), (b, 2), (a, 3), (b, 4), (a, 5), (b, 6), (b, 6.5)$ ...

does $\mathcal{A}$ **accept** it?

New: acceptance criterion is
visiting accepting state infinitely often.

Timed Languages

Definition. A **time sequence** $\tau = \tau_1, \tau_2, \ldots$ is an infinite sequence of time values $\tau_i \in \mathbb{R}_0^+$, satisfying the following constraints:

(i) **Monotonicity**:
τ increases **strictly** monotonically, i.e. $\tau_i < \tau_{i+1}$ for all $i \ge 1$.

(ii) **Progress**: For every $t \in \mathbb{R}_0^+$, there is some $i \ge 1$ such that $\tau_i > t$.

Definition. A **timed word** over an alphabet Σ is a pair (σ, τ) where

- $\sigma = \sigma_1, \sigma_2, \cdots \in \Sigma^\omega$ is an infinite word over Σ, and
- τ is a time sequence.

Definition. A **timed language** over an alphabet Σ is a set of timed words over Σ.

Example: Timed Language

Timed word over alphabet Σ: a pair (σ, τ) where

- $\sigma = \sigma_1, \sigma_2, \ldots$ is an infinite word over Σ, and
- τ is a time sequence (strictly (!) monotonic, non-Zeno).

z.B. a: flash light
b: beep

$$L_{crt} = \{((ab)^\omega, \tau) \mid \exists\, i\ \forall\, j \ge i : (\tau_{2j} < \tau_{2j-1} + 2)\}$$

a b a b a b ... a b a b ...

τ_1 τ_2 / 1,0 100

$j = 0$ $\tau_{2j} = \tau_2 = 10.0$

$\tau_{2j-1} + 2 = 1.0 + 2 = 3.0$

τ_{i-1} τ_i

$1 \cdot 1 < 2$ $1 \cdot 1 < 2$

Timed Büchi Automata

not simple!

Definition. The set $\Phi(X)$ of **clock constraints** over X is defined inductively by

$$\delta ::= x \le c \mid c \le x \mid \neg\delta \mid \delta_1 \wedge \delta_2$$

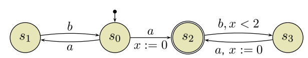
where $x \in X$ and $c \in \mathbb{Q}$ is a rational constant.

Definition. A **timed Büchi automaton** (TBA) $\mathcal{A}$ is a tuple $(\Sigma, S, S_0, X, E, F)$, where

- Σ is an alphabet,
- S is a finite set of states, $S_0 \subseteq S$ is a set of start states,
- X is a finite set of clocks, and
- $E \subseteq S \times S \times \Sigma \times 2^X \times \Phi(X)$ gives the set of transitions.
  An edge $(s, s', a, \lambda, \delta)$ represents a transition from state s to state s' on input symbol a. The set $\lambda \subseteq X$ gives the clocks to be reset with this transition, and δ is a clock constraint over X.
- $F \subseteq S$ is a set of **accepting states**.

Example: TBA

$$\mathcal{A} = (\Sigma, S, S_0, X, E, F)$$
$$(s, s', a, \lambda, \delta) \in E$$



(Accepting) TBA Runs

Definition. A **run** r, denoted by $(\bar{s}, \bar{\nu})$, of a TBA $(\Sigma, S, S_0, X, E, F)$ over a timed word (σ, τ) is an **infinite** sequence of the form

$$r : \langle s_0, \nu_0\rangle \xrightarrow[\tau_1]{\sigma_1} \langle s_1, \nu_1\rangle \xrightarrow[\tau_2]{\sigma_2} \langle s_2, \nu_2\rangle \xrightarrow[\tau_3]{\sigma_3} \ldots$$

with $s_i \in S$ and $\nu_i : X \to \mathbb{R}_0^+$, satisfying the following requirements:

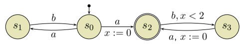
- **Initiation**: $s_0 \in S_0$ and $\nu(x) = 0$ for all $x \in X$.
- **Consecution**: for all $i \ge 1$, there is an edge in E of the form $(s_{i-1}, s_i, \sigma_i, \lambda_i, \delta_i)$ such that
  - $(\nu_{i-1} + (\tau_i - \tau_{i-1}))$ satisfies δ_i and
  - $\nu_i = (\nu_{i-1} + (\tau_i - \tau_{i-1}))[\lambda_i := 0]$.

The set $inf(r) \subseteq S$ consists of those states $s \in S$ such that $s = s_i$ for infinitely many $i \ge 0$.

Definition. A run $r = (\bar{s}, \bar{\nu})$ of a TBA over timed word (σ, τ) is called (an) **accepting** (run) if and only if $inf(r) \cap F \neq \emptyset$.

Example: (Accepting) Runs

$r : \langle s_0, \nu_0 \rangle \xrightarrow[\tau_1]{\sigma_1} \langle s_1, \nu_1 \rangle \xrightarrow[\tau_2]{\sigma_2} \langle s_2, \nu_2 \rangle \xrightarrow[\tau_3]{\sigma_3} \dots$ initial and $(s_{i-1}, s_i, \sigma_i, \lambda_i, \delta_i) \in E$, s.t. $(\nu_{i-1} + (\tau_i - \tau_{i-1})) \models \delta_i$, $\nu_i = (\nu_{i-1} + \underbrace{(\tau_i - \tau_{i-1})}_{\text{delay}})[\lambda_i := 0]$. Accepting iff $inf(r) \cap F \neq \emptyset$.



Timed word: $(a,1), (b,2), (a,3), (b,4), (a,5), (b,6), \dots$

- Can we construct **any run**? Is it accepting?

  $\langle s_0, x=0 \rangle \xrightarrow[1.0]{a} \langle s_2, 0 \rangle \xrightarrow[2.0]{b} \langle s_3, 1.0 \rangle \dots$ ✓

- Can we construct a **non-run**?

- Can we construct a **(non-)accepting run**?

  $\langle s_0, 0 \rangle \xrightarrow[1.0]{a} \langle s_1, 1 \rangle \xrightarrow[2.0]{b} \langle s_0, 2 \rangle \xrightarrow[3.0]{a} \langle s_1, 3 \rangle \dots$

The Language of a TBA

Definition. For a TBA $\mathcal{A}$, the **language** $L(\mathcal{A})$ of timed words it accepts is defined to be the set

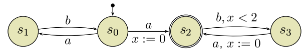
$$\{(\sigma, \tau) \mid \mathcal{A} \text{ has an accepting run over } (\sigma, \tau)\}.$$

For short: $L(\mathcal{A})$ is the **language of $\mathcal{A}$**.

Definition. A timed language L is a **timed regular language** if and only if $L = L(\mathcal{A})$ for **some** TBA $\mathcal{A}$.

Example: Language of a TBA

$L(\mathcal{A}) = \{(\sigma, \tau) \mid \mathcal{A} \text{ has an accepting run over } (\sigma, \tau)\}.$



Claim:

$$L(\mathcal{A}) = L_{crt} \; (= \{((ab)^\omega, \tau) \mid \exists\, i \; \forall\, j \geq i : (\tau_{2j} < \tau_{2j-1} + 2)\})$$

Question: Is L_{crt} timed regular or not?

The Universality Problem is Undecidable for TBA

[Alur and Dill, 1994]

The Universality Problem

- **Given:** A TBA $\mathcal{A}$ over alphabet Σ.
- **Question:** Does $\mathcal{A}$ accept all timed words over Σ?

  In other words: Is $L(\mathcal{A}) = \{(\sigma, \tau) \mid \sigma \in \Sigma^\omega, \tau \text{ time sequence}\}$.

$\Sigma = \{a, b, c\}$ $\mathcal{A}$:

The Universality Problem

- **Given:** A TBA $\mathcal{A}$ over alphabet Σ.
- **Question:** Does $\mathcal{A}$ accept all timed words over Σ?

  In other words: Is $L(\mathcal{A}) = \{(\sigma, \tau) \mid \sigma \in \Sigma^\omega, \tau \text{ time sequence}\}$.

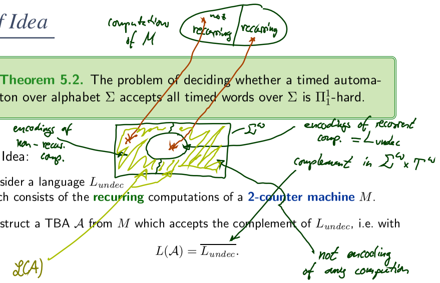
Theorem 5.2. The problem of deciding whether a timed automaton over alphabet Σ accepts all timed words over Σ is Π^1_1-hard.

("The class Π^1_1 consists of highly undecidable problems, including some nonarithmetical sets (for an exposition of the analytical hierarchy consult, see for instance [Rogers, 1967].)

Recall: With classical Büchi Automata (untimed), this is different:

- Let $\mathcal{B}$ be a Büchi Automaton over Σ.
- $\mathcal{B}$ is universal if and only if $\overline{L(\mathcal{B})} = \emptyset$.
- $\mathcal{B}'$ such that $L(\mathcal{B}') = \overline{L(\mathcal{B})}$ is effectively computable.
- Language emptyness is decidable for Büchi Automata.

Proof Idea



Theorem 5.2. The problem of deciding whether a timed automaton over alphabet Σ accepts all timed words over Σ is Π^1_1-hard.

Proof Idea:

- Consider a language L_{undec} which consists of the **recurring** computations of a **2-counter machine** M.
- Construct a TBA $\mathcal{A}$ from M which accepts the complement of L_{undec}, i.e. with
$$L(\mathcal{A}) = \overline{L_{undec}}.$$
- Then $\mathcal{A}$ is universal if and only if L_{undec} is empty...

...which is the case if and only if M **doesn't have** a recurring computation.

Once Again: 2-Counter Mach. (Different Flavour)

A **two-counter machine** M

- has two **counters** C, D and
- a finite **program** consisting of n instructions.
- An **instruction increments or decrements** one of the counters, or **jumps**, here even non-deterministically.
- A **configuration** of M is a triple $\langle i, c, d\rangle$:

  program counter $i \in \{1, \dots, n\}$, values $c, d \in \mathbb{N}_0$ of C and D.

- A **computation** of M is an infinite consecutive sequence
$$\langle 1, 0, 0\rangle = \langle i_0, c_0, d_0\rangle, \langle i_1, c_1, d_1\rangle, \langle i_2, c_2, d_2\rangle, \dots$$
  that is, $\langle i_{j+1}, c_{j+1}, d_{j+1}\rangle$ is a result executing instruction i_j at $\langle i_j, c_j, d_j\rangle$.

A computation of M is called **recurring** iff $i_j = 1$ for infinitely many $j \in \mathbb{N}_0$.

Step 1: The Language of Recurring Computations

- Let M be a 2CM with n instructions.

Wanted: A timed language L_{undec} (over some alphabet) representing exactly the recurring computations of M.
(In particular s.t. $L_{undec} = \emptyset$ if and only if M has no recurring computation.)

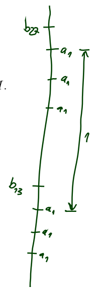
- Choose $\Sigma = \{b_1, \dots, b_n, a_1, a_2\}$ as alphabet.
- We represent a configuration $\langle i, c, d\rangle$ of M by the sequence
$$b_i \underbrace{a_1 \dots a_1}_{c \text{ times}} \underbrace{a_2 \dots a_2}_{d \text{ times}} = b_1 a_1^c a_2^d$$

Step 1: The Language of Recurring Computations

Let L_{undec} be the set of the timed words (σ, τ) with

- σ is of the form $b_{i_1} a_1^{c_1} a_2^{d_1} b_{i_2} a_1^{c_2} a_2^{d_2} \dots$
- $\langle i_1, c_1, d_1\rangle, \langle i_2, c_2, d_2\rangle, \dots$ is a recurring computation of M.
- For all $j \in \mathbb{N}_0$,
  - the time of b_{i_j} is j.
  - if $c_{j+1} = c_j$:
    for every a_1 at time t in the interval $[j, j+1]$
    there is an a_1 at time $t+1$,
  - if $c_{j+1} = c_j + 1$:
    for every a_1 at time t in the interval $[j+1, j+2]$,
    except for the last one, there is an a_1 at time $t-1$,
  - if $c_{j+1} = c_j - 1$:
    for every a_1 at time t in the interval $[j, j+1]$,
    except for the last one, there is an a_1 at time $t+1$,

And analogously for the a_2's.



Step 2: Construct "Observer" for $\overline{L_{undec}}$

Wanted: A TBA $\mathcal{A}$ such that $L(\mathcal{A}) = \overline{L_{undec}}$,
i.e., $\mathcal{A}$ accepts a timed word (σ, τ) if and only if $(\sigma, \tau) \notin L_{undec}$.

Approach: What are the reasons for a timed word **not to be** in L_{undec}?

Recall: (σ, τ) is **in** L_{undec} if and only if:

- $\sigma = b_{i_1} a_1^{c_1} a_2^{d_1} b_{i_2} a_1^{c_2} a_2^{d_2}$
- $\langle i_1, c_1, d_1\rangle, \langle i_2, c_2, d_2\rangle, \dots$
  is a recurring computation of M.
- the time of b_{i_j} is j,
- if $c_{j+1} = c_j$ $(= c_j + 1, = c_j - 1)$: ...

(i) The b_i at time $j \in \mathbb{N}$ is missing, or there is a spurious b_i at time $t \in]j, j+1[$.
(ii) The prefix of the timed word with times $0 \le t < 1$ doesn't encode $\langle 1, 0, 0\rangle$.
(iii) The timed word is not recurring, i.e. it has only finitely many b_i.
(iv) The configuration encoded in $[j+1, j+2[$ doesn't faithfully represent the effect of instruction b_i on the configuration encoded in $[j, j+1[$.

Step 2: Construct "Observer" for $\overline{L_{undec}}$

Wanted: A TBA $\mathcal{A}$ such that $L(\mathcal{A}) = \overline{L_{undec}}$,
i.e., $\mathcal{A}$ accepts a timed word (σ, τ) if and only if $(\sigma, \tau) \notin L_{undec}$.

Approach: What are the reasons for a timed word **not to be** in L_{undec}?

(i) The b_i at time $j \in \mathbb{N}$ is missing, or there is a spurious b_i at time $t \in]j, j+1[$.
(ii) The prefix of the timed word with times $0 \le t < 1$ doesn't encode $\langle 1, 0, 0\rangle$.
(iii) The timed word is not recurring, i.e. it has only finitely many b_i.
(iv) The configuration encoded in $[j+1, j+2[$ doesn't faithfully represent the effect of instruction b_i on the configuration encoded in $[j, j+1[$.

Plan: Construct a TBA $\mathcal{A}_0$ for case (i), a TBA $\mathcal{A}_{init}$ for case (ii), a TBA $\mathcal{A}_{recur}$ for case (iii), and one TBA $\mathcal{A}_i$ for each instruction for case (iv).

Then set
$$\mathcal{A} = \mathcal{A}_0 \cup \mathcal{A}_{init} \cup \mathcal{A}_{recur} \cup \bigcup_{1 \le i \le n} \mathcal{A}_i$$

Step 2.(i): Construct $\mathcal{A}_0$

(i) The b_i at time $j \in \mathbb{N}$ is missing, or there is a spurious b_i at time $t \in]j, j+1[$.

[Alur and Dill, 1994]: "It is easy to construct such a timed automaton."

Step 2.(ii): Construct $\mathcal{A}_{init}$

(ii) The prefix of the timed word with times $0 \le t < 1$ doesn't encode $\langle 1, 0, 0 \rangle$.

- It accepts

$$\{(\sigma_j, \tau_j)_{j \in \mathbb{N}_0} \mid (\sigma_0 \neq b_1) \vee (\tau_0 \neq 0) \vee (\tau_1 \neq 1)\}.$$

Step 2.(iii): Construct $\mathcal{A}_{recur}$

(iii) The timed word is not recurring, i.e. it has only finitely many b_i.

- $\mathcal{A}_{recur}$ accepts words with only finitely many b_i.

Step 2.(iv): Construct $\mathcal{A}_i$

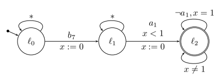
(iv) The configuration encoded in $[j+1, j+2[$ doesn't faithfully represent the effect of instruction b_i on the configuration encoded in $[j, j+1[$.

Example: assume instruction 7 is:

Increment counter D and jump non-deterministically to instruction 3 or 5.

Once again: stepwise. $\mathcal{A}_7$ is $\mathcal{A}_7^1 \cup \dots \cup \mathcal{A}_7^6$.

- $\mathcal{A}_7^1$ accepts words with b_7 at time j but neither b_3 nor b_5 at time $j+1$. "Easy to construct."
- $\mathcal{A}_7^2$ is



- $\mathcal{A}_7^3$ accepts words which encode unexpected increment of counter C.
- $\mathcal{A}_7^4, \dots, \mathcal{A}_7^6$ accept words with missing decrement of D.

Aha, And...?

Consequences: Language Inclusion

- **Given:** Two TBAs $\mathcal{A}_1$ and $\mathcal{A}_2$ over alphabet B.
- **Question:** Is $\mathcal{L}(\mathcal{A}_1) \subseteq \mathcal{L}(\mathcal{A}_2)$?

Possible applications of a decision procedure:

- Characterise the allowed behaviour as $\mathcal{A}_2$ and model the design as $\mathcal{A}_1$.
- Automatically check whether the behaviour of the design is a subset of the allowed behaviour.

- If **language inclusion** was decidable, then we could use it to decide universality of $\mathcal{A}$ by checking

$$\mathcal{L}(\mathcal{A}_{univ}) \subseteq \mathcal{L}(\mathcal{A})$$

where $\mathcal{A}_{univ}$ is **any** universal TBA (which is easy to construct).

Consequences: Complementation

- **Given:** A timed regular language W over B (that is, there is a TBA $\mathcal{A}$ such that $\mathcal{L}(\mathcal{A}) = W$).
- **Question:** Is $\overline{W}$ timed regular?

Possible applications of a decision procedure:

- Characterise the allowed behaviour as $\mathcal{A}_2$ and model the design as $\mathcal{A}_1$.
- Automatically construct $\mathcal{A}_3$ with $L(\mathcal{A}_3) = \overline{L(\mathcal{A}_2)}$ and check

$$L(\mathcal{A}_1) \cap L(\mathcal{A}_3) = \emptyset,$$

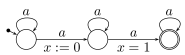
  that is, whether the design has any non-allowed behaviour.
- Taking for granted that:
  - The intersection automaton is effectively computable.
  - The emptyness problem for Büchi automata **is decidable**. (Proof by construction of region automaton [Alur and Dill, 1994].)

Consequences: Complementation

- **Given:** A timed regular language W over B (that is, there is a TBA $\mathcal{A}$ such that $\mathcal{L}(\mathcal{A}) = W$).
- **Question:** Is $\overline{W}$ timed regular?

- If the class of timed regular languages were closed under **complementation**, "the complement of the inclusion problem is recursively enumerable. This contradicts the Π_1^1-hardness of the inclusion problem." [Alur and Dill, 1994]

A non-complementable TBA $\mathcal{A}$:



$$\mathcal{L}(\mathcal{A}) = \{(a^\omega, (t_i)_{i\in\mathbb{N}_0}) \mid \exists\, i \in \mathbb{N}_0\ \exists\, j > i : (t_j = t_i + 1)\}$$

Complement language:

$$\overline{\mathcal{L}(\mathcal{A})} = \{(a^\omega, (t_i)_{i\in\mathbb{N}_0}) \mid \text{no two } a \text{ are separated by distance } 1\}.$$

Beyond Timed Regular

Beyond Timed Regular

With clock constraints of the form

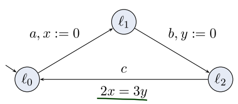
$$x + y \le x' + y'$$

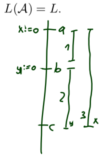
we can describe timed languages which are not timed regular.

In other words:

- There are strictly more timed languages than timed **regular** languages.
- There exists timed languages L such that there exists no $\mathcal{A}$ with $L(\mathcal{A}) = L$.

Example:





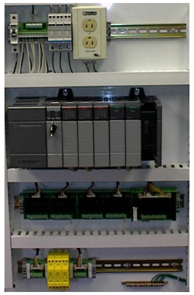
$$\{((abc)^\omega, \tau) \mid \forall\, j.(\tau_{3j} - \tau_{3j-1}) = 2(\tau_{3j-1} - \tau_{3j-2})\}$$

hat is a PLC?

hat's special about PLC?

- microprocessor, memory, timers
- digital (or analog) I/O ports
- possibly RS 232, fieldbuses, networking
- robust hardware
- reprogrammable
- **standardised programming model** (IEC 61131-3)

here are PLC employed?



- mostly **process automatisation**
  - production lines
  - packaging lines
  - chemical plants
  - power plants
  - electric motors, pneumatic or hydraulic cylinders
  - ...
- not so much: **product automatisation**, there
  - tailored or OTS controller boards
  - embedded controllers
  - ...

o are PLC programmed?

- PLC have in common that they operate in a cyclic manner:

read inputs
compute
write outputs

- Cyclic operation is repeated until external interruption (such as shutdown or reset).
- Cycle time: typically a few milliseconds. [?]
- Programming for PLC means providing the "**compute**" part.
- Input/output values are available via designated local variables.

hy study PLC?

- **Note:**
  the discussion here is **not limited** to PLC and IEC 61131-3 languages.
- Any programming language on an operating system with **at least one** real-time clock will do.
  (Where a **real-time clock** is a piece of hardware such that,
  - we can program it to wait for t time units,
  - we can query whether the set time has elapsed,
  - if we program it to wait for t time units, it does so with negligible deviation.)
- And strictly speaking, we don't even need "full blown" operating systems.
- PLC are just a formalisation on a good level of abstraction:
  - there are inputs **somehow** available as local variables,
  - there are outputs **somehow** available as local variables,
  - **somehow**, inputs are polled and outputs updated atomically,
  - there is **some** interface to a real-time clock.

References

[Alur and Dill, 1994] Alur, R. and Dill, D. L. (1994). A theory of timed automata. *Theoretical Computer Science*, 126(2):183–235.

[Olderog and Dierks, 2008] Olderog, E.-R. and Dierks, H. (2008). *Real-Time Systems - Formal Specification and Automatic Verification*. Cambridge University Press.