

Softwaretechnik / Software-Engineering

*Lecture 18: Runtime Verification, Review &
Wrapup*

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Topic Area Code Quality Assurance: Content

- VL 15
 - Introduction and Vocabulary
- VL 16
 - Limits of Software Testing
 - Glass-Box Testing
 - └─● Statement-, branch-, term-coverage.
- ⋮
 - Other Approaches
 - └─● Model-based testing,
 - └─● Runtime verification.
- VL 17
 - Software quality assurance in a larger scope.
- ⋮
 - Program Verification
 - └─● partial and total correctness,
 - └─● Proof System PD.
- VL 18
 - Runtime Verification
- ⋮
 - Review
 - Code QA: Discussion

- **Runtime-Verification**

- └ (● Idea
- └ (● Assertions
- └ (● LSC-Observers

- **Reviews**

- └ (● **Roles** and **artefacts**
- └ (● Review **procedure**
- └ (● Stronger and weaker **variants**

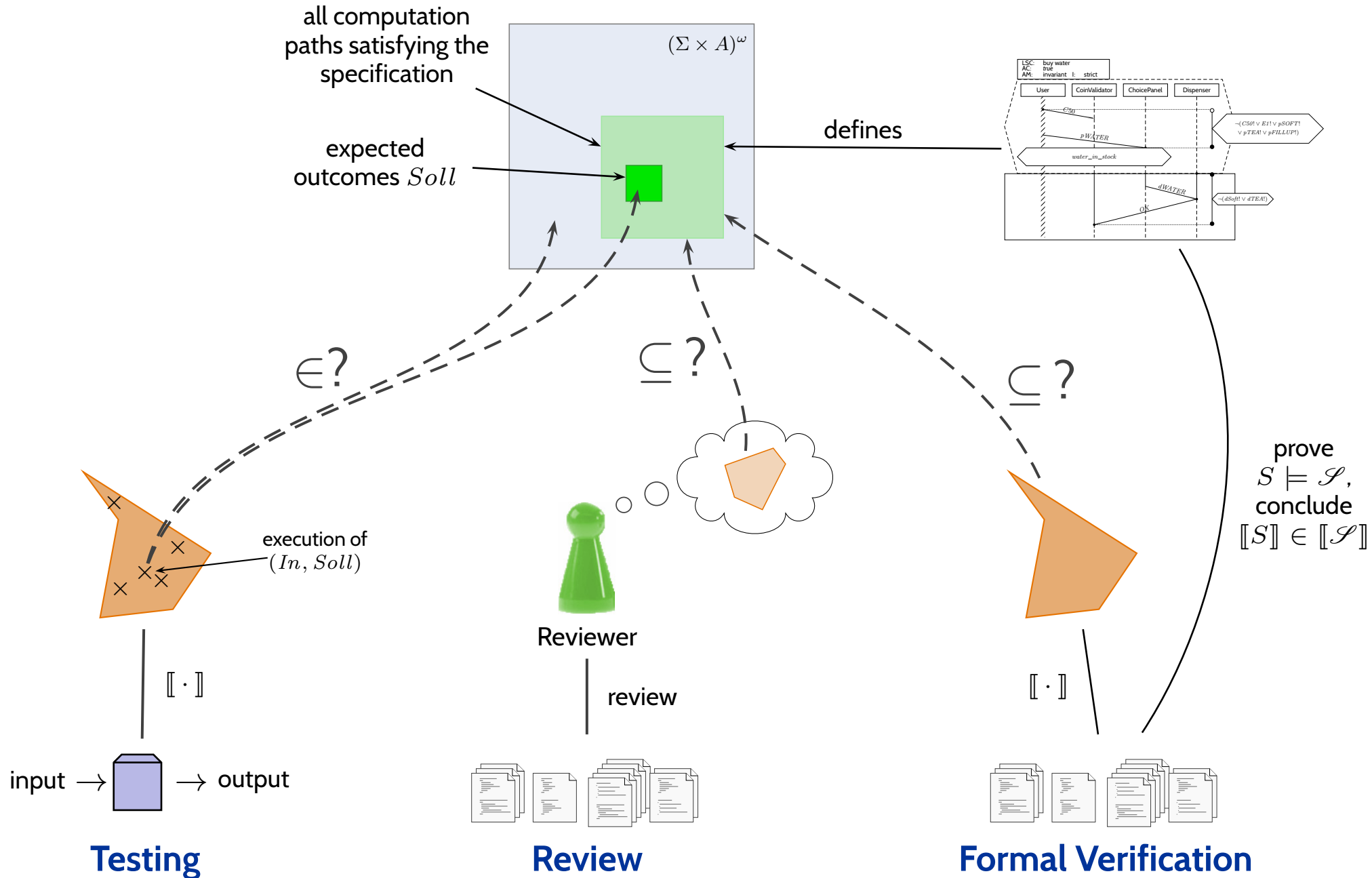
- **Do's and Don'ts** in Code QA

- **Code QA Techniques** Revisited

- └ (● **Test**
- └ (● **Runtime-Verification**
- └ (● **Review**
- └ (● **Static Checking**
- └ (● **Formal Verification**

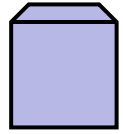
- **Dependability**

Recall: Three Basic Directions



Run-Time Verification

Run-Time Verification: Idea



Software S

- Assume, there is a function f in software S with the following specification:
 - pre-condition:** p , **post-condition:** q .
- Computation paths of S may look like this:

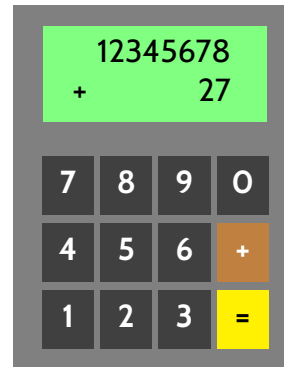
$$\sigma_0 \xrightarrow{\alpha_1} \sigma_1 \xrightarrow{\alpha_2} \sigma_2 \cdots \xrightarrow{\alpha_{n-1}} \sigma_n \xrightarrow{\text{call } f} \sigma_{n+1} \cdots \sigma_m \xrightarrow{f \text{ returns}} \sigma_{m+1} \cdots$$

- Assume there are functions $check_p$ and $check_q$, which **check** whether p and q hold at the current program state, and which **do not modify the program state** (except for program counter).
- Idea:** create software S' by
 - (i) extending S by implementations of $check_p$ and $check_q$,
 - (ii) call $check_p$ right after entering f ,
 - (iii) call $check_q$ right before returning from f .
- For S' , obtain computation paths like:

$$\sigma_0 \xrightarrow{\alpha_1} \sigma_1 \xrightarrow{\alpha_2} \sigma_2 \cdots \xrightarrow{\alpha_{n-1}} \sigma_n \xrightarrow{\text{call } f} \sigma_{n+1} \xrightarrow{check_p} \sigma'_{n+1} \cdots \sigma_m \xrightarrow{check_q} \sigma'_m \xrightarrow{f \text{ returns}} \sigma_{m+1} \cdots$$

- If $check_p$ and $check_q$ notify us of violations of p or q , then we are **notified of f violating its specification** when running S' (= at run-time).

Run-Time Verification: Example



```
1  int main() {  
2  
3      while (true) {  
4          int x = read_number();  
5          int y = read_number();  
6  
7          int sum = add( x, y );  
8  
9          verify_sum( x, y, sum );  
10         display(sum);  
11     }  
12 }  
13 }
```

```
1  void verify_sum( int x, int y,  
2                  int sum )  
3  {  
4      if (sum != (x+y)  
5          || (x + y > 99999999  
6              && !(sum < 0)))  
7      {  
8          fprintf( stderr ,  
9                  "verify_sum:_error\n" );  
10         abort();  
11     }  
12 }
```

A Very Useful Special Case: Assertions

- Maybe the simplest instance of runtime verification: **Assertions**.
- Available in standard libraries of many programming languages (C, C++, Java, ...).
- For example, the C standard library manual reads:

```
1  ASSERT(3)          Linux Programmer's Manual          ASSERT(3)
2
3  NAME
4      assert — abort the program if assertion is false
5
6  SYNOPSIS
7      #include <assert.h>
8
9      void assert(scalar expression);
10
11 DESCRIPTION
12      [...] the macro assert() prints an error message to stanâ
13      dard error and terminates the program by calling abort(3) if expression
14      is false (i.e., compares equal to zero).
15
16      The purpose of this macro is to help the programmer find bugs in his or her
17      program. The message "assertion failed in file foo.c, function
18      do_bar(), line 1287" is of no help at all to a user.
```

- In C code, assert can be **disabled** in **production code** (-D NDEBUG).

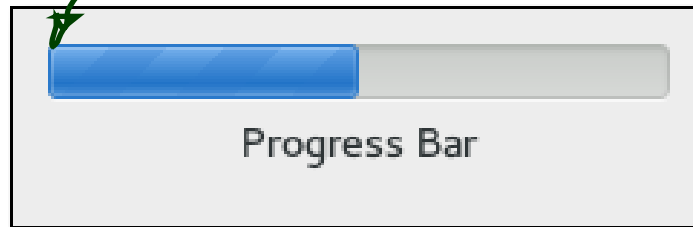
Assertions At Work

- The abstract f -example from **run-time verification**:

```
1 void f( ... ) {  
2     assert( p );  
3     ...  
4     assert( q );  
5 }
```

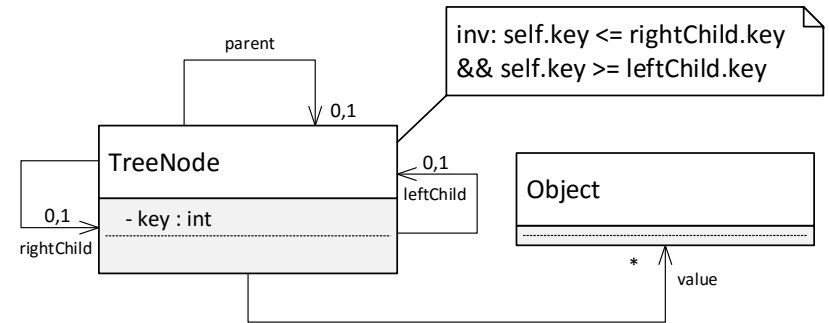
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1 ASSERT(3)      Linux Programmer's Manual      ASSERT(3)  
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17     program. The message "assertion failed in file foo.c, function  
18     do_bar(), line 1287" is of no help at all to a user.
```

- Compute the width of a progress bar. *window_left*



```
1  
2 int progress_bar_width( int progress , int window_left , int window_right )  
3 {  
4     assert( window_left <= window_right ); /* pre-condition */  
5     ...  
6     /* treat special cases 0 and 100 */  
7     ...  
8     assert( 0 < progress && progress < 100 ); // extremal cases already treated  
9     ...  
10    assert( window_left <= r && r <= window_right ); /* post-condition */  
11    return r;  
12 }
```

Assertions At Work II



- Recall the **structure model** with Proto-OCL constraint from Exercise Sheet 4.
- Assume, we add a method `set_key()` to class **TreeNode**:

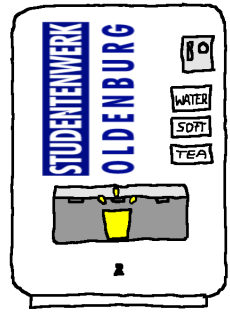
```
1 class TreeNode {
2
3     private int key;
4     TreeNode parent, leftChild, rightChild;
5
6     public int get_key() { return key; }
7
8     public void set_key( int new_key ) {
9         key = new_key;
10    }
11 }
```

- We can **check consistency** with the Proto-OCL constraint at runtime by using assertions:

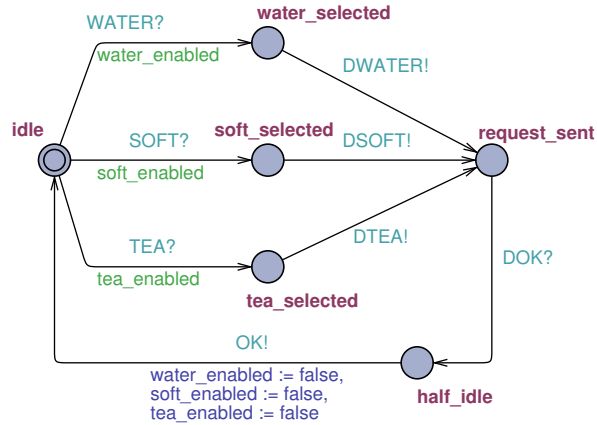
```
1     public void set_key( int new_key ) {
2         assert ( parent == null || parent.get_key() <= new_key );?
3         assert ( leftChild == null || new_key <= leftChild.get_key() );
4         assert ( rightChild == null || new_key <= rightChild.get_key() );
5
6         key = new_key;
7     }
```

- Use `java -ea ...` to **enable assertion checking** (disabled by default).
(cf. <https://docs.oracle.com/javase/8/docs/technotes/guides/language/assert.html>)

More Complex Run-Time Verification: LSC Observers



ChoicePanel:

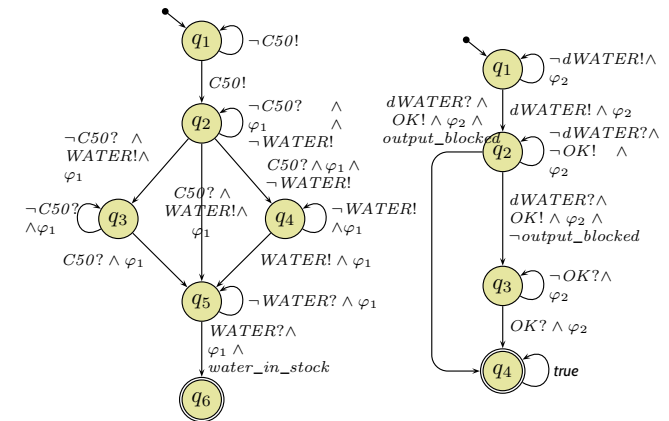
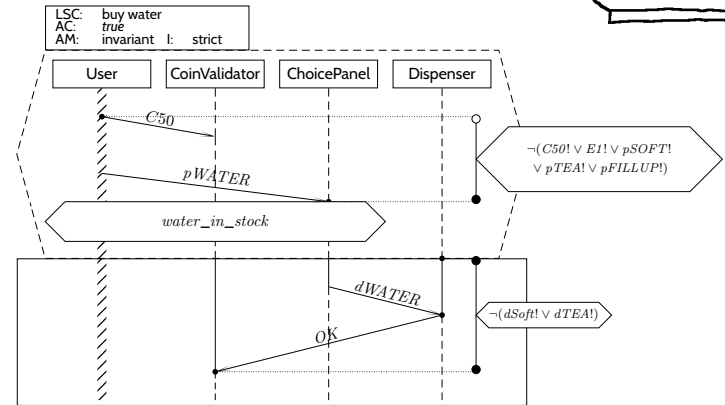


```

st : { idle, wsel, ssel, tsel, reqs, half };

take_event( E : { TAU, WATER, SOFT, TEA, ... } ) {
  bool stable = 1;
  switch (st) {
    case idle :
      switch (E) {
        case WATER :
          if (water_enabled) { st := wsel; stable := 0; }
          ;;
        case SOFT :
          ;;
        case TEA :
          ;;
      }
    case wsel:
      switch (E) {
        case TAU :
          send_DWATER(); st := reqs;
          hey_observer_I_just_sent_DWATER();
          ;;
      }
      ;;
  }
}

```



Run-Time Verification: Discussion

- **Experience:**

During development, **assertions** for pre/post conditions and intermediate invariants are an **extremely powerful** tool with a **very attractive gain/effort ratio** (low effort, high gain).

- Assertions effectively work as **safe-guard against unexpected use** of functions and **regression**, e.g. during later maintenance or efficiency improvement.
- Can serve as **formal** (support of) **documentation**:

“Dear reader, at this point in the program, I expect condition *expr* to hold, because...”

- **Development- vs. Release Versions:**

- Common practice:
 - development version **with** run-time verification enabled (cf. `assert(3)`),
 - release version **without** run-time verification.

If run-time verification is enabled in a release version,

- ▷ ● software should **terminate as gracefully as possible** (e.g. try to save data),
- **save information from assertion failure** if possible for future analysis.

Risk: with bad luck, the software only behaves well **because of** the run-time verification code...

Then disabling run-time verification “breaks” the software. Yet very complex run-time verification may significantly slow down the software, so needs to be disabled...

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- **Do's and Don'ts** in Code QA

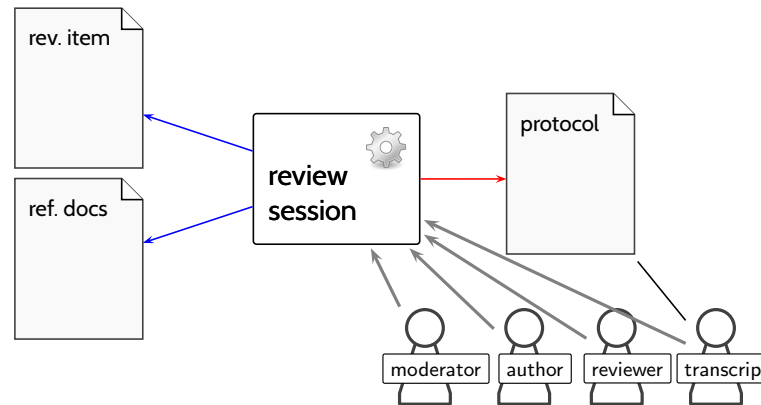
- **Code QA Techniques** Revisited

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- **Dependability**

Review

Reviews



- **Input to Review Session:**

- **Review item:** can be every closed, human-readable part of software (documentation, module, test data, installation manual, etc.)

/// **Social aspect:** it is an **artefact** which is examined, not the **human** (who created it).

- **Reference documents:** need to enable an assessment (requirements specification, guidelines (e.g. coding conventions), catalogue of questions (“all variables initialised?”), etc.)

- **Roles:**

Moderator: leads session, responsible for properly conducted procedure.

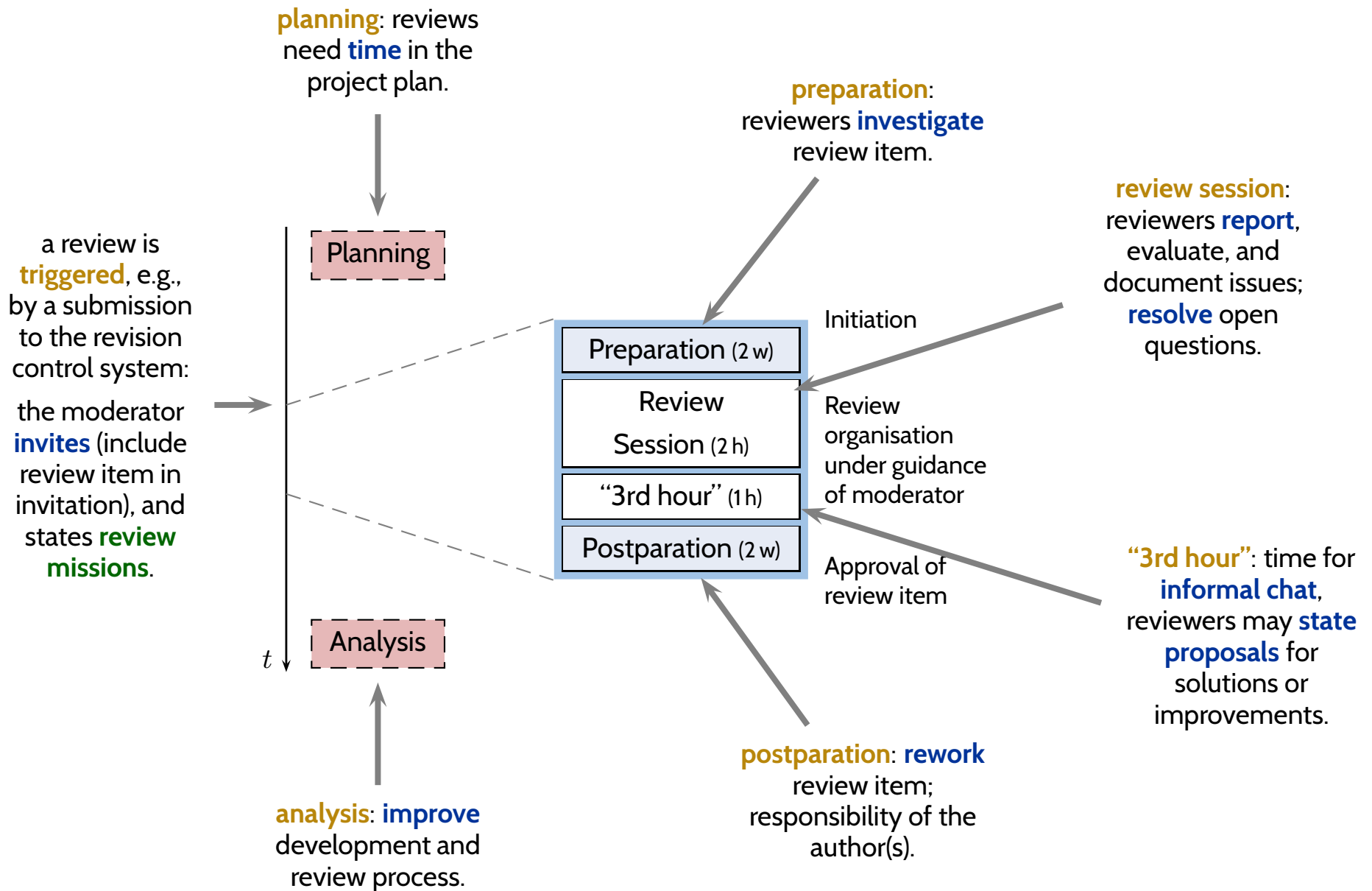
Author: (representative of the) creator(s) of the artefact under review; is present to listen to the discussions; can answer questions; does not speak up if not asked.

Reviewer(s): person who is able to judge the artefact under review; maybe different reviewers for different aspects (programming, tool usage, etc.), at best experienced in detecting inconsistencies or incompleteness.

Transcript Writer: keeps minutes of review session, can be assumed by author.

- The **review team** consists of everybody but the author(s).

Review Procedure Over Time



- Reviewers **re-assess** reworked review item (until **approval** is declared).

Review Rules (*Ludewig and Lichter, 2013*)

- (i) The **moderator** organises the review, issues invitations, supervises the review session.
- (ii) The **moderator** may terminate the review if conduction is not possible, e.g., due to inputs, preparation, or people missing.
- (iii) The review session is **limited to 2 hours**.
If needed: organise more sessions.
- (iv) The **review item** is under review, not the author(s).
Reviewers choose their words accordingly.
Authors neither defend themselves nor the review item.
- (v) Roles are **not mixed up**, e.g., the moderator does not act as reviewer.
(Exception: author may write transcript.)
- (vi) **Style** issues (outside fixed conventions) are **not discussed**.

- (vii) The **review team** is **not** supposed to **develop solutions**.
Issues are **not** noted down in form of **tasks** for the **author(s)**.
- (viii) Each **reviewer** gets the opportunity to present her/his findings appropriately.
- (ix) **Reviewers** need to reach **consensus** on issues, consensus is noted down.
- (x) **Issues** are classified as:
 - **critical** (review unusable for purpose),
 - **major** (usability severely affected),
 - **minor** (usability hardly affected),
 - **good** (no problem).
- (xi) The **review team** declares:
 - **accept without changes**,
 - **accept with changes**,
 - **do not accept**.
- (xii) The **protocol** is signed by all participants.

Stronger and Weaker Review Variants

- **Design and Code Inspection** (Fagan, 1976, 1986)

- deluxe variant of review,
- **approx. 50% more time, approx. 50% more errors found.**

- **Review**

- **Structured Walkthrough**

- simple variant of review:
 - **developer** moderates walkthrough-session,
 - **developer** presents artefact(s),
 - **reviewer** poses (prepared or spontaneous) questions,
 - issues are noted down,
- **Variation point:** do reviewers see the artefact before the session?
- **less effort, less effective.**

→ **disadvantages:** unclear responsibilities; “salesman”-**developer** may trick **reviewers**.

- **Comment** (“Stellungnahme”)

- **colleague(s)** of **developer** read artefacts,
- **developer** considers feedback.

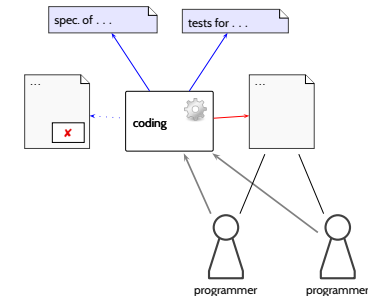
→ **advantage:** low organisational effort;

→ **disadvantages:** choice of colleagues may be biased; no protocol; consideration of comments at discretion of developer.

- **Careful Reading** (“Durchsicht”)

- done by **developer**,
- **recommendation:** “away from screen” (use print-out or different device and situation)

XP's pair programming ("on-the-fly review"?)



Some Final, General Guidelines

Do's and Don'ts in Code Quality Assurance



Avoid using special **examination versions** for examination.
(Test-harness, stubs, etc. **may have errors** which may cause false positives and (!) negatives.)



Avoid to **stop examination** when the first error is detected.

Clear: Examination should be aborted if the examined program is not executable at all.

Do not modify the artefact under examination **during** examination.

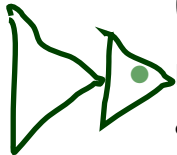
- otherwise, it is **unclear what exactly** has been examined (“moving target”),
(examination results need to be uniquely traceable to one artefact version.)

- fundamental flaws are sometimes **easier to detect**
with a **complete picture** of unsuccessful/successful tests,



- **changes are particularly error-prone**, should not happen “en passant” in examination,

- fixing flaws during examination may cause them to **go uncounted** in the **statistics**
(which we need for all kinds of estimation),



- roles **developer** and **examinor** are different anyway:
an **examinor** fixing flaws would **violate the role assignment**.



Do not switch (fine grained) between **examination** and **debugging**.

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Code Quality Assurance Techniques Revisited

Techniques Revisited

	auto- matic	prove “can run”	toolchain considered	exhaus- tive	prove correct	partial results	entry cost
Test	(✓)	✓	✓	✗	✗	✓	✓
Runtime- Verification							
Review							
Static Checking							
Verification							

Strengths:

- can be **fully automatic** (yet not easy for GUI programs);
- negative test **proves “program not completely broken”**, “can run” (or positive scenarios);
- **final product is examined**, thus toolchain and platform considered;
- one can stop at any time and take **partial results**;
- few, simple test cases are usually **easy to obtain**; *“one test case per feature”*
- provides **reproducible counter-examples** (good starting point for repair).

Weaknesses:

- (in most cases) **vastly incomplete**, thus no proofs of correctness;
- creating test cases for complex functions (or complex conditions) **can be difficult**;
- **maintenance** of many, complex test cases be **challenging**.
- executing many tests may need **substantial time** (but: can sometimes be run in parallel);

Techniques Revisited

	auto- matic	prove “can run”	toolchain considered	exhaus- tive	prove correct	partial results	entry cost
Test	(✓)	✓	✓	✗	✗	✓	✓
Runtime- Verification	✓	(✓)	✓	(✗)	✗	✓	(✓)
Review							
Static Checking							
Verification							

Strengths:

- **fully automatic** (once observers are in place);
- **provides counter-example**;
- (nearly) **final product is examined**, thus toolchain and platform considered;
- one can stop at any time and take **partial results**;
- **assert-statements have a very good effort/effect ratio**.

Weaknesses:

- counter-examples **not necessarily reproducible**;
- may negatively affect **performance**;
- code is changed, program may only run **because of** the observers;
- completeness depends on usage,
may also be **vastly incomplete**, so no correctness proofs;
- constructing observers for complex properties may be **difficult**,
one needs to learn how to construct observers.

Techniques Revisited

	auto- matic	prove “can run”	toolchain considered	exhaus- tive	prove correct	partial results	entry cost
Test	(✓)	✓	✓	✗	✗	✓	✓
Runtime- Verification	✓	(✓)	✓	(✗)	✗	✓	(✓)
Review	✗	✗	✗	(✓)	(✓)	✓	(✓)
Static Checking							
Verification							

Strengths:

- human readers can **understand the code**, may spot point errors;
- reported to be **highly effective**;
- one can stop at any time and take **partial results**;
- intermediate **entry costs**;
good effort/effect ratio achievable.

Weaknesses:

- no **tool support**;
- no results on actual execution, **toolchain not reviewed**;
- human readers may **overlook** errors; usually not aiming at proofs.
- does (in general) **not provide counter-examples**,
developers may deny existence of error.

Techniques Revisited

	auto- matic	prove “can run”	toolchain considered	exhaus- tive	prove correct	partial results	entry cost
Test	(✓)	✓	✓	✗	✗	✓	✓
Runtime- Verification	✓	(✓)	✓	(✗)	✗	✓	(✓)
Review	✗	✗	✗	(✓)	(✓)	✓	(✓)
Static Checking	✓	(✗)	✗	✓	(✓)	✓	(✗)
Verification							

Strengths:

- there are (commercial), **fully automatic** tools (lint, Coverity, Polyspace, etc.);
- some tools are **complete** (relative to assumptions on language semantics, platform, etc.);
- can be **faster than testing**;
- one can stop at any time and take **partial results**.

Weaknesses:

- no results on actual execution, **toolchain not reviewed**;
- can be very **resource consuming** (if few false positives wanted), e.g., code may need to be “designed for static analysis”.
- many false positives can be very **annoying to developers** (if fast checks wanted);
- distinguish **false from true positives** can be challenging;
- **configuring the tools** (to limit false positives) can be challenging.

Techniques Revisited

	auto- matic	prove “can run”	toolchain considered	exhaus- tive	prove correct	partial results	entry cost
Test	(✓)	✓	✓	✗	✗	✓	✓
Runtime- Verification	✓	(✓)	✓	(✗)	✗	✓	(✓)
Review	✗	✗	✗	(✓)	(✓)	✓	(✓)
Static Checking	✓	(✗)	✗	✓	(✓)	✓	(✗)
Verification	(✓)	✗	✗	✓	✓	(✗)	✗

Strengths:

- some **tool support** available (few commercial tools);
- **complete** (relative to assumptions on language semantics, platform, etc.);
- thus can provide **correctness proofs**;
- can prove correctness for **multiple language semantics and platforms** at a time;
- can be **more efficient than other techniques**.

Weaknesses:

- no results on actual execution, **toolchain not reviewed**;
- not many **intermediate results**: “half of a proof” may not allow any useful conclusions;
- **entry cost high**: significant training is useful to know how to deal with tool limitations;
- proving things is challenging: failing to find a proof does not allow any useful conclusion;
- **false negatives** (broken program “proved” correct) hard to detect.

Quality Assurance — Concluding Discussion

Proposal: Dependability Cases (*Jackson, 2009*)

- A **dependable** system is one you can **depend** on – that is, you can place your trust in it.

“Developers [should] **express the critical properties** and **make an explicit argument** that the system satisfies them.”

quality assurance – (1) A planned and systematic pattern of all actions necessary to provide **adequate confidence** that an item or product conforms to established technical requirements.
IEEE 610.12 (1990)

Proposed Approach:

- Identify the **critical requirements**, and determine what **level of confidence** is needed.
Most systems do also have **non-critical** requirements.
- Construct a **dependability case**:
 - an argument, that the software, in concert with other components, establishes the critical properties.
- The case should be
 - **auditable**: can (easily) be evaluated by third-party certifier.
 - **complete**: no holes in the argument, any assumptions that are not justified should be noted (e.g. assumptions on compiler, on protocol obeyed by users, etc.)
 - **sound**: e.g. should not claim full correctness [...] based on nonexhaustive testing; should not make unwarranted assumptions on independence of component failures; etc.

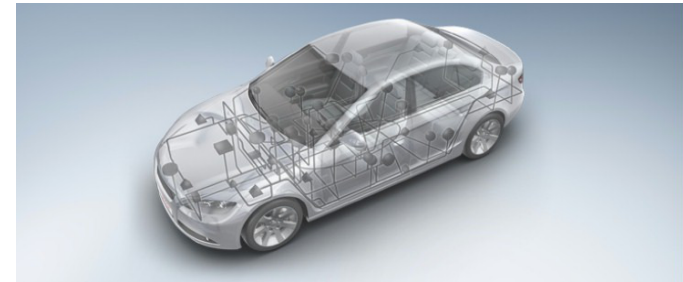
Critical Systems

Still, it seems like computer systems more or less **inevitably have errors**.

Then why...



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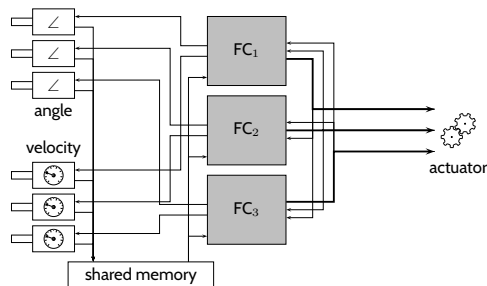


Robert Bosch GmbH

- ... do modern planes fly at all?

- (i) **very careful development**,
- (ii) **very thorough analysis**,
- (iii) **strong regulatory obligations**.

Plus: classical engineering wisdom for high reliability, like **redundancy**.

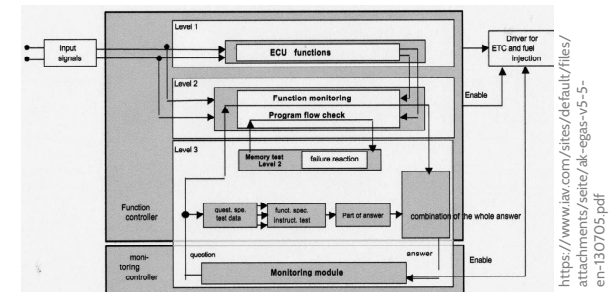


(Mrugalla et al., 2005)

- ... do modern cars drive at all?

- (i) **careful development**,
- (ii) **thorough analysis**,
- (iii) **regulatory obligations**.

Plus: classical engineering wisdom for high reliability, like **monitoring**.



Tell Them What You've Told Them...

- **Runtime Verification**
 - (as the name suggests) checks properties at **program run-time**,
 - a good **pinch of assert's** can be a valuable safe-guard against
 - **regressions**,
 - usage **outside specification**,
 - etc.

and serve as **formal documentation** of assumptions.
- **Review** (structured examination of artefacts by humans)
 - (mild variant) advocated in the XP approach,
 - **not uncommon**:
 - lead programmer reviews **all commits** from team members,
 - literature reports good effort/effect ratio achievable.
- All approaches to **code quality assurance** have their
 - **advantages** and **drawbacks**.
 - Which to use? It depends!
- **Dependability Cases**
 - an (auditable, complete, sound) argument, that a software has the **critical properties**.

References

References

Fagan, M. (1976). Design and code inspections to reduce errors in program development. *IBM Systems Journal*, 15(3):182–211.

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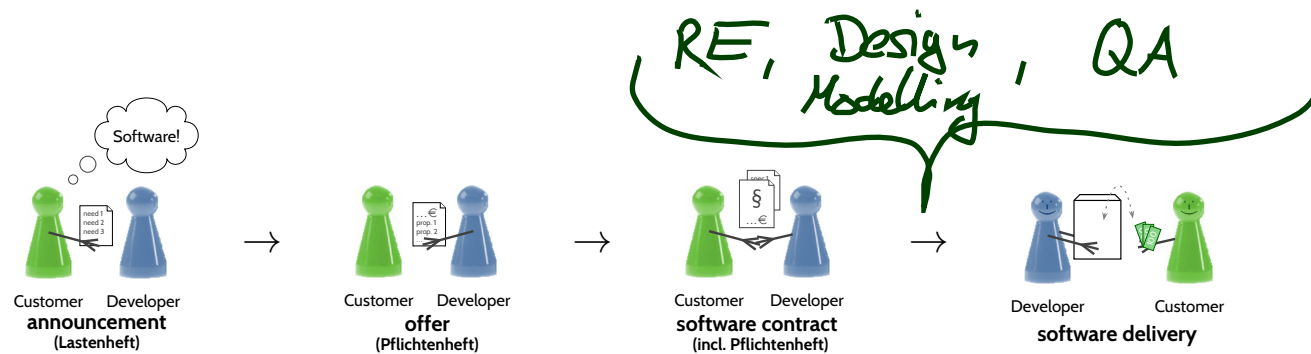
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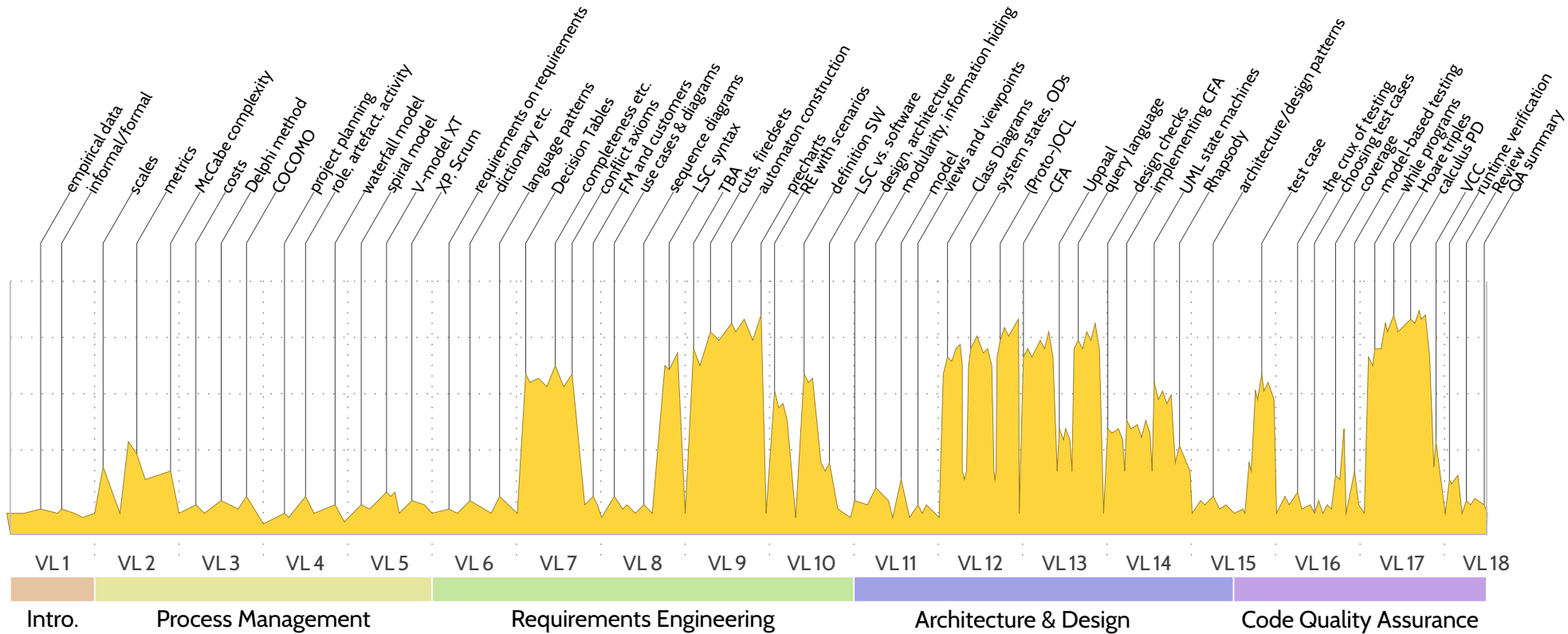
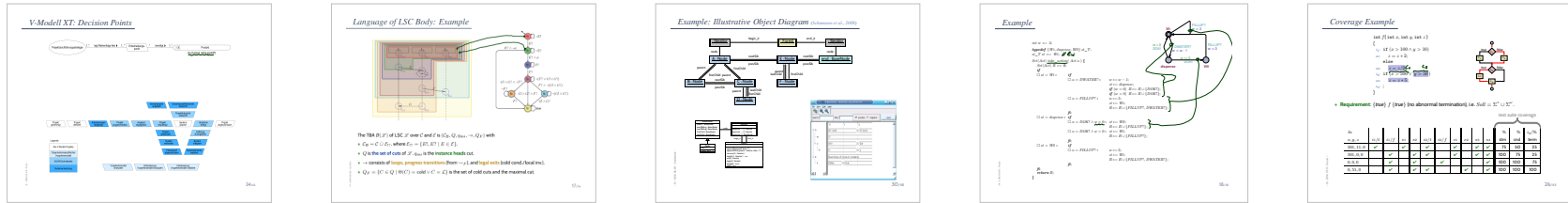
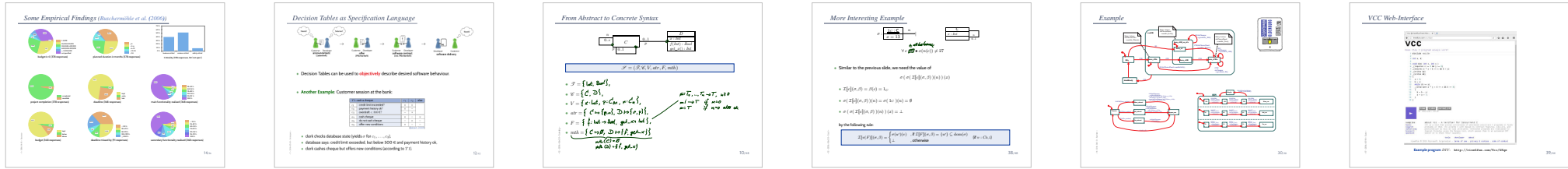
*Looking Back:
18 Lectures on Software Engineering*

Contents of the Course



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What Did We Do?



Expectations

- none, because mandatory course
- **overall**
 - ✓ well-structured lectures
 - (✓) praxis oriented
 - ✗ practical knowledge about planning, designing and testing software
 - ✓ improve skills in scientific work
 - (✓) more about scientific methods
- **other courses**
 - ✗ more on how courses are linked together
 - ✗ skills we need to organise SoPra
 - ✓ maybe transfer knowledge in SoPra



“real world”

- ✓ vocabulary and methods in professional software development
- ✓ learn how things work in a company, to easier integrate into teams, e.g., communication

• kinds of software

- ✓ embedded systems and software
- ✗ how to combine HW and SW parts

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Expectations Cont'd

• software development

- ✓ understand how software development practically works
- ✓ developing, maintaining software at bigger scale
- ✓ aspects of software development

• software project management

- ✓ learn what is important to plan
- ✓ how to structure the process of a project
- ✓ how to keep control of project, measure success
- ✗ which projects need full-time project manager
- ✗ which kind of documentation is really necessary
- ✗ want to get better in leading a team; how to lead team of engineers

• cost estimation

- ✓ how to estimate time and effort
- (✗) formal methods for better planning of projects
- ✗ tools which help planning

• quality

- ✓ learn ways how to judge quality based on the requirements
- ✓ avoid mistakes during software development
- ✓ make better programs, or make programs more efficiently

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Expectations Cont'd

● requirements

- ✓ formal ways to specify requirements
- ✓ learn techniques to reduce misunderstandings
- ✓ understand types of requirements
- (✓) learn how requirements are to be stated
- (✓) how to create requirements/specification document

● design

- ✓ techniques for design
- ✓ predict potential risks and crucial design errors
- (✗) come up with good design, learn how to design
- (✗) practical knowledge on application of design patterns
 - ✗ how to structure, compose components, how to define interfaces
 - ✗ standards for keeping parts of project compatible
 - ✗ how to guarantee a particular reliability

● Implementation

- (✓) modular programming, better documentation of big projects
 - ✗ more of computers and programming, write faster better programs
 - ✗ strengths and weaknesses of standards, training in their application
 - ✗ improve coding skills
 - ✗ how to increase (software) performance

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Expectations Cont'd

- **code quality assurance**

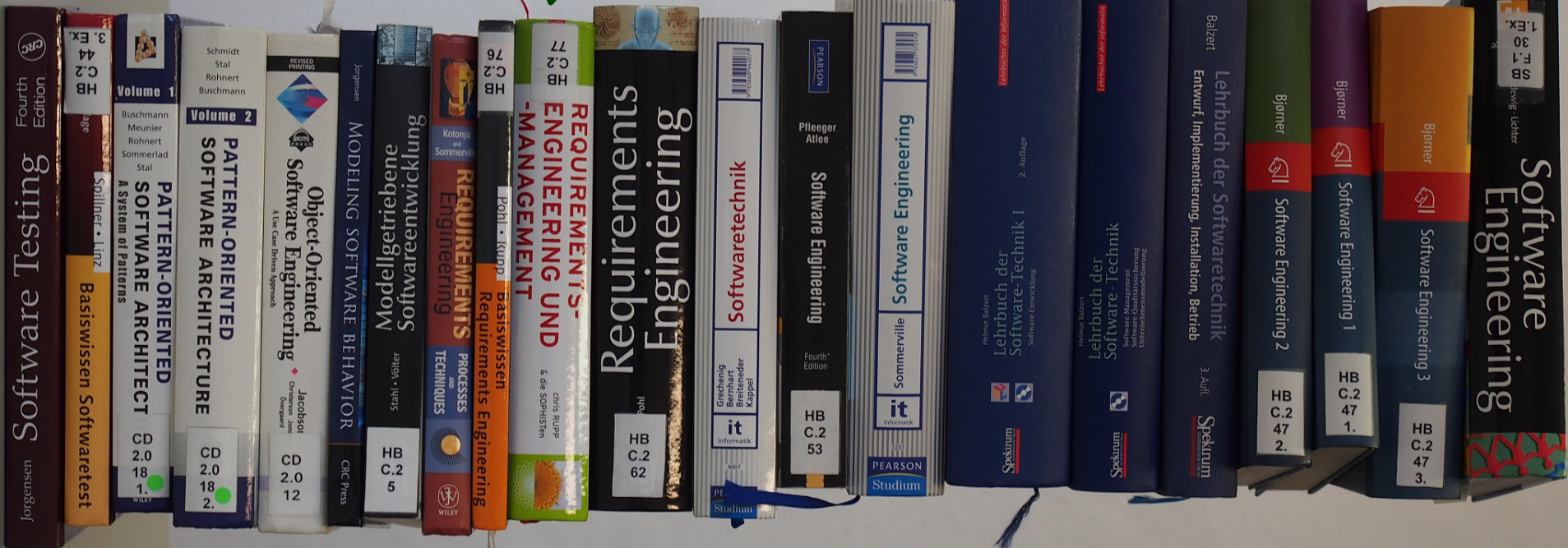
- ✓ methods for testing to guarantee high level of quality
- (✓) how to conduct most exhaustive test as possible in reasonable time
- ✓ formal methods like program verification
- ✗ learn about practical implementation of these tools

- **extra information**

- “will work as teacher”
- “want to work on medical software”
- “want to work in automotive industry”
- “worked as software-engineer”

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That's Today's Software Engineering — More or Less...



Coming Soon to Your Local Lecture Hall...

-18 - 2016-07-18 - Ssoon -



Thursday, 2016-07-21, 1200 to 1400:

Plenary Tutorial 6 & Questions Session

in 101-0-026 (right here)