

Prof. Dr. Andreas Podelski Dr. Matthias Heizmann Christian Schilling

Tutorial for Program Verification Bonus Sheet 13 (not discussed in the tutorial)

Exercise 1: Transition predicate abstraction

Consider the following modified version of the TPA algorithm. (Modification underlined: T composed with the *abstraction* of ρ_{τ})

 $\begin{array}{l} \boldsymbol{Algorithm} \ (\boldsymbol{TPA}^{cl}) \\ \textbf{Input:} \ \ \text{program} \ P = (\Sigma, \mathcal{T}, \rho) \\ \quad \text{set of transition predicates } \mathcal{P} \\ \quad \text{abstraction } \alpha \ \text{defined by } \mathcal{P} \\ \textbf{Output:} \ \ \text{set of abstract transitions } P^{\#} = \{T_1, \ldots, T_n\} \\ \quad \text{such that } T_1 \cup \cdots \cup T_n \ \text{is a transition invariant} \\ P^{\#} := \{\alpha(\rho_{\tau}) \mid \tau \in \mathcal{T}\} \\ \textbf{repeat} \\ P^{\#} := P^{\#} \cup \{\alpha(T \circ \underline{\alpha}(\rho_{\tau})) \mid T \in P^{\#}, \ \tau \in \mathcal{T}, \ \underline{\alpha}(T \circ \alpha(\rho_{\tau})) \neq \emptyset\} \\ \textbf{until no change} \end{array}$

(a) Prove or refute the following claim:

The set of abstract transitions computed by TPA^{cl} is a disjunctively well-founded transition invariant iff the set of abstract transitions computed by TPA is a disjunctively well-founded transition invariant.

(b) Think about a setting where we reapply the algorithm multiple times for the same set of transition predicates. What can be a possible advantage of TPA^{cl} over TPA?

Exercise 2: TPA with initial states

So far we considered only programs $P = (\Sigma, \mathcal{T}, \rho)$ where every state is an initial state. Let us now consider programs $P = (\Sigma, \Sigma_{\text{init}}, \mathcal{T}, \rho)$ where only the states in $\Sigma_{\text{init}} \subseteq \Sigma$ are initial states.

- (a) Give a program $P = (\Sigma, \Sigma_{\text{init}}, \mathcal{T}, \rho)$ whose transition relation R_P is not well-founded, but R_P restricted to the reachable states of P is well-founded. Give an informal explanation why for each set of transition predicates \mathcal{P} the set of abstract transitions $P^{\#}$ is not disjunctively well-founded.
- (b) Assume you have a tool that does a reachability analysis and you have a tool that computes the TPA algorithm. Describe a termination analysis that uses both tools and can be used to show termination of your program stated in part (a).