## Exercise Sheet 2

Early submission: Monday, 2020-06-08, 14:00 Regular submission: Tuesday, 2020-06-09, 14:00
Exercise 1 - Validating LSC Specifications

(a) LSC 'move down 1'.

(b) Pre-chart TBA of 'move down 1'.

(c) Main-chart TBA of 'move down 1'.

Figure 1: Powerwindow specification.
Consider the LSC requirements specification for a power window in Figure 1a and the TBAs of preand main-chart given in Figures 1b and 1c. For the following computation paths, state whether it trivially satisfies (main-chart not activated), non-trivially satisfies with legal exit or without (in complete or partial traversal, i.e., final state reached or not), or violates the LSC, where the violation can be due to violated condition, violated progress condition, or either message sending or reception missing and argue your claim with the TBAs.
All computation paths end in an infinite sequence of $\tau$-transitions (they 'do not do anything' after the last given transition). We use the following abbreviations: p/position, bot/bottom, up/up_pressed, dp/down_pressed, dr/down_released, md/move_down, br/bottom_reached, st/stop.
(i) $\pi_{1} \equiv \sigma_{0} \xrightarrow{\tau} \underset{(p=b o t)}{\sigma_{1}} \xrightarrow{u p_{!?}^{B, C}} \underset{(p=b o t)}{\sigma_{2}} \xrightarrow{\tau} \cdots$
(ii) $\pi_{2} \equiv \sigma_{0} \xrightarrow{\tau} \underset{\substack{ \\(p \neq b o t)}}{\sigma_{1}} \xrightarrow{\frac{d p_{!?}^{B, C}}{\longrightarrow}} \sigma_{2} \xrightarrow{\tau} \sigma_{3} \xrightarrow{m d_{!?}^{C, M}} \sigma_{4} \xrightarrow{\tau} \cdots$
(iii) $\pi_{3} \equiv \sigma_{0} \xrightarrow{\tau} \underset{(p \neq b o t)}{\sigma_{1}} \xrightarrow{d p_{!?}^{B, C}} \sigma_{2} \xrightarrow{\tau} \sigma_{3} \xrightarrow{m d_{!?}^{C, M}} \sigma_{4} \xrightarrow{b r_{!?}^{S, C}} \sigma_{5} \xrightarrow{\tau} \cdots$
(iv) $\pi_{4} \equiv \sigma_{0} \xrightarrow{\tau} \underset{(p \neq b o t)}{\sigma_{1}} \xrightarrow{d p_{!?}^{B, C}} \sigma_{2} \xrightarrow{\tau} \sigma_{3} \xrightarrow{m d_{!?}^{C, M}} \sigma_{4} \xrightarrow{b r_{!?}^{S, C}} \sigma_{5} \xrightarrow{s t_{!?}^{C, M}} \sigma_{6} \xrightarrow{\tau} \cdots$
(v) $\pi_{5} \equiv \sigma_{0} \xrightarrow{\tau} \underset{(p \neq b o t)}{\sigma_{1}} \xrightarrow{d p_{!!?}^{B, C}} \sigma_{2} \xrightarrow{\tau} \sigma_{3} \xrightarrow{m d_{!?}^{C, M}} \sigma_{4} \xrightarrow{d r_{!?}^{B, C}} \sigma_{5} \xrightarrow{s t_{!?}^{C, M}} \sigma_{6} \xrightarrow{\tau} \cdots$

## Exercise 2 - Validating LSC Specifications Continued



Figure 2: Powerwindow specification continued.

Figure 2 gives a second LSC requirements specification for the same power window.
(i) Does computation path $\pi_{5}$ satisfy or violate (trivial, non-trivial, etc.) Figure 2a? Argue your statement on the TBAs.
(ii) Discuss the following computation path wrt. both LSC.
$\pi_{6} \equiv \sigma_{0} \xrightarrow{\tau} \underset{(p \neq b o t)}{\sigma_{1}} \xrightarrow{d p_{!?}^{B, C}} \sigma_{2} \xrightarrow{\tau} \sigma_{3} \xrightarrow{m d_{!?}^{C, M}} \sigma_{4} \xrightarrow{\text { br } r_{!}^{S, C}, d r_{!?}^{B, C}} \sigma_{5} \xrightarrow{s t_{!?}^{C, M}} \sigma_{6} \xrightarrow{\tau} \cdots$
(iii) Could a power window system on which we can observe computation paths $\pi_{5}$ and $\pi_{6}$ satisfy the overall specification consisting of the LSCs in Figures 1a and 2a? Should it?

## Exercise 3 - LSC Syntax

To investigate the semantics of the LSC shown in Figure 3, we need its abstract syntax.
(i) Provide the set of locations $\mathcal{L}$ including their temperature (make sure to clearly indicate where in the diagram which of your locations occurs).
(ii) Give the partial order relation $\preceq$ on the locations (only direct predecessors/successors need to be given; the full relation $\preceq$ is then the transitive, reflexive closure).
(iii) Give one example each of (a) an element of the simultaneity relation $\sim$, (b) a message from Msg (including temperature), (c) a local invariant from Loclnv (including temperature), and (d) a condition from Cond (including temperature).

Use the formal notations from the lecture; you can indicate the temperature with colors.
Hint: Note that all locations in a coregion obtain the same temperature from the relevant instance line segment adjacent to the whole coregion.


Figure 3: A Live Sequence Chart.

## Exercise 4 - LSC Semantics

(i) Give the cuts (including their temperatures) and corresponding fired sets for the body of the chart in Figure 3.
For at least one hot and cold cut prove why they are hot/cold.
(ii) Compute the Büchi automaton for the body of the chart in Figure 3.
(iii) Provide the strictness conditions $\psi_{\text {strict }}$ for one progress transition.

## Exercise 5 - LSC Acceptance

(i) Consider the following computation path $\pi_{1}$.

$$
\pi_{1} \equiv \sigma_{0} \xrightarrow{A_{1}^{I_{3}, I_{2}} \wedge A_{?}^{I_{3}, I_{2}}} \sigma_{1} \xrightarrow{B_{1}^{I_{2}, I_{3}} \wedge B_{?}^{I_{2}, I_{3}}} \sigma_{2} \xrightarrow{\tau} \sigma_{3} \xrightarrow{\tau} \cdots
$$

where $\sigma_{i} \models c_{1} \wedge \neg c_{2}$ for all $i \geq 0$.
Intuitively, the first event in the path is sending and receiving the message $A$. The second event is sending and receiving the message $B$. After that, no event happens ever again. During the whole computation path the condition $c_{1}$ holds and the condition $c_{2}$ is violated.
Is the path $\pi_{1}$ accepted by the chart or does it violate the chart? If it is accepted, does it take a legal exit or not?
(ii) Give two further computation paths $\pi_{2}$ and $\pi_{3}$ such that each of the following conditions is satisfied by one of the paths $\pi_{1}, \pi_{2}$, or $\pi_{3}$ :

- One of the paths violates the chart.
- One of the paths is accepted and takes a legal exit.
- One of the paths is accepted without taking a legal exit.

Use $\pi_{1}, \pi_{2}$, and $\pi_{3}$ to explain in your own words the intuition behind the concepts of acceptance with and without taking a 'legal exit', and of violation of a chart.

Exercise 6 - Use Cases and Use Case Diagrams
\(\left.$$
\begin{array}{|l|l|}\hline \text { name } & \text { exercise submission (preliminaries) } \\
\hline \text { goal } & \text { exercise sheet solution uploaded } \\
\hline \text { precondition } & \text { student navigated to exercise group screen in ILIAS } \\
\hline \text { postcondition } & \text { student has submitted their solutions } \\
\hline \text { actors } & \text { student (main actor), tutor } \\
\hline \text { open questions } & \text { none } \\
\hline \text { normal case } & \begin{array}{l}\text { 1. student opens the current exercise sheet tab } \\
\text { 2. Ilias shows 'Hand In' button }\end{array}
$$ <br>

\hline exception case 1a student clicks 'Hand In' button\end{array}\right\}\)| submission deadline exceeded |
| :--- |
| 2a.1 Ilias does not show 'Hand In' button due to deadline |
| 2a.2 student sends solutions to tutor via e-mail and explains why |
| the submission is late |

Figure 4: Example Use Case.
(i) Provide the use case diagram for the simple use case in Figure 4.

Hint: As usual, convince your tutor of the correctness of your solution.
(ii) Why is Task (i) only worth one single, lonely exercise point?

