Software Design, Modelling and Analysis in UML

Lecture 02: Semantical Model

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Why (of all things) UML?

- Note: being a modelling languages doesn't mean being graphical (or: being a visual formalism [Harel]).
- For instance, [Kastens and Büning, 2008] also name:
- Terms and Algebras Sets, Relations, Functions
- Propositional and Predicate Logic
- Graphs
- XML Schema, Entity Relation Diagrams, UML Class Diagrams
- Finite Automata, Petri Nets, UML State Machines
- Pro: visual formalisms are found appealing and easier to grasp.
 Yet they are not necessarily easier to write!
- Beware: you may meet people who dislike visual formalisms just for being graphical maybe because it is easier to "trick" people with a meaningless picture than with a meaningless formula.

More serious: it's maybe easier to misunderstand a picture than a formula.

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Contents & Goals

Last Lecture:

- Motivation: model-based development of things (houses, software) to cope with complexity, detect errors early
- Model-based (or -driven) Software Engineering
- UML Mode of the Lecture: Blueprint.

This Lecture:

- Educational Objectives: Capabilities for these tasks/questions:

 Why is UML of the form it is!

 Shall one feel had if not using all diagrams during software development?

 What is a signature, an object, a system state, etc.?

 What's the purpose of signature, object, etc. in the counse?

 How do Basic Object System Signatures relate to UML class diagrams?

- Brief history of UML
 Course map revisited
 Basic Object System Signature, Structure, and System State

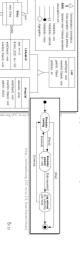
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Why (of all things) UML?

Boxes/lines a
1970's, Soft
Idea: lear A Brief Histo Early 1990's
 Inflation Mid 1980's Languages: Object-Modeling Technique (OMT) [Rumbaugh et al., 1990]
 Booch Method and Notation [Booch, 1993] il., 1990] mming

A Brief History of UML

- Boxes/lines and finite automata are used to visualise software for ages.
- 1970's, Software CrisisTM
 Idea: learn from engineering disciplines to handle growing complexity. Languages: Flowcharts, Nassi-Shneiderman, Entity-Relation Diagrams
- Mid 1980's: Statecharts [Harel, 1987], StateMateTM [Harel et al., 1990]
- Early 1990's, advent of Object-Oriented-Analysis/Design/Programming
 Inflation of notations and methods, most prominent:
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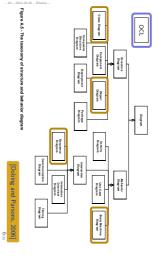
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 Booch Method and Notation [Booch, 1993]
 Object-Oriented Software Engineering (OOSE) [Jacobson et al., 1992]
- Each "persuasion" selling books, tools, seminars...
- Late 1990's: joint effort UML 0.x, 1.x Standards published by **Object Management Group** (OMG), "international, open membership, not-for-profit computer industry consortium".
- Since 2005: UML 2.x

UML Overview (OMG, 2007b, 684)



Course Map Revisited

(i) Common semantical domain.
(ii) UML fragments as syntax.
(iii) Abstract regressmants in of diagrams.
(iv) Informal semantics:
(iv) Informal semantics:
(iv) assign meaning to diagram
(iv) assign meaning to diagram
(iv) Define, e.g., comistency. The Plan | Control | Cont 9/23

Common Expectations on UML

- Easily writeable, readable even by customers
- Powerful enough to bridge the gap between idea and implementation
 Means to tame complexity by separation of concerns ("views")
- Unambiguous
- Standardised, exchangeable between modelling tools
 UML standard says how to develop software
 Using UML leads to better software

We will see...

Seriously: After the course, you should have an own opinion on each of these claims. In how far/in what sense does it hold? Why? Why not? How can it be achieved? Which ones are really only hopes and expectations? \dots ?

UML: Semantic Areas



[OMG, 2007b, 11]

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Common Semantical Domain

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Basic Object System Signature Another Example

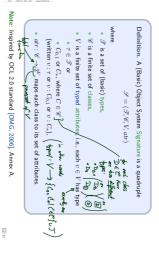
 $\mathcal{S} = (\mathcal{T}, \mathcal{C}, V, atr)$ where

```
\bullet (basic) types {\mathscr T} and classes {\mathscr C}, (both finite),
• atr: \mathscr{C} \to 2^V mapping classes to attributes.
                                                                             • typed attributes V, \tau from \mathscr T or C_{0,1} or C_*, C\in\mathscr C,
```



Q: What about or city with built x: ht
- class Q with afficient x: Affind?
- class Q with afficient x: Affind?
A: Reviewe consistently.

Basic Object System Signature

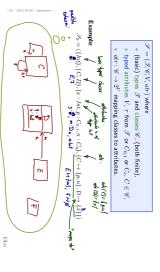


Basic Object System Structure

```
We use \mathscr{D}(\mathscr{C}) to denote \bigcup_{C\in\mathscr{C}}\mathscr{D}(C); analogously \mathscr{D}(\mathscr{C}_*).
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                      Definition. A Basic Object System Structure of
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                             is a domain function {\mathscr D} which assigns to each type a domain, i.e.
                                                                                                                                                                                         • C \in \mathscr{C} is mapped to an infinite set \mathscr{D}(C) of (\circ b)ect) identities. Note: Object identities only have the "x" operation; object identities only have the "x" operation; object identities of different classes are disjoint, i.e. \forall C, D \in \mathscr{C}: C \neq D \to \mathscr{D}(C) \cap \mathscr{D}(D) = \emptyset.
                                                                                                    • C_* and C_{0,1} for C \in \mathscr{C} are mapped to 2^{\mathscr{D}(C)}.
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                       • \tau \in \mathscr{T} is mapped to \mathscr{D}(\tau).
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                         \mathscr{S} = (\mathscr{T}, \mathscr{C}, V, atr)
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Note: We identify objects and object identities, because both uniquely determine each other (cf. OCL 2.0 standard).

Basic Object System Signature Example

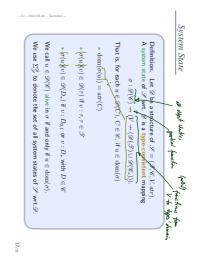


Basic Object System Structure Example

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Wanted: a structure for signature
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\mathcal{S}_0 = (\{Int\}, \{C, D\}, \{x : Int, p : C_{0,1}, n : C_*\}, \{C \mapsto \{p, n\}, D \mapsto \{x\}\}))
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Recall: by definition, seek a @ which maps
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                          • \tau \in \mathscr{T} to some \mathscr{D}(\tau),
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                             C_* and C_{0,1} for C\in \mathscr{C} to \mathscr{D}(C_{0,1})=\mathscr{D}(C_*)=2^{\mathscr{D}(C)}.
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                            c\in\mathscr{C} to some identities \mathscr{D}(C) (infinite, disjoint for different classes),
\mathcal{D}(C_{0,1}) = \mathcal{D}(C_*) = 2 \frac{\mathcal{D}(\mathcal{C})}{\mathcal{D}(D_{0,1})}
\mathcal{D}(D_{0,1}) = \mathcal{D}(D_*) = 2 \frac{\mathcal{D}(D_*)}{\mathcal{D}(D_*)}
                                                                                                                                                                                                          \mathscr{D}(D) = \mathbb{N}^{+} \times \{\mathfrak{D}\} \simeq \{\eta_{\mathbf{b}}, 2\mathfrak{d}, 3\mathfrak{d}, ...\} / {}^{*} \{ \mathfrak{t}, \mathfrak{e}, ... \}
                                                                                                                                                                                                                                                                                                                                                                                    \mathscr{D}(Int) = \mathbf{Z}
                                                                                                                                                                                                                                                                                               \mathscr{D}(C) = \mathbb{N}^{4} \times \{\mathcal{C}\} \simeq \{\mathcal{C}_{\epsilon}, 2_{\epsilon}, 3_{\epsilon}, ...\} \quad \left| z \in \mathcal{I}, 3, S, ... \right\}
                                                                                                                                                                                                                                                                                                                                                                                                                    also rectid: Do
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You Are Here.

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System State Example

System State Example Signature, Structure:

Wanted: $\sigma: \mathscr{D}(\mathscr{C}) \nrightarrow (V \nrightarrow (\mathscr{D}(\mathscr{T}) \cup \mathscr{D}(\mathscr{C}_*)))$ such that

 $\mathscr{S}_{0} = (\{Int\}, \{C, D\}, \{x: Int, p: C_{0,1}, n: C_{*}\}, \{C \mapsto \{p, n\}, D \mapsto \{x\}\})$ $\mathscr{D}(Int) = \mathbb{Z}, \quad \mathscr{D}(C) = \{1_C, 2_C, 3_C, \ldots\}, \quad \mathscr{D}(D) = \{1_D, 2_D, 3_D, \ldots\}$

$$\begin{split} & \bullet \ \mathrm{dom}(\sigma(u)) = atr(C), \\ & \bullet \ \sigma(u)(v) \in \mathscr{D}(\tau) \ \text{if} \ v : \tau, \tau \in \mathscr{T}, \\ & \bullet \ \sigma(u)(v) \in \mathscr{D}(C_*) \ \text{if} \ v : D_* \ \text{with} \ D \in \mathscr{C} \ . \end{split}$$

Signature, Structure:

 $\mathscr{S}_0 = (\{\mathit{Int}\}, \{C, D\}, \{x : \mathit{Int}, p : C_{0,1}, n : C_{\bullet}\}, \{C \mapsto \{p, n\}, D \mapsto \{x\}\})$ $\mathscr{D}(Int) = \mathbb{Z}, \quad \mathscr{D}(C) = \{1_C, 2_C, 3_C, \ldots\}, \quad \mathscr{D}(D) = \{1_D, 2_D, 3_D, \ldots\}$



Alternative: symbolic system state

assuming $c_1, c_2 \in \mathscr{D}(C), d \in \mathscr{D}(D), c_1 \neq c_2.$

 $\sigma = \{c_1 \mapsto \{p \mapsto \emptyset, n \mapsto \{c_2\}\}, c_2 \mapsto \{p \mapsto \emptyset, n \mapsto \emptyset\}, d \mapsto \{x \mapsto 23\}\}$

 $\sigma = \{1_C \mapsto \{p \mapsto \emptyset, n \mapsto \{5_C\}\}, 5_C \mapsto \{\underline{p} \mapsto \emptyset, \underline{n} \mapsto \emptyset\}, 1_D \mapsto \{x \mapsto 23\}\}.$

dan (σ(Sz))= {ρ,+}= a+(e) ν

Concrete, explicit:

Course Map $\varphi \in \mathsf{OCL}$ G = (N, E, f)OD Mathematics $w_{\pi} = ((\sigma_i, cons_i, Snd_i))_{i \in \mathbb{N}}$

References

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References

[Booch, 1903] Booch, G. (1903). Object-oriented Analysis and Design with Applications.

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