

# *Software Design, Modelling and Analysis in UML*

## *Lecture 02: Semantical Model*

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# Contents & Goals

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## Last Lecture:

- Motivation: model-based development of things (houses, software) to cope with complexity, detect errors early
- Model-based (or -driven) Software Engineering
- UML Mode of the Lecture: Blueprint.

## This Lecture:

- **Educational Objectives:** Capabilities for these tasks/questions:
  - Why is UML of the form it is?
  - Shall one feel bad if not using all diagrams during software development?
  - What is a signature, an object, a system state, etc.?  
What's the purpose of signature, object, etc. in the course?
  - How do Basic Object System Signatures relate to UML class diagrams?
- **Content:**
  - Brief history of UML
  - Course map revisited
  - Basic Object System Signature, Structure, and System State

*Why (of all things) UML?*

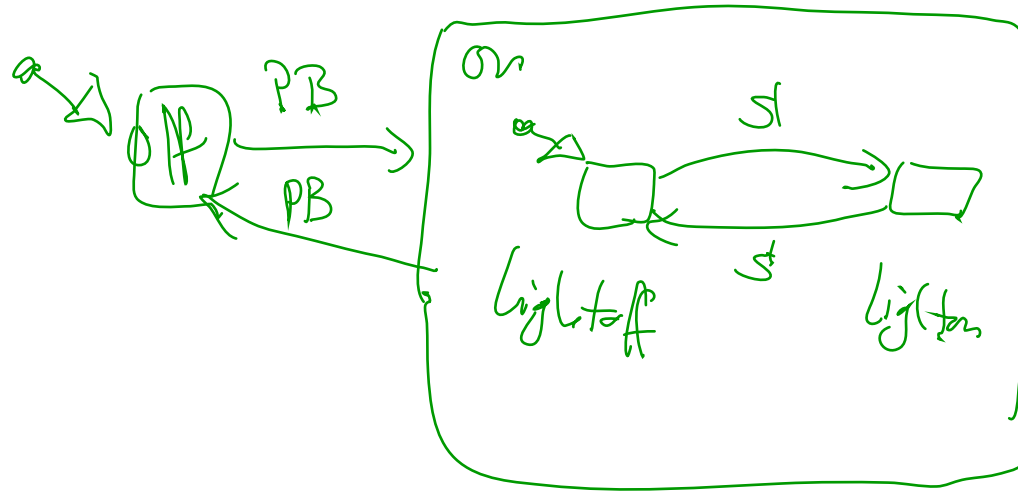
# *Why (of all things) UML?*

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- Note: being a **modelling** languages doesn't mean being graphical (or: being a visual formalism [Harel]).
- For instance, [Kastens and Büning, 2008] also name:
  - Sets, Relations, Functions
  - Terms and Algebras
  - Propositional and Predicate Logic
  - Graphs
  - XML Schema, Entity Relation Diagrams, UML Class Diagrams
  - Finite Automata, Petri Nets, UML State Machines
- **Pro**: visual formalisms are found appealing and easier to **grasp**.  
Yet they are not necessarily easier to **write**!
- **Beware**: you may meet people who dislike visual formalisms just for being graphical — maybe because it is easier to “trick” people with a meaningless picture than with a meaningless formula.  
More serious: it's maybe easier to misunderstand a picture than a formula.

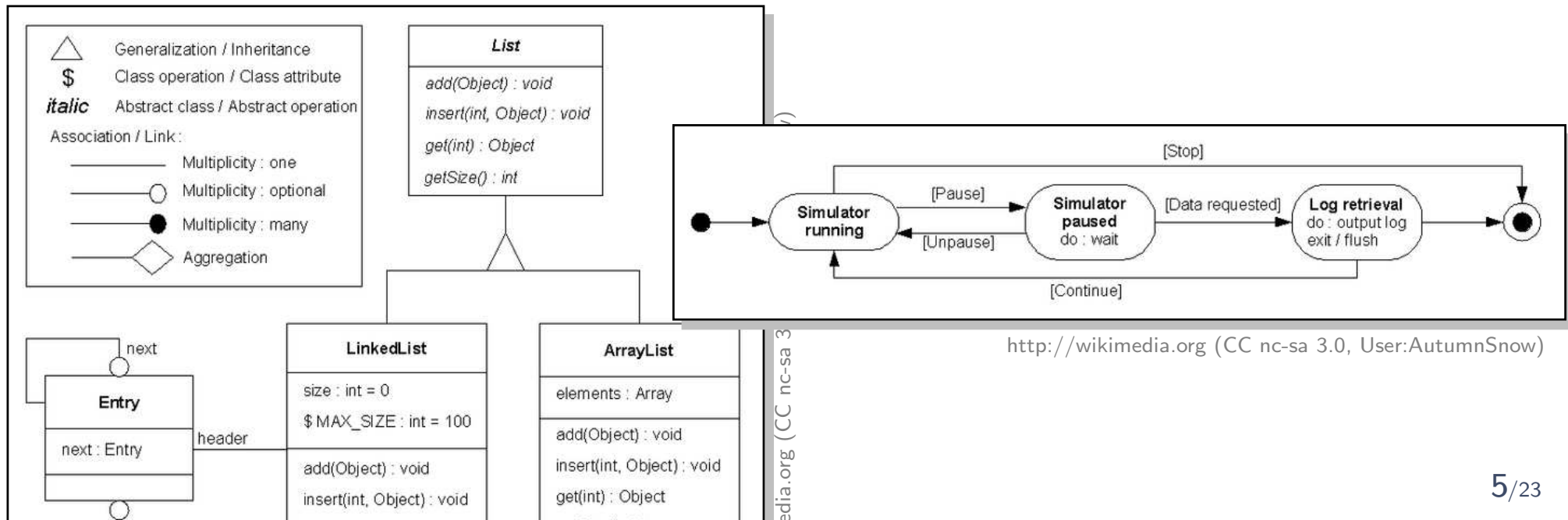
# A Brief History of UML

- Boxes/lines and finite automata are used to visualise software **for ages**.
- **1970's, Software Crisis**<sup>TM</sup>
  - Idea: learn from engineering disciplines to handle growing complexity.
  - Languages: **Flowcharts, Nassi-Shneiderman, Entity-Relation Diagrams**
- Mid **1980's**: **Statecharts** [Harel, 1987], **StateMate**<sup>TM</sup> [Harel et al., 1990]



# A Brief History of UML

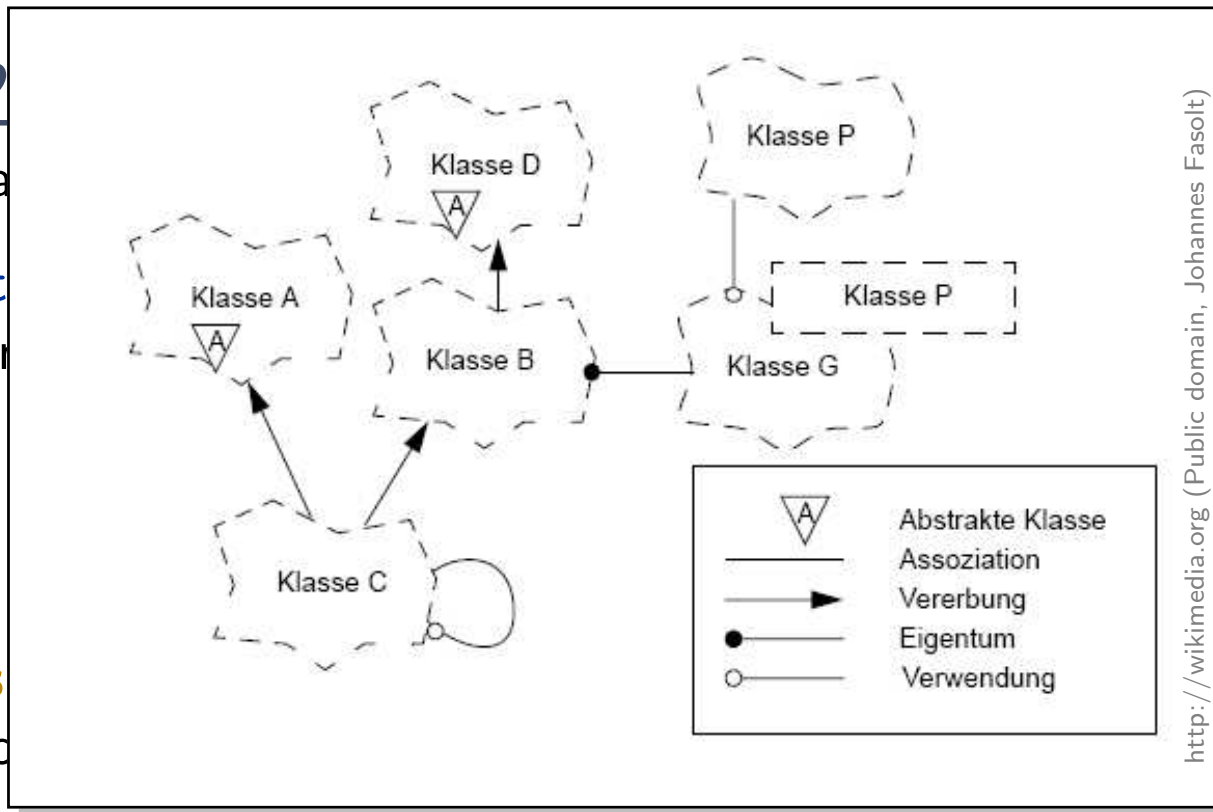
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- Mid **1980's**: **Statecharts** [Harel, 1987], **StateMate**<sup>TM</sup> [Harel et al., 1990]
- Early **1990's**, advent of **Object-Oriented**-Analysis/Design/Programming
  - Inflation of notations and methods, most prominent:
    - **Object-Modeling Technique** (OMT) [Rumbaugh et al., 1990]



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# A Brief History

- Boxes/lines and arrows
- **1970's, Software Engineering**
  - Idea: learn from programming languages
- Mid **1980's**:
  - Languages: Smalltalk, C++, Pascal
- Early **1990's**
  - Inflation of notations



- **Object-Modeling Technique (OMT)** [Rumbaugh et al., 1990]
- **Booch Method and Notation** [Booch, 1993]

Images.

Complexity.

Games

al., 1990]

Programming

# A Brief History of UML

- Boxes/lines and finite automata are used to visualise software **for ages**.
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    - **Booch Method and Notation** [Booch, 1993]
    - **Object-Oriented Software Engineering** (OOSE) [Jacobson et al., 1992]

Each “persuasion” selling books, tools, seminars. . .

- Late **1990's**: joint effort **UML 0.x, 1.x**
  - Standards published by **Object Management Group** (OMG), “*international, open membership, not-for-profit computer industry consortium*”.
- Since **2005**: **UML 2.x**



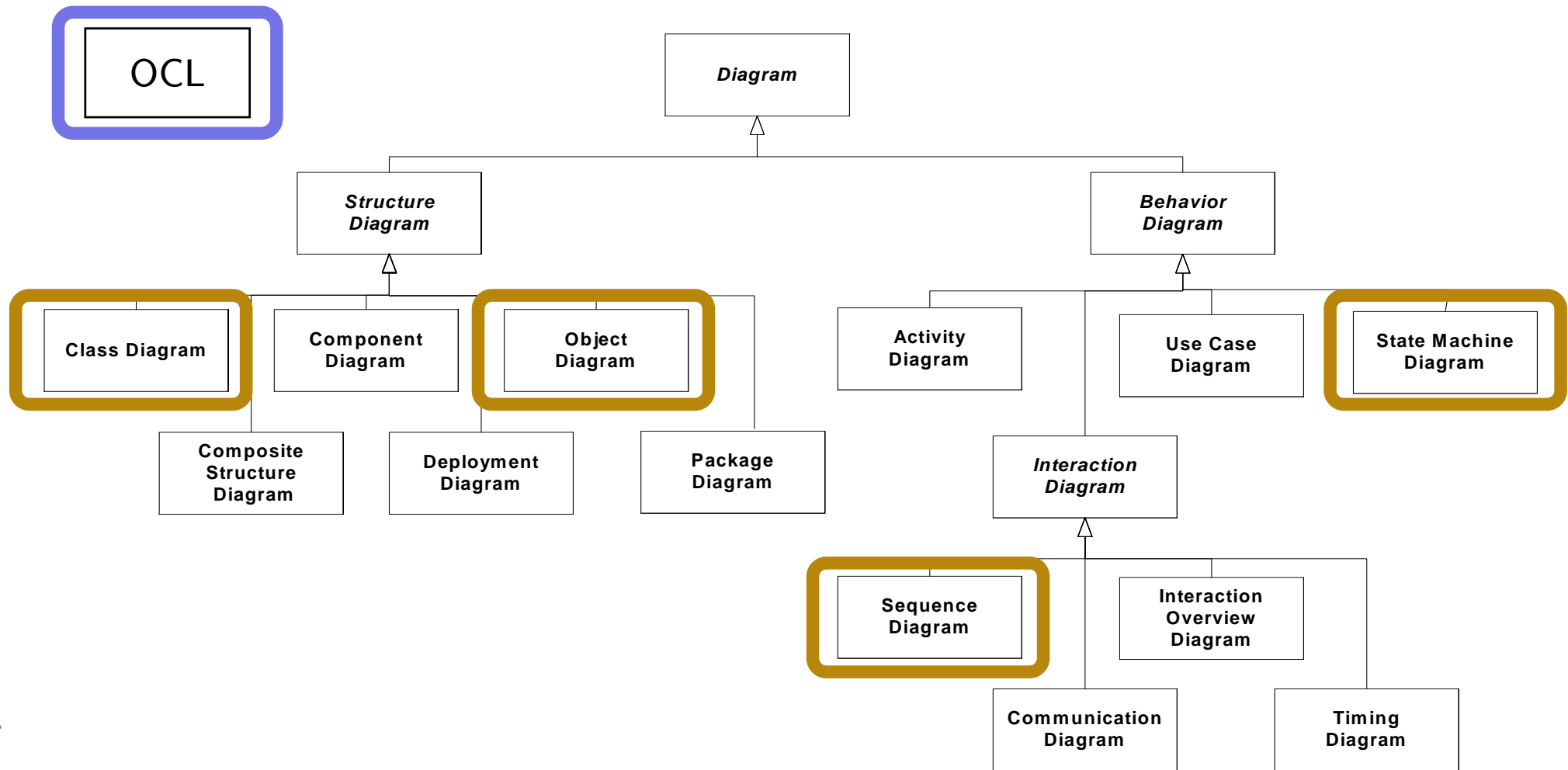


Figure A.5 - The taxonomy of structure and behavior diagram

[Dobing and Parsons, 2006]

# *Common Expectations on UML*

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- Easily writeable, readable even by customers
- Powerful enough to bridge the gap between idea and implementation
- Means to tame complexity by separation of concerns ( “views” )
- Unambiguous
- Standardised, exchangeable between modelling tools
- UML standard says how to develop software
- Using UML leads to better software
- . . .

## We will see...

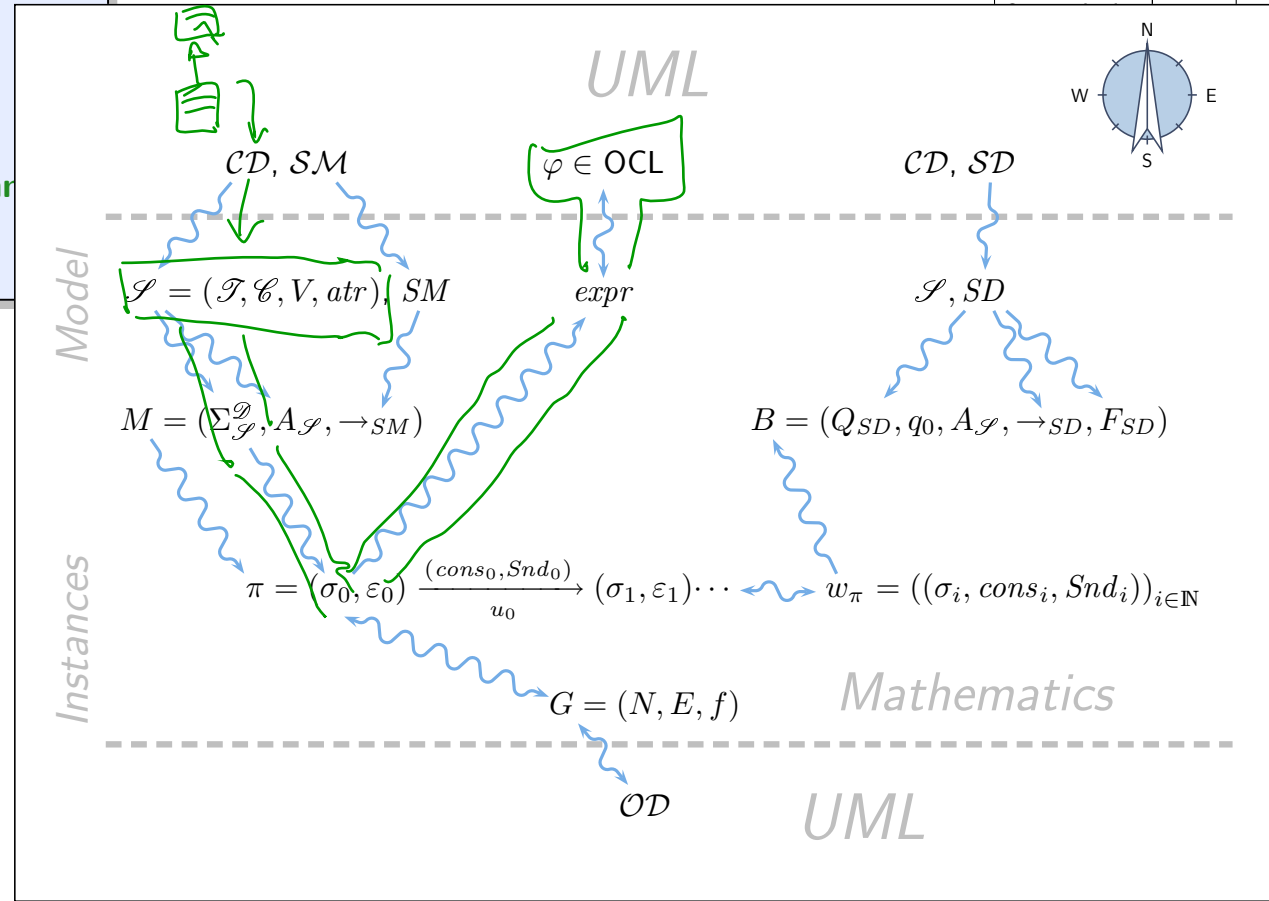
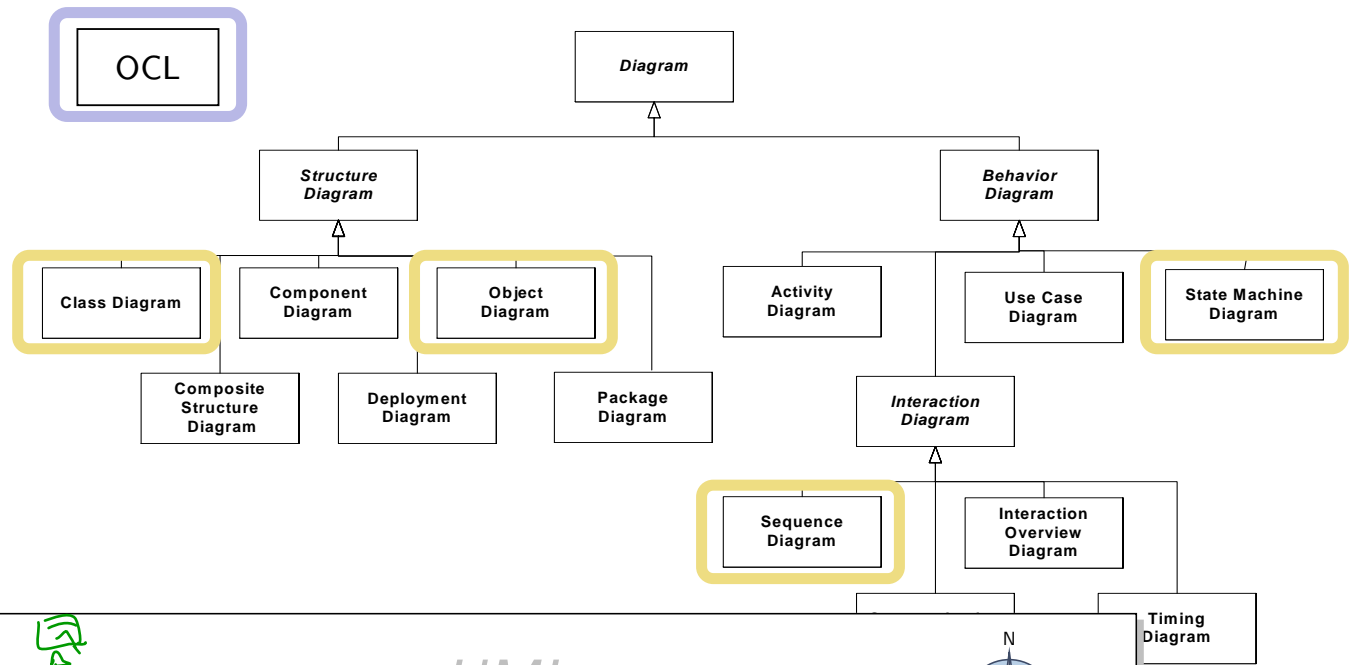
Seriously: After the course, you should have an own opinion on each of these claims. In how far/in what sense does it hold? Why? Why not? How can it be achieved? Which ones are really only hopes and expectations? . . . ?

# *Course Map Revisited*

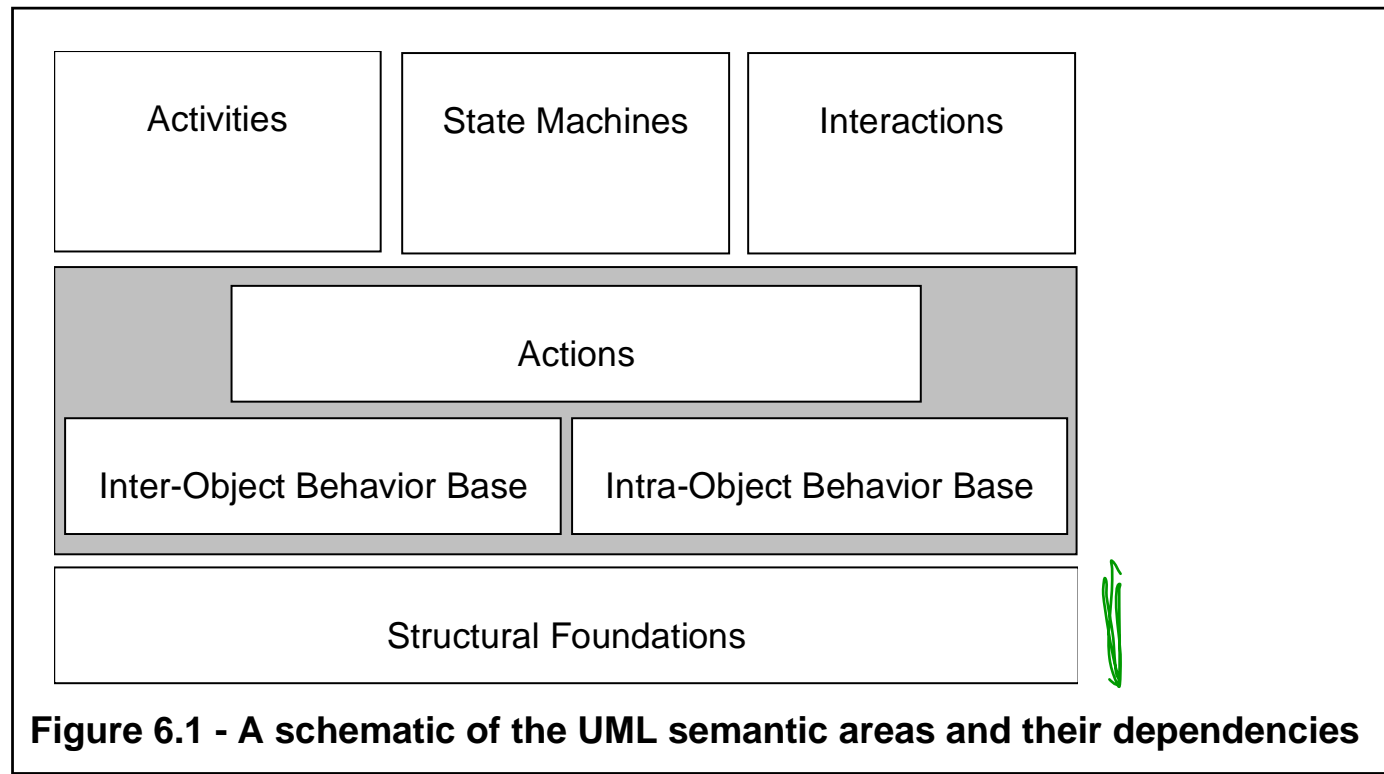
# The Plan

Recall:

- **Overall aim:** a formal language for software blueprints.
- **Approach:**
  - Common semantical domain.
  - UML fragments as **syntax**.
  - Abstract **representation of diagrams**.
  - Informal semantics:** UML standard
  - assign meaning to diagram**
  - Define, e.g., **consistency**.



# UML: Semantic Areas



[OMG, 2007b, 11]

# *Common Semantical Domain*

# Basic Object System Signature

**Definition.** A (Basic) Object System **Signature** is a quadruple

$$\mathcal{S} = (\mathcal{T}, \mathcal{C}, V, \text{atr})$$

where

- $\mathcal{T}$  is a set of (basic) **types**,
- $\mathcal{C}$  is a finite set of **classes**,
- $V$  is a finite set of **typed attributes**, i.e., each  $v \in V$  has type
  - $\tau \in \mathcal{T}$  or
  - $C_{0,1}$  or  $C_*$ , where  $C \in \mathcal{C}$
 (written  $v : \tau$  or  $v : C_{0,1}$  or  $v : C_*$ ),
- $\text{atr} : \mathcal{C} \rightarrow 2^V$  maps each class to its set of attributes.

for each class  $D \in \mathcal{C}$   
there are two different  
types:

$D_{0,1}$       $D_\Delta$

$D_*$       $D_{\odot}$

$\alpha:$   


$\alpha:$   


in other words:

$$\text{typeof} : V \rightarrow \mathcal{T} \cup \{D_{0,1}, D_*, D_{\odot} \mid D \in \mathcal{C}\}$$

total function

powerset of  $V$

**Note:** Inspired by OCL 2.0 standard [OMG, 2006], Annex A.

# Basic Object System Signature Example

$\mathcal{S} = (\mathcal{T}, \mathcal{C}, V, atr)$  where

- (basic) **types**  $\mathcal{T}$  and **classes**  $\mathcal{C}$ , (both finite),
- **typed attributes**  $V$ ,  $\tau$  from  $\mathcal{T}$  or  $C_{0,1}$  or  $C_*$ ,  $C \in \mathcal{C}$ ,
- $atr : \mathcal{C} \rightarrow 2^V$  mapping classes to attributes.

**Example:**

$\mathcal{S}_0 = (\{\text{Int}\}, \{C, D\}, \{x : \text{Int}, p : C_{0,1}, n : C_*\}, \{C \mapsto \{p, n\}, D \mapsto \{x\}\})$

*Handwritten annotations:*  
 - "basic types" points to  $\{\text{Int}\}$   
 - "classes" points to  $\{C, D\}$   
 - "attributes" points to  $\{x : \text{Int}, p : C_{0,1}, n : C_*\}$   
 - "attribute x has type Int" points to  $x : \text{Int}$   
 - "atr" points to the mapping set  $\{C \mapsto \{p, n\}, D \mapsto \{x\}\}$   
 - "atr(C) = {p, n}" points to  $C \mapsto \{p, n\}$   
 - "atr(D) = {x}" points to  $D \mapsto \{x\}$   
 - "maps to" points to the mapping set.



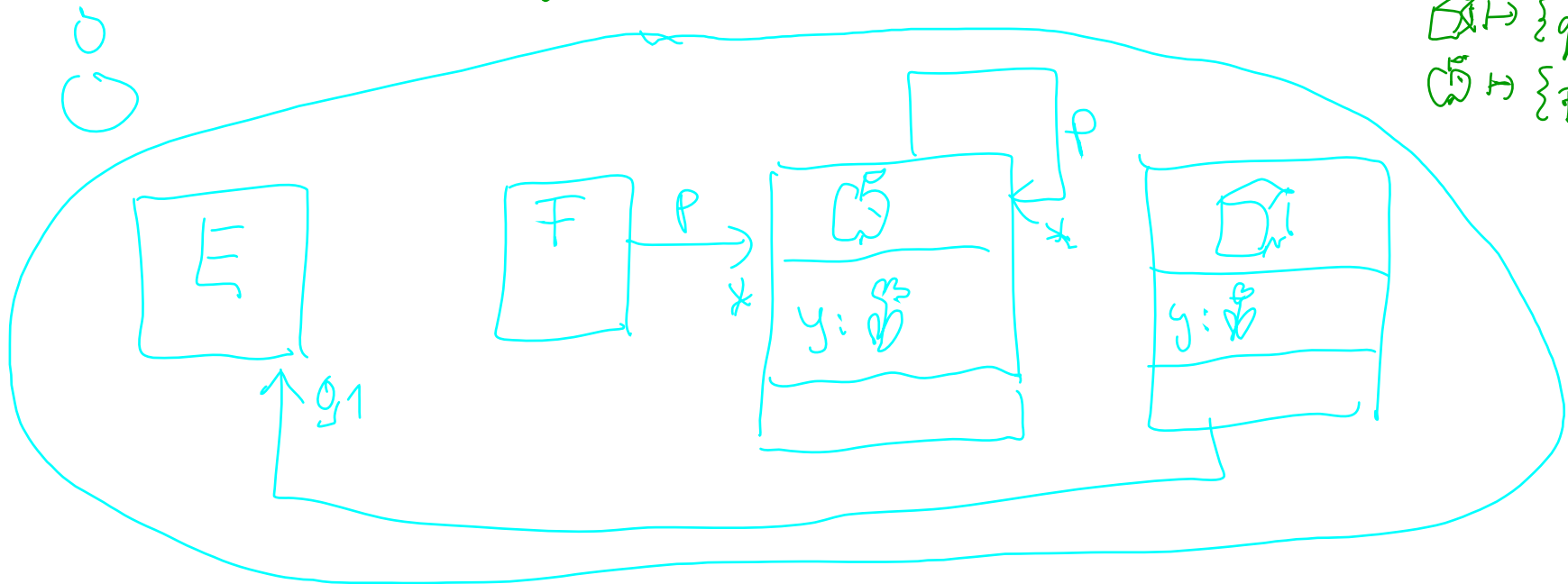
# Basic Object System Signature Another Example

$\mathcal{S} = (\mathcal{T}, \mathcal{C}, V, atr)$  where

- (basic) **types**  $\mathcal{T}$  and **classes**  $\mathcal{C}$ , (both finite),
- **typed attributes**  $V$ ,  $\tau$  from  $\mathcal{T}$  or  $C_{0,1}$  or  $C_*$ ,  $C \in \mathcal{C}$ ,
- $atr : \mathcal{C} \rightarrow 2^V$  mapping classes to attributes.

**Example:**

$$\mathcal{S}_n = (\{\emptyset\}, \{E, F, \boxed{\square}\}, \{y: \emptyset, p: \boxed{\square}_*, q: E_{0,1}\}, \{E \mapsto \emptyset, F \mapsto \{p\}, \boxed{\square} \mapsto \{q, y\}, \emptyset \mapsto \{p, y\}\})$$



**Definition.** A Basic Object System **Structure** of

$$\mathcal{S} = (\mathcal{T}, \mathcal{C}, V, atr)$$

is a domain function  $\mathcal{D}$  which assigns to each type a domain, i.e.

- $\tau \in \mathcal{T}$  is mapped to  $\mathcal{D}(\tau)$ ,
- $C \in \mathcal{C}$  is mapped to an infinite set  $\mathcal{D}(C)$  of **(object) identities**.

Note: Object identities only have the “=” operation;  
object identities of different classes are disjoint, i.e.

$$\forall C, D \in \mathcal{C} : C \neq D \rightarrow \mathcal{D}(C) \cap \mathcal{D}(D) = \emptyset.$$

- $C_*$  **and**  $C_{0,1}$  for  $C \in \mathcal{C}$  are mapped to  $2^{\mathcal{D}(C)}$ .

We use  $\mathcal{D}(\mathcal{C})$  to denote  $\bigcup_{C \in \mathcal{C}} \mathcal{D}(C)$ ; analogously  $\mathcal{D}(\mathcal{C}_*)$ .

**Note:** We identify objects and object identities, because both uniquely determine each other (cf. OCL 2.0 standard).

# Basic Object System Structure Example

**Wanted:** a structure for signature

$$\mathcal{S}_0 = (\{Int\}, \{C, D\}, \{x : Int, p : C_{0,1}, n : C_*\}, \{C \mapsto \{p, n\}, D \mapsto \{x\}\})$$

Recall: by definition, seek a  $\mathcal{D}$  which maps

- $\tau \in \mathcal{T}$  to **some**  $\mathcal{D}(\tau)$ ,
- $c \in \mathcal{C}$  to **some** identities  $\mathcal{D}(C)$  (infinite, disjoint for different classes),
- $C_*$  and  $C_{0,1}$  for  $C \in \mathcal{C}$  to  $\mathcal{D}(C_{0,1}) = \mathcal{D}(C_*) = 2^{\mathcal{D}(C)}$ .

alternative  
choice:

$$\begin{aligned} \mathcal{D}(Int) &= \mathbb{Z} \\ \mathcal{D}(C) &= \mathbb{N}^+ \times \{C\} = \{1_C, 2_C, 3_C, \dots\} \\ \mathcal{D}(D) &= \mathbb{N}^+ \times \{D\} = \{1_D, 2_D, 3_D, \dots\} \\ \mathcal{D}(C_{0,1}) = \mathcal{D}(C_*) &= 2^{\mathbb{N}^+ \times \{C\}} \\ \mathcal{D}(D_{0,1}) = \mathcal{D}(D_*) &= 2^{\mathbb{N}^+ \times \{D\}} \end{aligned}$$

$$\begin{aligned} \mathcal{D}_2(Int) &= \{-127, \dots, 127\} \\ \mathcal{D}_2(C) &= \{1, 3, 5, \dots\} \\ \mathcal{D}_2(D) &= \{2, 4, 6, \dots\} \\ 2^{\mathcal{D}_2(C)} & \\ 2^{\mathcal{D}_2(D)} & \end{aligned}$$

**Definition.** Let  $\mathcal{D}$  be a structure of  $\mathcal{S} = (\mathcal{I}, \mathcal{C}, V, atr)$ .  
A **system state** of  $\mathcal{S}$  wrt.  $\mathcal{D}$  is a **type-consistent** mapping

$$\sigma : \mathcal{D}(\mathcal{C}) \rightarrow (V \rightarrow (\mathcal{D}(\mathcal{I}) \cup \mathcal{D}(\mathcal{C}_*))).$$

That is, for each  $u \in \mathcal{D}(\mathcal{C})$ ,  $C \in \mathcal{C}$ , if  $u \in \text{dom}(\sigma)$

- $\text{dom}(\sigma(u)) = atr(C)$  ||
- $(\sigma(u))(v) \in \mathcal{D}(\tau)$  if  $v : \tau, \tau \in \mathcal{I}$
- $(\sigma(u))(v) \in \mathcal{D}(D_*)$  if  $v : D_{0,1}$  or  $v : D_*$  with  $D \in \mathcal{C}$

We call  $u \in \mathcal{D}(\mathcal{C})$  **alive** in  $\sigma$  if and only if  $u \in \text{dom}(\sigma)$ .

We use  $\Sigma_{\mathcal{S}}^{\mathcal{D}}$  to denote the set of all system states of  $\mathcal{S}$  wrt.  $\mathcal{D}$ .

all object identities

partial function

partial function from  $V$  to types' domains

# System State Example

## Signature, Structure:

$$\mathcal{S}_0 = (\{Int\}, \{C, D\}, \{x : Int, p : C_{0,1}, n : C_*\}, \{C \mapsto \{p, n\}, D \mapsto \{x\}\})$$

$$\mathcal{D}(Int) = \mathbb{Z}, \quad \mathcal{D}(C) = \{1_C, 2_C, 3_C, \dots\}, \quad \mathcal{D}(D) = \{1_D, 2_D, 3_D, \dots\}$$

**Wanted:**  $\sigma : \mathcal{D}(\mathcal{C}) \rightarrow (V \rightarrow (\mathcal{D}(\mathcal{T}) \cup \mathcal{D}(\mathcal{C}_*)))$  such that

- $\text{dom}(\sigma(u)) = \text{atr}(C)$ ,
- $\sigma(u)(v) \in \mathcal{D}(\tau)$  if  $v : \tau, \tau \in \mathcal{T}$ ,
- $\sigma(u)(v) \in \mathcal{D}(C_*)$  if  $v : D_*$  with  $D \in \mathcal{C}$ .

•  $\sigma_1 = \emptyset$  ← empty function

$$\bullet \sigma_2 = \{ 1_C \mapsto \{ p \mapsto \{ 1_C \}, n \mapsto \{ 5_C, 6_C \} \}, \\ 2_D \mapsto \{ x \mapsto 3 \} \}$$

Wrt.  $\mathcal{D}_2$ :

$$\bullet \sigma_3 = \{ 5 \mapsto \{ p \mapsto \{ 1_C \}, n \mapsto \emptyset \} \}$$

one way to read out:

= object  $1_C$  has a  
p-link to  $1_C$   
(i.e. to itself)

= object  $1_C$  refers to  
objects  $5_C, 6_C$   
via link  $n$

# System State Example

## Signature, Structure:

$$\mathcal{S}_0 = (\{Int\}, \{C, D\}, \{x : Int, p : C_{0,1}, n : C_*\}, \{C \mapsto \{p, n\}, D \mapsto \{x\}\})$$

$$\mathcal{D}(Int) = \mathbb{Z}, \quad \mathcal{D}(C) = \{1_C, 2_C, 3_C, \dots\}, \quad \mathcal{D}(D) = \{1_D, 2_D, 3_D, \dots\}$$

**Wanted:**  $\sigma : \mathcal{D}(\mathcal{C}) \rightarrow (V \rightarrow (\mathcal{D}(\mathcal{T}) \cup \mathcal{D}(\mathcal{C}_*)))$  such that

- $\text{dom}(\sigma(u)) = \text{atr}(C)$ ,
- $\sigma(u)(v) \in \mathcal{D}(\tau)$  if  $v : \tau, \tau \in \mathcal{T}$ ,
- $\sigma(u)(v) \in \mathcal{D}(C_*)$  if  $v : D_*$  with  $D \in \mathcal{C}$ .

### Concrete, explicit:

$$\sigma = \{1_C \mapsto \{p \mapsto \emptyset, n \mapsto \{5_C\}\}, 5_C \mapsto \{p \mapsto \emptyset, n \mapsto \emptyset\}, 1_D \mapsto \{x \mapsto 23\}\}.$$

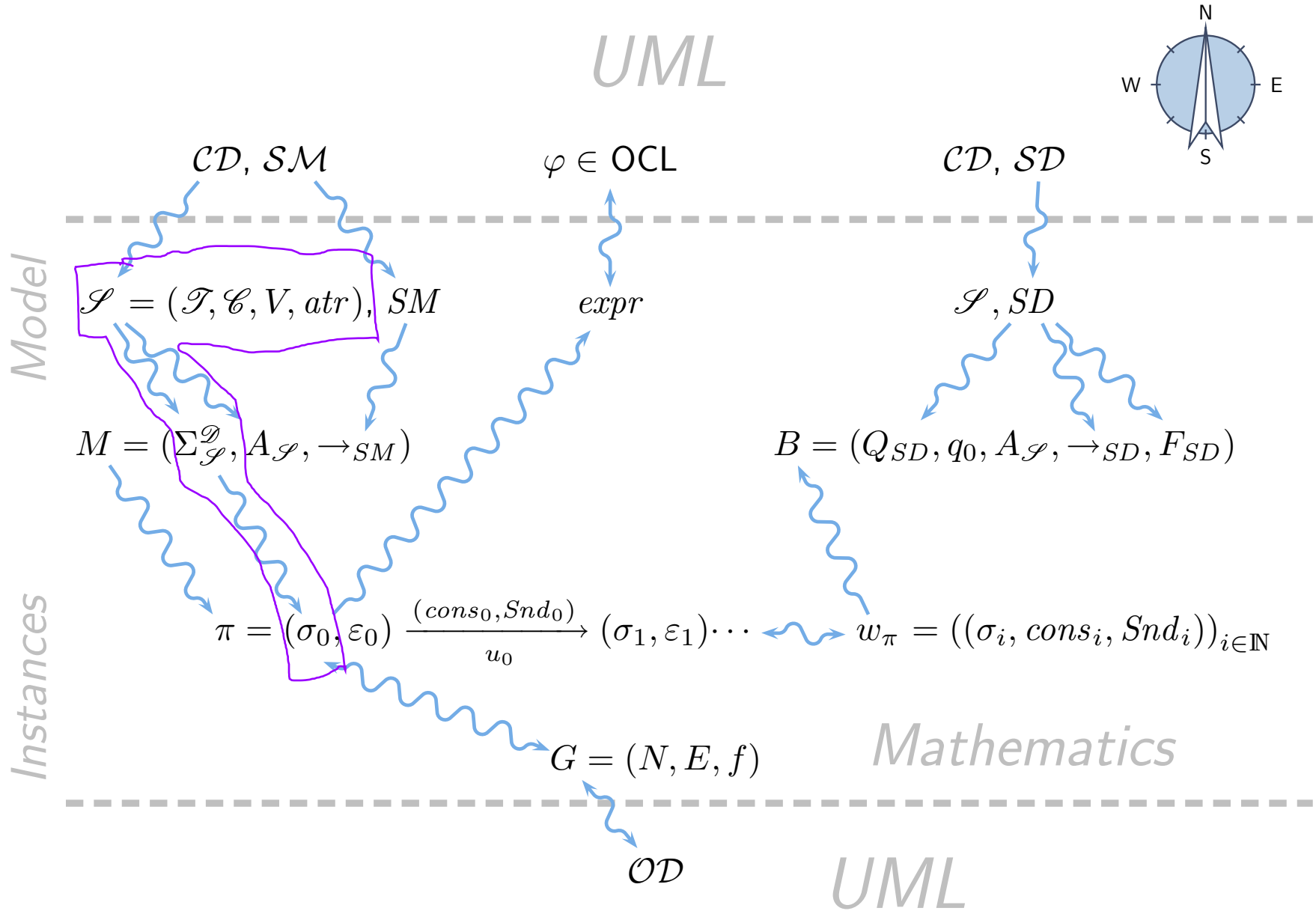
### Alternative: symbolic system state

$$\sigma = \{c_1 \mapsto \{p \mapsto \emptyset, n \mapsto \{c_2\}\}, c_2 \mapsto \{p \mapsto \emptyset, n \mapsto \emptyset\}, d \mapsto \{x \mapsto 23\}\}$$

assuming  $c_1, c_2 \in \mathcal{D}(C)$ ,  $d \in \mathcal{D}(D)$ ,  $c_1 \neq c_2$ .

*You Are Here.*

# Course Map





# *References*

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