Software Design, Modelling and Analysis in UML

Lecture 02: Semantical Model

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Why (of all things) UML?

A Brief History of UML

Boxes/lines and finite automata are used to visualise software for ages.

1970's, Software CrisisTM
 Idea: learn from engineering disciplines to handle growing complexity.

 Mid 1980's: Statecharts [Harel, 1987], StateMateTM [Harel et al., 1990] Languages: Flowcharts, Nassi-Shneiderman, Entity-Relation Diagrams

- Note: being a modelling languages doesn't mean being graphical (or: being a visual formalism [Harel]).
- For instance, [Kastens and Büning, 2008] also name:
- Terms and Algebras Sets, Relations, Functions
- Propositional and Predicate Logic
- Graphs
- XML Schema, Entity Relation Diagrams, UML Class Diagrams
 Finite Automata, Petri Nets, UML State Machines
- Pro: visual formalisms are found appealing and easier to grasp.
 Yet they are not necessarily easier to write!

Beware: you may meet people who dislike visual formalisms just for being graphical — maybe because it is easier to "trick" people with a meaningless picture than with a meaningless formula.

More serious: it's maybe easier to misunderstand a picture than a formula.

Contents & Goals

Last Lecture:

- Motivation: model-based development of things (houses, software) to cope with complexity, detect errors early
- Model-based (or -driven) Software Engineering

This Lecture: UML Mode of the Lecture: Blueprint.

- Educational Objectives: Capabilities for these tasks/questions:
 Why is UML of the form it is!
 Shall one feel had if not using all diagrams during software development?
 What is a signature, an object, a system state, etc.?
 What's the purpose of signature, object, etc. in the counse?
 How do Basic Object System Signatures relate to UML class diagrams?

- Brief history of UML
 Course map revisited
 Basic Object System Signature, Structure, and System State

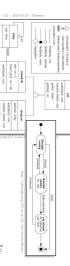
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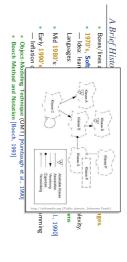
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 Inflation of notations and methods, most prominent:
- Object-Modeling Technique (OMT) [Rumbaugh et al., 1990]



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Common Expectations on UML

- Easily writeable, readable even by customers
- Powerful enough to bridge the gap between idea and implementation
- Means to tame complexity by separation of concerns ("views")

Course Map Revisited

- Unambiguous
- UML standard says how to develop software
 Using UML leads to better software Standardised, exchangeable between modelling tools

We will see...

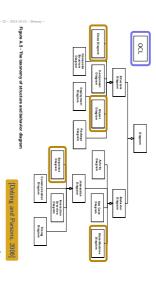
Seriously: After the course, you should have an own opinion on each of these claims. In how far/in what sense does it hold? Why? Why not? How can it be achieved? Which ones are really only hopes and expectations? ...?

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 Booch Method and Notation [Booch, 1993]
 Object-Oriented Software Engineering (OOSE) [Jacobson et al., 1992]
- Each "persuasion" selling books, tools, seminars...
- Late 1990's: joint effort UML 0.x, 1.x
- Standards published by **Object Management Group** (OMG), "international, open membership, not-for-profit computer industry consortium".
- Since 2005: UML 2.x

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UML Overview (OMG, 2007b, 684)



Approach:
(i) Common sensatical domain.
(ii) UMA fragments as syntax.
(iii) UMA fragments as syntax.
(iii) Matricat representation of diagrams.
(iv) Matricat sensantics:
(iv) Mat. standard
(iv) assign measuring to diagram (iv) Define, e.g., comistency. The Plan Rayer Engance Countries Countries Countries Countries 9/23

UML: Semantic Areas



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Common Semantical Domain

Basic Object System Signature

Definition. A (Basic) Object System Signa-

re is a quadruple

 $\mathcal{S} = (\mathcal{T}, \mathcal{C}, V, \mathit{atr})$

 $m{\mathscr{T}}$ is a set of (basic) types,

there are two different types: or or or Ø 8

 V is a finite set of typed attribu % is a finite set of classes, • $\tau \in \mathscr{T}$ or

y/i.e., each $v \in V$ has type $D_{\mathbf{y}} = D_{\mathbf{y}}$

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Hotal function parest of V Note: Inspired by OCL 2.0 standard [OMG, 2006], Annex A.

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• $C_{0,1}$ or C_{∞} , where $C \in \mathscr{C}_{1}$ in this ariself: (written $v : \tau$ or $v : C_{0,1}$ or $v : C_{\infty}$).

• $atr : \mathscr{C}_{M-2} \overset{\mathcal{D}'}{\longrightarrow} \text{maps} \text{ each class to its set of attributes.}$

Basic Object System Signature Example

```
\mathcal{S} = (\mathcal{T}, \mathcal{C}, V, atr) where
                                                                                                                 \bullet typed attributes V,\,\tau from \mathcal F or C_{0,1} or C_*,\,C\in\mathcal E,
                                                                                                                                       ullet (basic) types {\mathscr T} and classes {\mathscr C}, (both finite),
                                                , atr(C) = Epm3 atr(D)=Ex}
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Basic Object System Signature Another Example

```
\mathcal{S} = (\mathcal{T}, \mathcal{C}, V, atr) where
                                                                                                                                                                                                                                                                       • atr: \mathcal{C} \rightarrow 2^V mapping classes to attributes.
                                                                                                                                                                                                                                                                                                                             \bullet (basic) types {\mathcal F} and classes {\mathcal C}, (both finite),
                                                                                                                                                                                                                                                                                                        • typed attributes V, \tau from \mathcal T or C_{0,1} or C_*, C\in \mathcal C,
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Basic Object System Structure

```
We use \mathscr{D}(\mathscr{C}) to denote \bigcup_{C \in \mathscr{C}} \mathscr{D}(C); analogously \mathscr{D}(\mathscr{C}_*)
                                                                                                                                                                                                                                                                                                                                                                                                                                                         is a domain function {\mathcal Q} which assigns to each type a domain, i.e. • \tau\in {\mathcal F} is mapped to {\mathcal D}(\tau),
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                     Definition. A Basic Object System Structure of
                                                                                                                                                                                            • C \in \mathscr{C} is mapped to an <u>infinite</u> set \mathscr{D}(C) of (object) identities. Note: Object identities only have the "=" operation; object identities of different classes are <u>disjoint</u>, i.e. \forall C, D \in \mathscr{C}: C \neq D \to \mathscr{D}(C) \cap \mathscr{D}(D) = \emptyset.
                                                                                                     • C_* and C_{0,1} for C \in \mathscr{C} are mapped to 2^{\mathscr{D}(C)}.
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                       \mathcal{S} = (\mathcal{T}, \mathcal{C}, V, atr)
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Note: We identify objects and object identities, because both uniquely determine each other (cf. OCL 2.0 standard).

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Basic Object System Structure Example

Wanted: a structure for signature

 $\mathcal{S}_0 = (\{\mathit{Int}\}, \{C, D\}, \{x : \mathit{Int}, p : C_{0,1}, n : C_*\}, \{C \mapsto \{p, n\}, D \mapsto \{x\}\})$

Recall: by definition, seek a ${\mathscr D}$ which maps • $\tau \in \mathscr{T}$ to some $\mathscr{D}(\tau)$,

• C_* and $C_{0,1}$ for $C\in \mathscr{C}$ to $\mathscr{D}(C_{0,1})=\mathscr{D}(C_*)=2^{\mathscr{D}(C)}$. $c\in\mathscr{C}$ to some identities $\mathscr{D}(C)$ (infinite, disjoint for different classes),

 $\mathcal{D}(C_{0,1}) = \mathcal{D}(C_*) = 2^{\mathbb{N}^{\dagger} \times \{c\}}$ $\mathcal{D}(D_{0,1}) = \mathcal{D}(D_*) = 2^{\mathbb{D}^{c}(b)}$
$$\begin{split} \mathcal{D}(D) &= \mathbb{N}^{d} \times \{ D_{s}^{2} = \{ A_{b_{s}} \lambda_{b_{s}}, \beta_{b_{s}}, \omega \} \\ &= \mathcal{D}(C_{s}) = 2^{\mathbb{N}^{d}} \times \{ c_{s}^{2} \} \\ \end{split}$$
$$\begin{split} \mathcal{D}(C) &= \mathbb{N}^{d} \times \mathcal{E}C_{\delta}^{2} = \left\{ \ell_{\alpha} \mathcal{L}_{\alpha} \mathcal{L}_{\alpha} \mathcal{L}_{\alpha} \right\} \\ \mathcal{D}(D) &= \mathbb{N}^{d} \times \mathcal{E}C_{\delta}^{2} = \left\{ \ell_{\alpha} \mathcal{L}_{\alpha} \mathcal{L}_{\alpha} \mathcal{L}_{\alpha} \right\} \\ \mathcal{D}_{\delta}(C^{1/2} \mathcal{E}_{\delta} \mathcal{I}, \mathcal{I}, \mathcal{I}, -1) \mathcal{H} \end{split}$$

System State Definition. Let \mathscr{D} by a structure of $\mathscr{S}=(\mathscr{T}\mathscr{K},V,atr)$. A system state of $\mathscr{S}(art,\mathscr{D})$ is a type-consistent mapping $\sigma:\mathscr{D}(\mathscr{D})\to V\to (\mathscr{D}(\mathscr{D})\cup\mathscr{D}(\mathscr{C}))$. We use $\Sigma_{\mathscr{S}}^{\mathscr{D}}$ to denote the set of all system states of \mathscr{S} wrt. \mathscr{D} . We call $u \in \mathscr{D}(\mathscr{C})$ alive in σ if and only if $u \in \text{dom}(\sigma)$. That is, for each $u \in \mathcal{Q}(C)$, $C \in \mathcal{C}$, if $u \in \text{dom}(\sigma)$ • $\operatorname{dom}(\sigma(u)) = \operatorname{der}(C)$ $\bullet \left| \sigma(u) \!\! \left| \!\! \left| \!\! \left(v \right) \in \mathcal{D}(D_*) \right. \right| \text{if } v : D_{0,1} \text{ or } v : D_* \text{ with } D \in \mathcal{C}$ $(\sigma(u)(v) \in \mathcal{D}(\tau) \text{ if } v : \tau, \tau \in \mathcal{F}$ all object identifies partial bundion form to types domains

· 07=0 6 graphy function

Wanted: $\sigma: \mathscr{D}(\mathscr{C}) \nrightarrow (V \nrightarrow (\mathscr{D}(\mathscr{S}) \cup \mathscr{D}(\mathscr{C}_*)))$ such that $\bullet \ \mathrm{dom}(\sigma(u)) = atr(C)$,

 $\quad \sigma(u)(v) \in \mathscr{D}(\tau) \text{ if } v: \tau, \tau \in \mathscr{T}, \quad \bullet \ \sigma(u)(v) \in \mathscr{D}(C_*) \text{ if } v: D_* \text{ with } D \in \mathscr{C}$

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System State Example

Signature, Structure:

 $\mathscr{S}_0 = (\{Int\}, \{C, D\}, \{x: Int, p: C_{0,1}, n: C_*\}, \{C \mapsto \{p, n\}, D \mapsto \{x\}\})$ $\mathscr{D}(Int) = \mathbb{Z}, \quad \mathscr{D}(C) = \{1_C, 2_C, 3_C, \ldots\}, \quad \mathscr{D}(D) = \{1_D, 2_D, 3_D, \ldots\}$

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System State Example

Signature, Structure:

 $\mathscr{S}_0 = (\{Int\}, \{C, D\}, \{x: Int, p: C_{0,1}, n: C_*\}, \{C \mapsto \{p, n\}, D \mapsto \{x\}\})$ $\mathscr{D}(Int) = \mathbb{Z}, \quad \mathscr{D}(C) = \{1_C, 2_C, 3_C, \ldots\}, \quad \mathscr{D}(D) = \{1_D, 2_D, 3_D, \ldots\}$

Wanted: $\sigma: \mathcal{D}(\mathcal{C}) \nrightarrow (V \nrightarrow (\mathcal{D}(\mathcal{T}) \cup \mathcal{D}(\mathcal{C}_*)))$ such that

$$\begin{split} &\sigma(u)(v)\in \mathscr{D}(\tau) \text{ if } v:\tau,\tau\in \mathscr{T},\\ &\sigma(u)(v)\in \mathscr{D}(C_*) \text{ if } v:D_* \text{ with } D\in \mathscr{C} \ . \end{split}$$

Concrete, explicit:

 $\sigma = \{\underbrace{1_C} \mapsto \{p \mapsto \emptyset, n \mapsto \{5_C\}\}, 5_C \mapsto \{p \mapsto \emptyset, n \mapsto \emptyset\}, 1_D \mapsto \{x \mapsto 23\}\}.$

Alternative: symbolic system state

$$\begin{split} \sigma &= \{c_1\} \cdots \{p \mapsto \emptyset, n \mapsto \{c_2\}\}, \{c_2\} \cdots \{p \mapsto \emptyset, n \mapsto \emptyset\}, d \mapsto \{x \mapsto 23\}\}\\ \text{assuming CFO}, d &\in \mathcal{D}(C), d \in \mathcal{D}(D), c_1 \neq c_2. \end{split}$$

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You Are Here.

Course Map $\mathcal{S} = (\mathcal{F}, \mathcal{C}, V, atr)$ $\varphi \in \mathsf{OCL}$ G = (N, E, f)OD $B = (Q_{SD}, q_0, A_{\mathscr{S}}, \rightarrow_{SD}, F_{SD})$ Mathematics $\dot{w}_{\pi} = ((\sigma_i, cons_i, Snd_i))_{i \in \mathbb{N}}$ 21/23

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[Hard et al., 1999] Harel, D., Lachover, H., et al. (1990). Statemate: A working environment for the development of complex seature systems. IEEE Transactions on Software Engineering. Below of Computer Seature Seature

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