Software Design, Modelling and Analysis in UML

Lecture 4: OCL Semantics

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Contents & Goals

Last Lecture:

- OCL Syntax

This Lecture:

- **Educational Objectives:** Capabilities for these tasks/questions:
  - Please un-abbreviate all abbreviations in this OCL expression. ✓
  - Please explain this OCL constraint.
  - Please formalise this constraint in OCL.
  - Does this OCL constraint hold in this system state?
  - Give a system state satisfying this constraint?
  - In what sense is OCL a three-valued logic? For what purpose?
  - How are $\mathscr{D}(C)$ and T_C related?

- **Content:**
  - OCL Semantics
  - OCL Consistency and Satisfiability

Recall

OCL Syntax 1/4: Expressions

$$
\begin{array}{lll}
expr ::= & & \\
& w & : \tau(w) \\
& \mid expr_1 =_\tau expr_2 & : \tau \times \tau \to Bool \\
& \mid \mathsf{oclIsUndefined}_\tau(expr_1) & : \tau \to Bool \\
& \mid \{expr_1, \dots, expr_n\} & : \overbrace{\tau \times \dots \times \tau}^{n\text{-times}} \to Set(\tau) \\
& \mid \mathsf{isEmpty}(expr_1) & : Set(\tau) \to Bool \\
& \mid \mathsf{size}(expr_1) & : Set(\tau) \to Int \\
& \mid \mathsf{allInstances}_C & : Set(\tau_C) \\
& \mid v(expr_1) & : \tau_C \to \tau(v) \\
& \mid r_1(expr_1) & : \tau_C \to \tau_D \\
& \mid r_2(expr_1) & : \tau_C \to Set(\tau_D)
\end{array}
$$

?

Where, given $\mathscr{S} = (\mathscr{T}, \mathscr{C}, V, atr)$,

- $W \supseteq \{self_C : \tau_C \mid C \in \mathscr{C}\}$ is a set of typed logical variables, w has type $\tau(w)$
- τ is any type from $\mathscr{T} \cup T_B \cup T_\mathscr{C} \cup \{Set(\tau_0) \mid \tau_0 \in \mathscr{T} \cup T_B \cup T_\mathscr{C}\}$
  - T_B is a set of (OCL) basic types, in the following we use $T_B = \{Bool, Int, String\}$
  - $T_\mathscr{C} = \{\tau_C \mid C \in \mathscr{C}\}$ is the set of object types,
  - $Set(\tau_0)$ denotes the set-of-τ_0 type for $\tau_0 \in T_B \cup T_\mathscr{C}$ (sufficient because of "flattening" (cf. standard))
- $v : T(v) \in atr(C)$, $T(v) \in \mathscr{T}$,
- $r_1 : D_{0,1} \in atr(C)$,
- $r_2 : D_* \in atr(C)$,
- $C, D \in \mathscr{C}$.

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OCL Syntax 2/4: Constants & Arithmetics

For example:

$$
\begin{array}{lll}
expr ::= & \dots & \\
& \mid \mathsf{true}, \mathsf{false} & : Bool \\
& \mid expr_1 \; \{\mathsf{and}, \mathsf{or}, \mathsf{implies}\} \; expr_2 & : Bool \times Bool \to Bool \\
& \mid \mathsf{not} \; expr_1 & : Bool \to Bool \\
& \mid 0, -1, 1, -2, 2, \dots & : Int \\
& \mid \mathsf{OclUndefined}_\tau & : \tau \\
& \mid expr_1 \; \{+, -, \dots\} \; expr_2 & : Int \times Int \to Int \\
& \mid expr_1 \; \{<, \leq, \dots\} \; expr_2 & : Int \times Int \to Bool
\end{array}
$$

?

Generalised notation:

$$expr ::= \omega(expr_1, \dots, expr_n) \qquad : \tau_1 \times \dots \times \tau_n \to \tau$$

$a + b \leadsto +(a, b)$

with $\omega \in \{+, -, \dots\}$

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OCL Syntax 3/4: Iterate

1/4 and 2/4

$$expr ::= \dots \mid expr_1 \text{->} \mathsf{iterate}(w_1 : \tau_1 \; ; \; w_2 : \tau_2 = expr_2 \mid expr_3)$$

or, with a little renaming,

$$expr ::= \dots \mid expr_1 \text{->} \mathsf{iterate}(iter : \tau_1; \; result : \tau_2 = \underset{\tau_2}{expr_2} \mid \underset{\tau_2}{expr_3})$$

where

- $expr_1$ is of a **collection type** (here: a set $Set(\tau_0)$ for some τ_0),
- $iter \in W$ is called **iterator**, gets type τ_1 (if τ_1 is omitted, τ_0 is assumed as type of $iter$)
- $result \in W$ is called **result variable**, gets type τ_2,
- $expr_2$ in an expression of type τ_2 giving the **initial value** for $result$, ($\mathsf{OclUndefined}_{\tau_2}$, if omitted)
- $expr_3$ is an expression of type τ_2 in which in particular $iter$ and $result$ may appear.

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OCL Syntax 4/4: Context

$$context ::= \mathsf{context} \; w_1 : \cancel{\tau_1}^{C_1}, \; \dots, \; w_n : \cancel{\tau_n}^{C_n} \; \mathsf{inv} : expr$$

where $w_i \in W$ and $\tau_i \in T_\mathscr{C}$ for all $1 \leq i \leq n$, $n \geq 0$.

$$\mathsf{context} \; w_1 : C_1, \dots, w_n : C_n \; \mathsf{inv} : expr$$

is an **abbreviation** for

```
allInstances_{C_1} -> forAll(w_1 : τ_{C_1} |
  ...
    allInstances_{C_n} -> forAll(w_n : τ_{C_n} |
      expr
    )
  ...
)
```

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OCL Semantics: The Task

- Given an OCL expression $expr$ (over signature $\mathscr{S}$), e.g.

$$expr_1 = \text{context } DD \text{ inv} : wen \text{ implies } win > 0$$

- and a system state $\sigma \in \Sigma_{\mathscr{S}}^{\mathscr{D}}$, e.g.

$$\sigma_1 = \{7_{VM} \mapsto \{dd \mapsto \{1_{DD}\}, cp \mapsto \{3_{DD}, 5_{DD}\}\}, \quad 1_{DD} \mapsto \{win \mapsto 13, wen \mapsto \mathit{true}\}, \\ 3_{CP} \mapsto \{wen \mapsto \mathit{true}\}, \quad 5_{CP} \mapsto \{wen \mapsto \mathit{false}\}\}$$

- and a valuation of logical variables $\beta : W \to I(\mathscr{T} \cup T_B \cup T_{\mathscr{C}})$,

- **define** the **interpretation of** $expr$ **in** σ **under** β

$$I[\![\cdot]\!](\cdot\,,\cdot) : OCLExpressions(\mathscr{S}) \times \Sigma_{\mathscr{S}}^{\mathscr{D}} \times (W \to I(\mathscr{T} \cup T_B \cup T_{\mathscr{C}})) \to I(Bool)$$

i.e.

$$I[\![expr]\!](\sigma, \beta) \in \{\mathit{true}, \mathit{false}, \bot_{Bool}\}.$$

OCL Semantics OMG (2006)

Basically business as usual...

(i) Equip each OCL (!) **type** with a reasonable **domain**, i.e. define function

$$I_{(i)} \text{ with } \operatorname{dom}(I) = \mathscr{T} \cup T_B \cup T_{\mathscr{C}}$$

(ii) Equip each **set type** $Set(\tau_0)$ with reasonable **domain**, i.e. define function

$$I_{(ii)} \text{ with } \operatorname{dom}(I) = \{Set(\tau_0) \mid \tau_0 \in \mathscr{T} \cup T_B \cup T_{\mathscr{C}}\}$$

(iii) Equip each **arithmetical operation** with a reasonable **interpretation** (that is, with a **function** operating on the corresponding **domains**).

$$I_{(iii)} \text{ with } \operatorname{dom}(I) = \{+, -, \leq, \dots\}, \text{ e.g., } I_{(iii)}(+) \in I(Int) \times I(Int) \to I(Int)$$

(iv) **Set operations** similar: $I_{(iv)}$ with $\operatorname{dom}(I) = \{\text{isEmpty}, \dots\}$

(v) Equip each **expression** with a reasonable **interpretation**, i.e. define function

$$I_{(v)} : Expr \times \Sigma_{\mathscr{S}}^{\mathscr{D}} \times (W \to I_{(i)}(\mathscr{T} \cup T_B \cup T_{\mathscr{C}})) \to I_{(i)}(Bool)$$

Basically business as usual...

(i) Equip each OCL (!) **type** with a reasonable **domain**, i.e. define function

$$I \text{ with } \mathrm{dom}(I) = \mathscr{T} \cup T_B \cup T_{\mathscr{C}}$$

(ii) Equip each **set type** $Set(\tau_0)$ with reasonable **domain**, i.e. define function

$$I \text{ with } \mathrm{dom}(I) = \{Set(\tau_0) \mid \tau_0 \in \mathscr{T} \cup T_B \cup T_{\mathscr{C}}\}$$

(iii) Equip each **arithmetical operation** with a reasonable **interpretation** (that is, with a **function** operating on the corresponding **domains**).

$$I \text{ with } \mathrm{dom}(I) = \{+, -, \leq, \dots\}, \text{ e.g., } I(+) \in I(Int) \times I(Int) \to I(Int)$$

(iv) **Set operations** similar: I with $\mathrm{dom}(I) = \{\mathsf{isEmpty}, \dots\}$

(v) Equip each **expression** with a reasonable **interpretation**, i.e. define function

$$I : Expr \times \Sigma_{\mathscr{S}}^{\mathscr{D}} \times (W \to I(\mathscr{T} \cup T_B \cup T_{\mathscr{C}})) \to I(Bool)$$

...except for OCL being a **three-valued logic**, and the “iterate” expression.

(i) Domains of OCL and (!) Model Basic Types

Recall: **OCL basic types**

$$T_B = \{Bool, Int, String\}$$

We set:

- $I_{(i)}(Bool) := \{true, false, \bot_{Bool}\}$! three-valued
- $I_{(i)}(Int) := \mathbb{Z} \dot{\cup} \{\bot_{Int}\}$
- $I_{(i)}(String) := \ldots \dot{\cup} \{\bot_{String}\}$

disjoint union

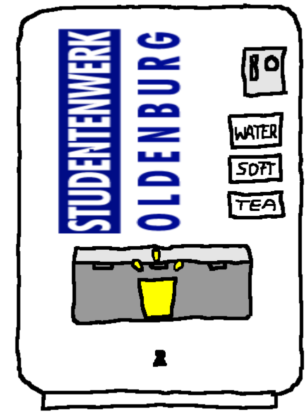
We may omit index τ of $\bot_\tau$ if it is clear from context.

Given signature $\mathscr{S}$ with **model basic types** $\mathscr{T}$ and domain $\mathscr{D}$, set

$$I(T) := \mathscr{D}(T) \dot{\cup} \{\bot_T\}$$

for each model basic type $T \in \mathscr{T}$.

OCL and Model Types?! An Example.



$$\mathscr{S} = (\{Bool, Nat\}, \{VM, CP, DD\},$$
$$\{cp : CP_*, dd : DD_{0,1}, wen : Bool, win : Nat\},$$
$$\{VM \mapsto \{cp, dd\}, CP \mapsto \{wen\}, DD \mapsto \{win, wen\}\})$$

Model Types:

$\mathscr{D}_1(Bool_M) = \{0,1\}$

$\mathscr{D}_1(Nat) = \{0,\dots,255\}$

$\mathscr{D}_1(VM) = \{1_{VM}, 2_{VM}, \dots\}$

OCL Types:

$I(Bool_{OCL}) = \{true, false, \bot_{Bool}\}$ fixed for T_B

$I(Int) = \mathbb{Z} \,\dot\cup\, \{\bot_{Int}\}$

$I(Nat) = \{0,\dots,255\} \,\dot\cup\, \{\bot_{Nat}\}$

$I(\tau_{VM}) = \{1_{VM}, 2_{VM}, \dots\} \,\dot\cup\, \{\bot_{\tau_{VM}}\}$

$I(Bool_M) = \{0,1\} \,\dot\cup\, \{\bot_{Bool_M}\}$

(i) Domains of Object and (ii) Set Types

- Let τ_C be an (OCL) **object type** for a class $C \in \mathscr{C}$.
- We set
$$I(\tau_C) := \mathscr{D}(C) \,\dot{\cup}\, \{\bot_{\tau_C}\}$$

- Let τ be a type from $\mathscr{T} \cup T_B \cup T_{\mathscr{C}}$.
- We set
$$I(Set(\tau)) := 2^{I(\tau)} \,\dot{\cup}\, \{\bot_{Set(\tau)}\}$$

Note: in the OCL standard, only **finite** subsets of $I(\tau)$.
But infinity doesn't scare **us**, so we simply allow it.

(iii) Interpretation of Arithmetic Operations

- **Literals** map to fixed values:

$$I(\underset{\substack{\cap \\ \mathit{OCLExpr}}}{\mathsf{true}}) := \underset{\substack{\cap \\ I(\mathit{Bool})}}{\mathit{true}}, \quad I(\mathsf{false}) := \mathit{false}, \quad I(\overset{\substack{\mathit{OCLExpr} \\ \cup}}{0}) := \overset{\substack{I(\mathit{Int}) \\ \cup}}{0}, \quad I(1) := 1, \ldots$$

$$I(\underset{\substack{\cap \\ \mathit{OCLExpr}}}{\mathsf{OclUndefined}_\tau}) := \underset{\substack{\cap \\ I(\tau)}}{\bot_\tau}$$

(iii) Interpretation of Arithmetic Operations

- **Literals** map to fixed values:

OCLExpr $\quad I(Int)$

$$I(\text{true}) := \mathit{true}, \quad I(\text{false}) := \mathit{false}, \quad I(0) := 0, \quad I(1) := 1, \ldots$$

OCLExpr $\quad I(Bool)$

$$I(\text{OclUndefined}_\tau) := \bot_\tau$$

OCLExpr $\quad I(\tau)$

- **Boolean operations** (defined point-wise for $x_1, x_2 \in I(\tau)$):

$I(\tau) \quad I(Bool)$

$$I(=_\tau)(x_1, x_2) := \begin{cases} \mathit{true} & \text{, if } x_1 \neq \bot_\tau \neq x_2 \text{ and } x_1 = x_2 \\ \mathit{false} & \text{, if } x_1 \neq \bot_\tau \neq x_2 \text{ and } x_1 \neq x_2 \\ \bot_{Bool} & \text{, otherwise} \end{cases}$$

- **Integer operations** (defined point-wise for $x_1, x_2 \in I(Int)$):

$$I(+)(x_1, x_2) := \begin{cases} x_1 + x_2 & \text{, if } x_1 \neq \bot \neq x_2 \\ \bot & \text{, otherwise} \end{cases}$$

but:

$I(\vee)(\mathit{false}, \bot) = \bot$

$I(\wedge)(\mathit{false}, \bot) = \mathit{false}$

and?

Note: There is a **common principle**.

The **interpretation** of an operation (symbol) $\omega : \tau_1 \times \ldots \tau_n \to \tau$

is a function $I(\omega) : I(\tau_1) \times \cdots \times I(\tau_n) \to I(\tau)$ on corresponding semantical domain(s).

(iii) Interpretation of OclIsUndefined

- The **is-undefined** predicate (defined point-wise for $x \in I(\tau)$):

$$I(\text{oclIsUndefined}_\tau)(x) := \begin{cases} \mathit{true} & \text{, if } x = \bot_\tau \\ \mathit{false} & \text{, otherwise} \end{cases}$$

Note: $I(\text{oclIsUndefined}_\tau)$ is **definite**, i.e., it never yields $\bot$.

(iv) Interpretation of Set Operations

Basically the same principle as with arithmetic operations...

Let $\tau \in \mathscr{T} \cup T_B \cup T_{\mathscr{C}}$.

- **Set comprehension** $(x_1, \dots, x_n \in I(\tau))$:

$$I(\{\}_n^\tau)(\underset{\substack{\Uparrow \\ I(\tau)}}{x_1}, \dots, \underset{\substack{\Uparrow \\ I(\tau)}}{x_n}) := \underbrace{\{x_1, \dots, x_n\}}_{\in I(Set(\tau))}$$

  for all $n \in \mathbb{N}_0$

- **Empty-ness check** $(x \in I(Set(\tau)))$:

$$I(\text{isEmpty}^\tau)(x) := \begin{cases} \mathit{true} & \text{, if } x = \emptyset \\ \bot_{Bool} & \text{, if } x = \bot_{Set(\tau)} \\ \mathit{false} & \text{, otherwise} \end{cases}$$

- **Counting** $(x \in I(Set(\tau)))$:

number of elements in x

$$I(\text{size}^\tau)(x) := \begin{cases} |x| & \text{, if } x \neq \bot_{Set(\tau)} \text{ and } x \text{ finite} \\ \bot_{Int} & \text{, otherwise} \end{cases}$$

(v) Putting It All Together

OCL Syntax 1/4: Expressions

$$
\begin{array}{lll}
expr ::= & & \\
\quad w & & : \tau(w) \\
\quad \mid expr_1 =_\tau expr_2 & & : \tau \times \tau \to Bool \\
\quad \mid \mathsf{oclIsUndefined}_\tau(expr_1) & & : \tau \to Bool \\
\quad \mid \{expr_1, \dots, expr_n\} & & : \overbrace{\tau \times \dots \times \tau}^{n\text{-times}} \to Set(\tau) \\
\quad \mid \mathsf{isEmpty}(expr_1) & & : Set(\tau) \to Bool \\
\quad \mid \mathsf{size}(expr_1) & & : Set(\tau) \to Int \\
\quad \mid \mathsf{allInstances}_C & & : Set(\tau_C) \\
\quad \mid v(expr_1) & & : \tau_C \to \tau(v) \\
\quad \mid r_1(expr_1) & & : \tau_C \to \tau_D \\
\quad \mid r_2(expr_1) & & : \tau_C \to Set(\tau_D)
\end{array}
$$

Where, given $\mathscr{S} = (\mathscr{T}, \mathscr{C}, V, atr)$,

- $W \supseteq \{self_C : \tau_C \mid C \in \mathscr{C}\}$ is a set of typed logical variables, w has type $\tau(w)$
- τ is any type from $\mathscr{T} \cup T_B \cup T_\mathscr{C} \cup \{Set(\tau_0) \mid \tau_0 \in \mathscr{T} \cup T_B \cup T_\mathscr{C}\}$
  - T_B is a set of (OCL) basic types, in the following we use $T_B = \{Bool, Int, String\}$
  - $T_\mathscr{C} = \{\tau_C \mid C \in \mathscr{C}\}$ is the set of object types,
  - $Set(\tau_0)$ denotes the set-of-τ_0 type for $\tau_0 \in T_B \cup T_\mathscr{C}$ (sufficient because of "flattening" (cf. standard))
- $v : T(v) \in atr(C)$, $T(v) \in \mathscr{T}$,
- $r_1 : D_{0,1} \in atr(C)$,
- $r_2 : D_* \in atr(C)$,
- $C, D \in \mathscr{C}$.

OCL Syntax 2/4: Constants & Arithmetics

For example:

$$
\begin{array}{ll}
expr ::= \dots & \\
\quad \mid \mathsf{true}, \mathsf{false} & : Bool \\
\quad \mid expr_1 \ \{\mathsf{and}, \mathsf{or}, \mathsf{implies}\} \ expr_2 & : Bool \times Bool \to Bool \\
\quad \mid \mathsf{not}\ expr_1 & : Bool \to Bool \\
\quad \mid 0, -1, 1, -2, 2, \dots & : Int \\
\quad \mid \mathsf{OclUndefined}_\tau & : \tau \\
\quad \mid expr_1 \ \{+, -, \dots\} \ expr_2 & : Int \times Int \to Int \\
\quad \mid expr_1 \ \{<, \leq, \dots\} \ expr_2 & : Int \times Int \to Bool
\end{array}
$$

Generalised notation:

$$expr ::= \omega(expr_1, \dots, expr_n) \qquad : \tau_1 \times \dots \times \tau_n \to \tau$$

with $\omega \in \{+, -, \dots\}$

$a + b \rightsquigarrow +(a, b)$

OCL Syntax 3/4: Iterate

1/4 and 2/4

$$expr ::= \dots \mid expr_1 \text{->}\mathsf{iterate}(w_1 : \tau_1 ; w_2 : \tau_2 = expr_2 \mid expr_3)$$

or, with a little renaming,

$$expr ::= \dots \mid expr_1 \text{->}\mathsf{iterate}(iter : \tau_1 ; result : \tau_2 = \underset{\tau_2}{expr_2} \mid \underset{\tau_2}{expr_3})$$

where

- $expr_1$ is of a **collection type** (here: a set $Set(\tau_0)$ for some τ_0),
- $iter \in W$ is called **iterator**, gets type τ_1 (if τ_1 is omitted, τ_0 is assumed as type of $iter$)
- $result \in W$ is called **result variable**, gets type τ_2,
- $expr_2$ in an expression of type τ_2 giving the **initial value** for $result$, ($\mathsf{OclUndefined}_{\tau_2}$, if omitted)
- $expr_3$ is an expression of type τ_2 in which in particular $iter$ and $result$ may appear.

OCL Syntax 4/4: Context

$$context ::= \mathsf{context}\ w_1 : \overset{C_1}{\cancel{\tau_1}}, \dots, w_n : \overset{C_n}{\cancel{\tau_n}}\ \mathsf{inv} : expr$$

where $w_i \in W$ and $\tau_i \in T_\mathscr{C}$ for all $1 \leq i \leq n$, $n \geq 0$.

context $w_1 : C_1, \dots, w_n : C_n$ inv : $expr$

is an **abbreviation** for

allInstances$_{C_1}$ -> forAll($w_1 : \tau_{C_1}$ |
…
allInstances$_{C_n}$ -> forAll($w_n : \tau_{C_n}$ |
$expr$
)
…
)

Valuations of Logical Variables

$\supseteq \{ self_C \mid C \in \mathscr{C} \}$

- **Recall**: we have typed logical variables $(w \in)\ W$, $\tau(w)$ is the type of w.
- By β, we denote a valuation of the logical variables, i.e. for each $w \in W$,

$$\beta(w) \in I(\tau(w)).$$

$$\beta : W \longrightarrow I(T_B \cup T_C \cup J)$$

$$I(\tau(\text{result})) = I(Bool)$$

$$\beta = \{ \text{result} \mapsto true, \; self_{VM} \mapsto 27_{VM} \}$$

$$w \in W$$

$$I(\tau(self_{VM})) = I(\tau_{self_{VM}})$$

(v) Putting It All Together...

$$
\begin{aligned}
expr ::= {} & w \mid \omega(expr_1, \dots, expr_n) \mid \mathsf{allInstances}_C \mid v(expr_1) \mid r_1(expr_1) \\
& \mid r_2(expr_1) \mid expr_1\text{->}\mathsf{iterate}(v_1 : \tau_1 \,;\, v_2 : \tau_2 = expr_2 \mid expr_3)
\end{aligned}
$$

double-u

$\Sigma_{\mathscr{D}}^{\mathscr{D}}$ ⇓

- $I[\![w]\!](\sigma, \beta) := \beta(w)$

e.g. $self_{vm}$ omega

- $I[\![\omega(expr_1, \dots, expr_n)]\!](\sigma, \beta) := I(\omega)(I[\![expr_1]\!](\sigma, \beta), \dots, I[\![expr_n]\!](\sigma, \beta))$

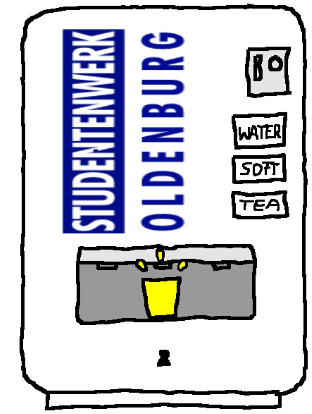
e.g. $+(13, result)$ e.g. $I(+)(I[\![13]\!](\sigma,\beta), I[\![result]\!](\sigma,\beta))$

- $I[\![\mathsf{allInstances}_C]\!](\sigma, \beta) := \underbrace{\mathrm{dom}(\sigma)}_{\text{all alive objects in } \sigma} \cap \underbrace{\mathscr{D}(C)}_{\text{objects of class } C}$

Note: in the OCL standard, $\mathrm{dom}(\sigma)$ is assumed to be **finite**.
Again: doesn't scare us.

Example

$\beta = \{ self_{VM} \mapsto 7_{VM},\ x \mapsto 27 \}$



$$\mathscr{S} = (\{Bool, Nat\}, \{VM, CP, DD\},$$
$$\{cp : CP_*, dd : DD_{0,1}, wen : Bool, win : Nat\},$$
$$\{VM \mapsto \{cp, dd\}, CP \mapsto \{wen\}, DD \mapsto \{win, wen\})$$

$$\sigma_1 = \{7_{VM} \mapsto \{dd \mapsto \{1_{DD}\}, cp \mapsto \{3_{DD}, 5_{DD}\}\}, \quad 1_{DD} \mapsto \{win \mapsto 13, wen \mapsto \mathsf{true}\},$$
$$3_{CP} \mapsto \{wen \mapsto \mathsf{true}\}, \quad 5_{CP} \mapsto \{wen \mapsto \mathsf{false}\}\}$$

- $I[\![w]\!](\sigma, \beta) := \beta(w)$ 1
- $I[\![\mathsf{allInstances}_C]\!](\sigma, \beta) := \mathrm{dom}(\sigma) \cap \mathscr{D}(C)$ 2
- $I[\![\omega(expr_1, \dots, expr_n)]\!](\sigma, \beta) := I(\omega)(I[\![expr_1]\!](\sigma, \beta), \dots I[\![expr_n]\!](\sigma, \beta))$ 3

- $I[\![self_{VM}]\!](\sigma_1, \beta) \overset{1}{=} \beta(self_{VM}) = 7_{VM}$
- $I[\![x]\!](\sigma_1, \beta) \overset{1}{=} \beta(x) = 27$ (*)
- $I[\![allInstances_{CP}]\!](\sigma_1, \beta) \overset{2}{=} \mathrm{dom}(\sigma_1) \cap \mathscr{D}(CP)$

$= \{7_{VM}, 1_{DD}, 3_{CP}, 5_{CP}\} \cap \{1_{CP}, 2_{CP}, \dots\} = \underbrace{\{3_{CP}, 5_{CP}\}}_{(**)}$

- $I[\![x + allInstances_{CP} \to size]\!](\sigma_1, \beta) = 29$

3 $(I[\![+(x, size(allInstances_{CP}))]\!](\sigma_1, \beta))$

$I(+)(I[\![x]\!](\sigma_1, \beta),$
$I[\![allInstances_{CP} \to size]\!](\sigma_1, \beta))$

$\overset{(*)}{=} I(+)(27, I(size)(I[\![allInstances_{CP}]\!](\sigma_1, \beta)))$

$\overset{(**)}{=} I(+)(27, I(size)(\{3_{CP}, 5_{CP}\}))$

$= I(+)(27, 2) = 29$ ✓

(v) Putting It All Together...

$$expr ::= w \mid \omega(expr_1, \ldots, expr_n) \mid \mathsf{allInstances}_C \mid v(expr_1) \mid r_1(expr_1)$$
$$\mid r_2(expr_1) \mid expr_1 \texttt{->iterate}(v_1 : \tau_1 \; ; v_2 : \tau_2 = expr_2 \mid expr_3)$$

Assume $expr_1 : \tau_C$ for some $C \in \mathscr{C}$. Set $u_1 := I[\![expr_1]\!](\sigma, \beta) \in \mathscr{D}(\tau_C)$.

- $I[\![v(expr_1)]\!](\sigma, \beta) := \begin{cases} (\sigma(u_1))(v), & \text{if } u_1 \in \mathrm{dom}(\sigma) \\ \bot, & \text{otherwise} \end{cases}$

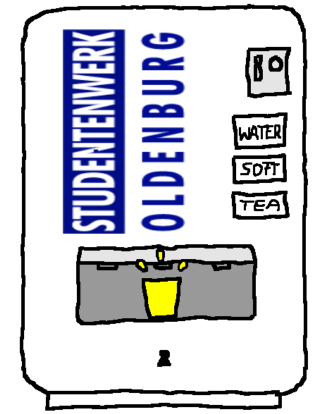
$\in atr$ $\; \vdots \tau_C$

$v : T,$

$T \in \mathscr{T}$

Example

$\beta = \{ self_{VM} \mapsto 7_{VM},\ self_{CP} \mapsto 3_{CP},\ y \mapsto 1_{CP} \}$



$$\begin{aligned}\mathscr{S} = (&\{Bool, Nat\}, \{VM, CP, DD\},\\ &\{cp : CP_*, dd : DD_{0,1}, wen : Bool, win : Nat\},\\ &\{VM \mapsto \{cp, dd\}, CP \mapsto \{wen\}, DD \mapsto \{win, wen\})\end{aligned}$$

$$\begin{aligned}\sigma_1 = \{&7_{VM} \mapsto \{dd \mapsto \{1_{DD}\}, cp \mapsto \{3_{DD}, 5_{DD}\}\}, \quad 1_{DD} \mapsto \{win \mapsto 13, wen \mapsto \underset{1}{\cancel{true}}\},\\ &3_{CP} \mapsto \{wen \mapsto \underset{1}{\cancel{true}}\}, \quad 5_{CP} \mapsto \{wen \mapsto \underset{0}{\cancel{false}}\}\}\end{aligned}$$

Assume $expr_1 : \tau_C$ for some $C \in \mathscr{C}$. Set $u_1 := I[\![expr_1]\!](\sigma, \beta) \in \mathscr{D}(\tau_C)$.

- $I[\![v(expr_1)]\!](\sigma, \beta) := \begin{cases} \sigma(u_1)(v) & \text{, if } u_1 \in \mathrm{dom}(\sigma) \\ \bot & \text{, otherwise} \end{cases}$

- $I[\![wen(\overbrace{self_{CP}}^{expr_1})]\!](\sigma_1, \beta) = (\sigma_1(u_1))(wen) = 1$

  $u_1 = I[\![self_{CP}]\!](\sigma_1, \beta) = 3_{CP}$ alive!

- $I[\![wen(y)]\!](\sigma_1, \beta) = \bot$

  $I[\![y]\!](\sigma_1, \beta) = 1_{CP}$ not alive!

(v) Putting It All Together...

$$
\begin{aligned}
expr ::= {} & w \mid \omega(expr_1, \dots, expr_n) \mid \mathsf{allInstances}_C \mid v(expr_1) \mid r_1(expr_1) \\
& \mid r_2(expr_1) \mid expr_1\texttt{->}\mathsf{iterate}(v_1 : \tau_1 \,;\, v_2 : \tau_2 = expr_2 \mid expr_3)
\end{aligned}
$$

Assume $expr_1 : \tau_C$ for some $C \in \mathscr{C}$. Set $u_1 := I[\![expr_1]\!](\sigma, \beta) \in \mathscr{D}(\tau_C)$.

- $I[\![v(expr_1)]\!](\sigma, \beta) := \begin{cases} \sigma(u_1)(v) & \text{, if } u_1 \in \mathrm{dom}(\sigma) \\ \bot & \text{, otherwise} \end{cases}$

- $I[\![r_1(expr_1)]\!](\sigma, \beta) := \begin{cases} u & \text{, if } u_1 \in \mathrm{dom}(\sigma) \text{ and } \sigma(u_1)(r_1) = \{u\} \\ \bot & \text{, otherwise} \end{cases}$

  $r_1 : D_{0,1}$

- $I[\![r_2(expr_1)]\!](\sigma, \beta) := \begin{cases} \sigma(u_1)(r_2) & \text{, if } u_1 \in \mathrm{dom}(\sigma) \\ \bot & \text{, otherwise} \end{cases}$

  $r_2 : D_*$

  Recall: σ evaluates r_2 of type C_* to a set.

(v) Putting It All Together...

$$\begin{aligned} expr ::= {} & w \mid \omega(expr_1, \dots, expr_n) \mid \mathsf{allInstances}_C \mid v(expr_1) \mid r_1(expr_1) \\ & \mid r_2(expr_1) \mid expr_1\texttt{->}\mathsf{iterate}(v_1 : \tau_1 \,;\, v_2 : \tau_2 = expr_2 \mid expr_3) \end{aligned}$$

$\perp_{\tau_2}$, if $I[\![expr_1]\!](\sigma, \beta) = \perp_{\tau},$

- $I[\![expr_1\texttt{->}\mathsf{iterate}(v_1 : \tau_1 \,;\, v_2 : \tau_2 = expr_2 \mid expr_3)]\!](\sigma, \beta)$

$$:= \begin{cases} I[\![expr_2]\!](\sigma, \beta) & \text{, if } I[\![expr_1]\!](\sigma, \beta) = \emptyset \\ iterate(hlp, v_1, v_2, expr_3, \sigma, \beta') & \text{, otherwise} \end{cases}$$

where $\beta' = \beta[hlp \mapsto I[\![expr_1]\!](\sigma, \beta), v_2 \mapsto I[\![expr_2]\!](\sigma, \beta)]$ and

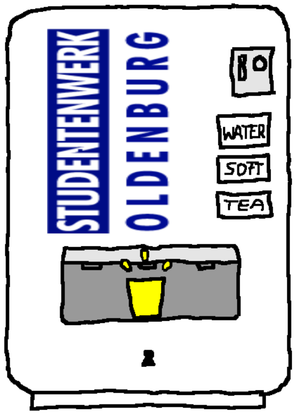
- $iterate(hlp, v_1, v_2, expr_3, \sigma, \beta')$

$$:= \begin{cases} I[\![expr_3]\!](\sigma, \beta'[v_1 \mapsto x]) & \text{, if } \beta'(hlp) = \{x\} \\ I[\![expr_3]\!](\sigma, \beta'') & \text{, if } \beta'(hlp) = X \,\dot\cup\, \{x\} \text{ and } X \neq \emptyset \end{cases}$$

where $\beta'' = \beta'[v_1 \mapsto x, v_2 \mapsto iterate(hlp, v_1, v_2, expr_3, \sigma, \beta'[hlp \mapsto X])]$

Quiz: Is (our) I a function?

Example



$$\mathscr{S} = (\{Bool, Nat\}, \{VM, CP, DD\},$$
$$\{cp : CP_*, dd : DD_{0,1}, wen : Bool, win : Nat\},$$
$$\{VM \mapsto \{cp, dd\}, CP \mapsto \{wen\}, DD \mapsto \{win, wen\})$$

$$\sigma_1 = \{7_{VM} \mapsto \{dd \mapsto \{1_{DD}\}, cp \mapsto \{3_{DD}, 5_{DD}\}\}, \quad 1_{DD} \mapsto \{win \mapsto 13, wen \mapsto \underline{\textsf{true}}\},$$
$$3_{CP} \mapsto \{wen \mapsto \underline{\textsf{true}}\}, \quad 5_{CP} \mapsto \{wen \mapsto \underline{\textsf{false}}\}\}$$

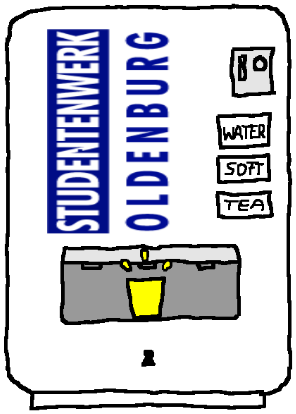
expr = context DD inv : wen implies $win > 0$

$I⟦ \mathit{allInstances}_{DD} \to \mathit{forAll}(\mathit{self}_{DD} : \tau_{DD} \mid \mathit{implies}(\mathit{wen}(\mathit{self}_{DD}), >(\mathit{win}(\mathit{self}_{DD}), 0)))⟧(\sigma, \vartheta)$

iterate(... ; result = true | result and ...)

= true

Another Example



$$\begin{aligned}\mathscr{S} = (&\{Bool, Nat\}, \{VM, CP, DD\},\\ &\{cp : CP_*, dd : DD_{0,1}, wen : Bool, win : Nat\},\\ &\{VM \mapsto \{cp, dd\}, CP \mapsto \{wen\}, DD \mapsto \{win, wen\})\end{aligned}$$

$$\begin{aligned}\sigma_1 = \{&7_{VM} \mapsto \{dd \mapsto \{1_{DD}\}, cp \mapsto \{3_{DD}, 5_{DD}\}\}, \quad 1_{DD} \mapsto \{win \mapsto 13, wen \mapsto \textsf{true}\},\\ &3_{CP} \mapsto \{wen \mapsto \textsf{true}\}, \quad 5_{CP} \mapsto \{wen \mapsto \textsf{false}\}\}\end{aligned}$$

$I[\![$ context VM inv: $cp \rightarrow forAll(c \mid wen(c) = true)$

or $cp \rightarrow forAll(c \mid wen(c) = false)]\!](\sigma_1, \emptyset) = false$

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