

Software Design, Modelling and Analysis in UML

Lecture 8: Class Diagrams III

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Contents & Goals

Last Lectures:

- completed class diagrams... except for associations.

This Lecture:

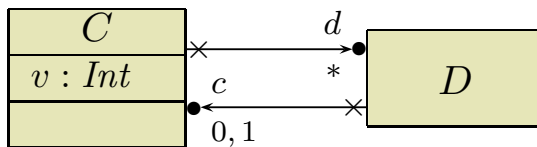
- **Educational Objectives:** Capabilities for following tasks/questions.
 - Please explain this class diagram with associations.
 - Which annotations of an association arrow are semantically relevant?
 - What's a role name? What's it good for?
 - What is "multiplicity"? How did we treat them semantically?
 - What is "reading direction", "navigability", "ownership", ...?
 - What's the difference between "aggregation" and "composition"?
- **Content:**
 - Study concrete syntax for "associations".
 - (**Temporarily**) extend signature, define mapping from diagram to signature.
 - Study effect on OCL.
 - Btw.: where do we put OCL constraints?

Overview

- **Class diagram:**



Alternative presentation:



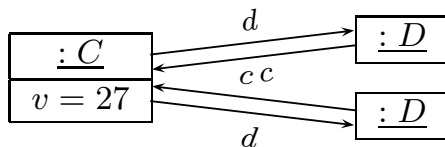
- **Signature:**

$$\mathcal{S} = (\{Int\}, \{C, D\}, \{v : Int, d : D_*, c : C_{0,1}\}, \{C \mapsto \{v, d\}, D \mapsto \{c\}\})$$

- **Example system state:**

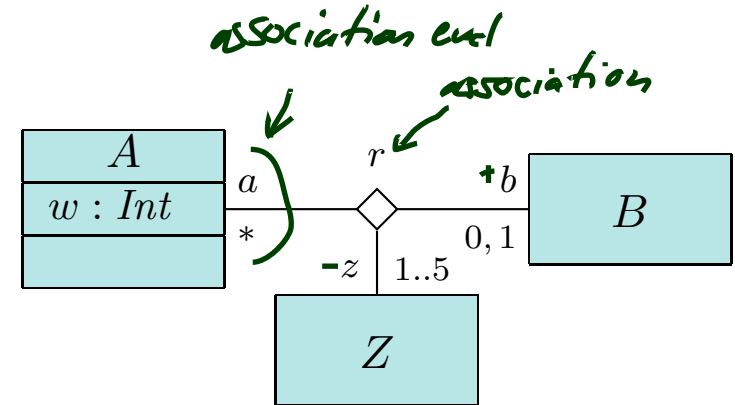
$$\sigma = \{1_C \mapsto \{v \mapsto 27, d \mapsto \{5_D, 7_D\}\}, 5_D \mapsto \{c \mapsto \{1_C\}\}, 7_D \mapsto \{c \mapsto \{1_C\}\}\}$$

- **Object diagram:**



control of $\lambda : a \rightarrow \mathcal{N}_r \rightarrow ?$

- **Class diagram (with ternary association):**



- **Signature:** extend again;

represent **association** r with **association ends** a , b , and z (each with multiplicity, visibility, etc.)

- **Example system state:**

$$\sigma = \{1_A \mapsto \{w \mapsto 13\}, 1_B \mapsto \emptyset, 1_Z \mapsto \emptyset\}$$

$$\lambda = \{ r \mapsto \{(1_A, 1_B, 1_Z)\} \}_1$$

Handwritten notes: $3_z \mapsto \emptyset$, $2_z \mapsto \emptyset$

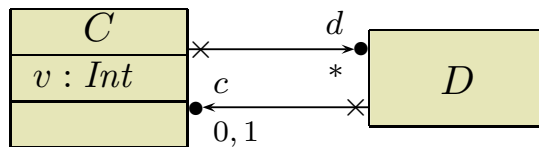
	a	b	z
λ_1	1_A	1_B	1_Z
λ_2	1_A	1_B	2_Z

Overview

- **Class diagram:**



Alternative presentation:



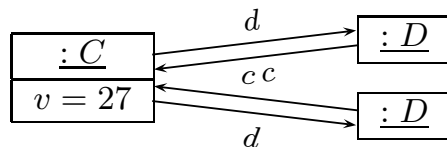
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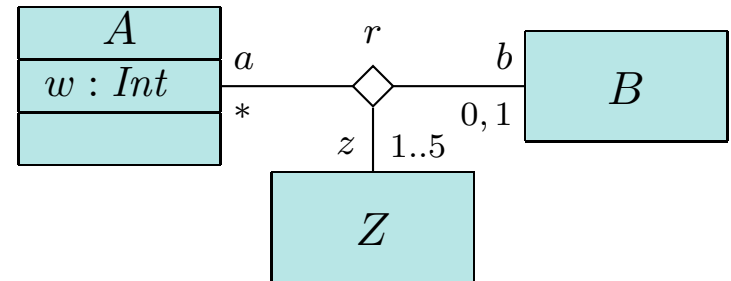
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- **Object diagram:**



- **Class diagram (with ternary association):**



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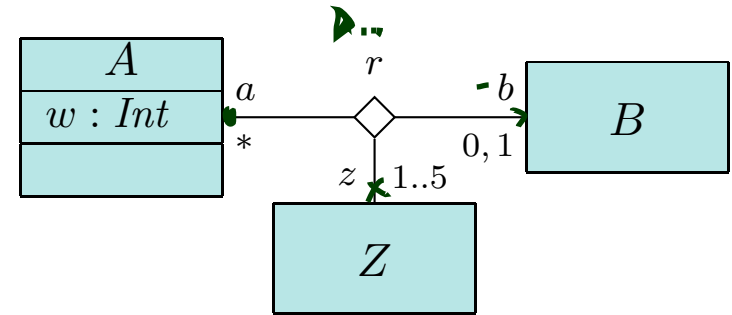
$$\lambda = \{ r \mapsto \{(1_A, 1_B, 1_Z)\} \}$$

- **Object diagram:** No...

Plan

- (i) Study association **syntax**.
- (ii) Extend **signature** accordingly.
- (iii) Define (σ, λ) **system states** with
 - **objects** in σ
(instances of classes),
 - **links** in λ
(instances of associations).
- (iv) Change **syntax** of OCL to refer to **association ends**.
- (v) Adjust **interpretation** I accordingly.
- (vi) ... go back to the special case of $C_{0,1}$ and C_* attributes.

- **Class diagram** (with ternary association):



- **Signature:** extend again;
represent **association** r
with **association ends** a , b , and z
(each with multiplicity, visibility, etc.)

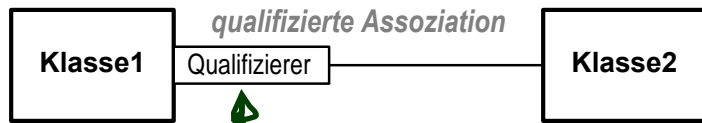
- **Example system state:**

$$\sigma = \{1_A \mapsto \{w \mapsto 13\}, 1_B \mapsto \emptyset, 1_Z \mapsto \emptyset\}$$

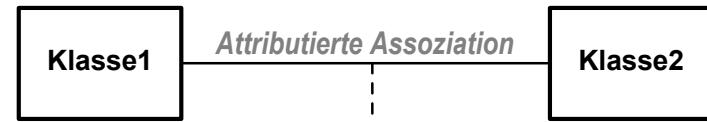
$$\lambda = \{ r \mapsto \{(1_A, 1_B, 1_Z)\} \}$$

- **Object diagram:** No...

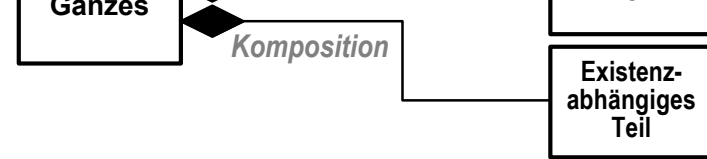
Associations: Syntax



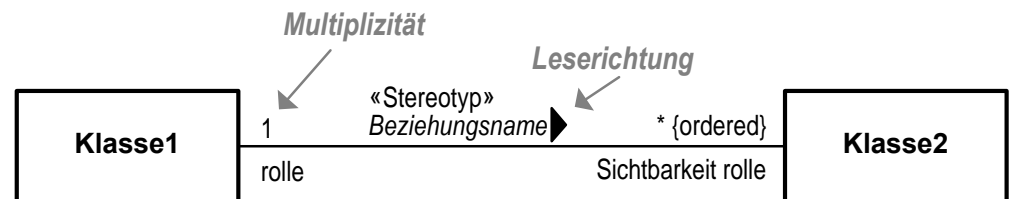
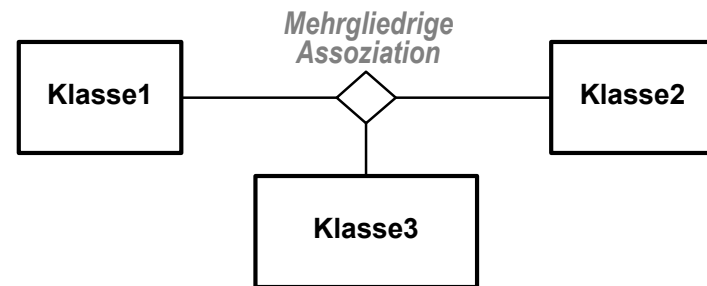
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exercises



More Association Syntax (OMG, 2011b, 61;43)

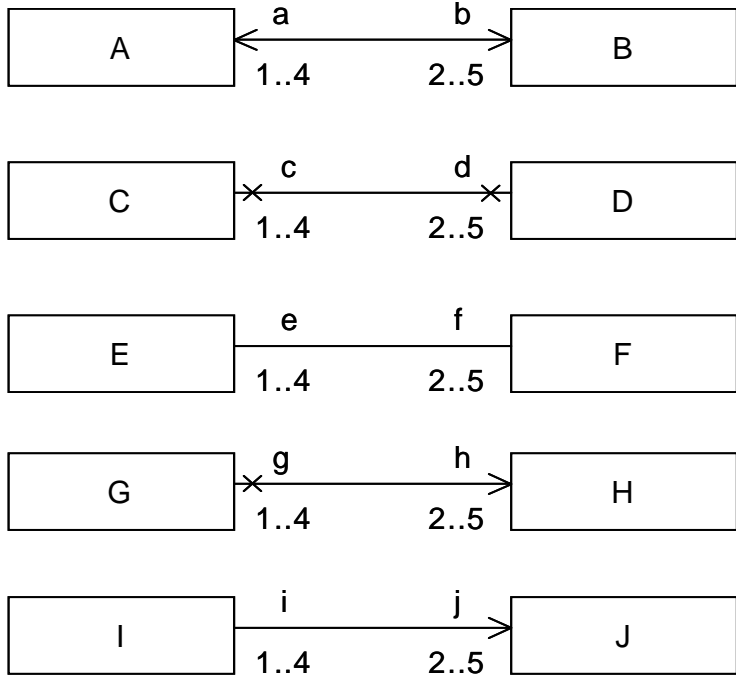


Figure 7.23 - Examples of navigable ends



Figure 7.19 - Graphic notation indicating exactly one association end owned by the association

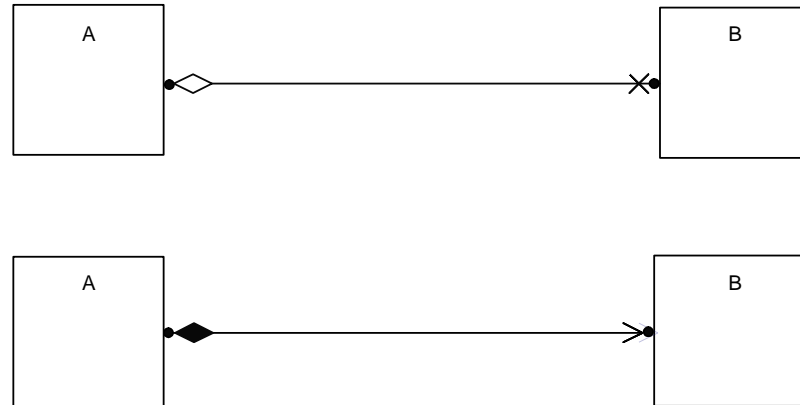


Figure 7.20 - Combining line path graphics

So, What Do We (Have to) Cover?

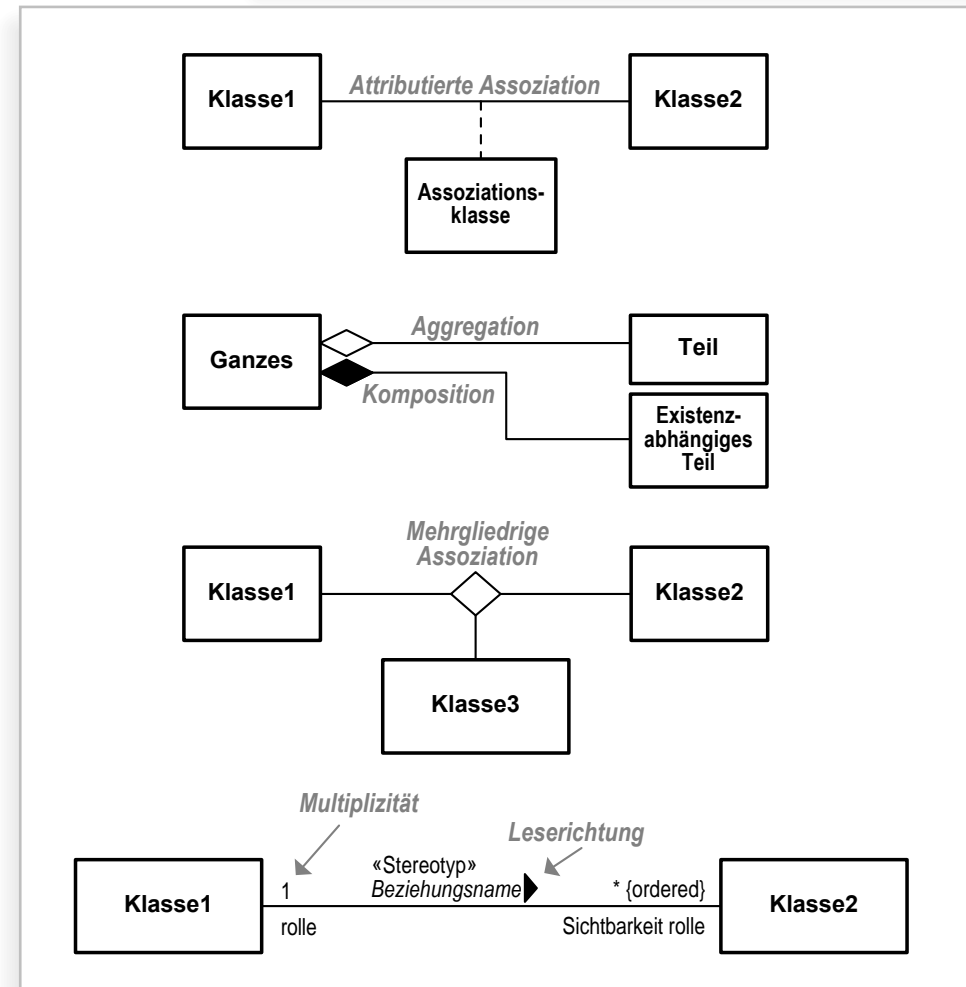
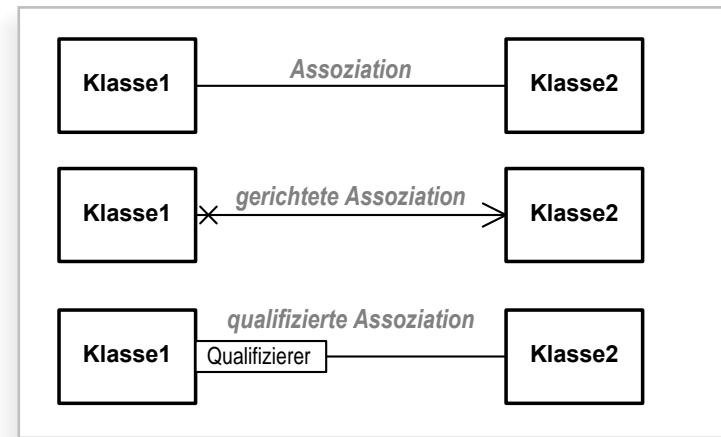
An **association** has

- a **name**,
- a **reading direction**, and
- at least two **ends**.

Each **end** has

- a **role name**,
- a **multiplicity**,
- a set of **properties**, such as **unique**, **ordered**, etc.
- a **qualifier**, (*not in lecture*)
- a **visibility**,
- a **navigability**,
- an **ownership**,
- and possibly a **diamond**. (*exercises*)

Wanted: places in the signature to represent the information from the picture.



(Temporarily) Extend Signature: Associations

Only for the course of Lectures 08/09 we assume that each element in V is

- either a **basic type attribute** $\langle v : T, \xi, expr_0, P_v \rangle$ with $T \in \mathcal{T}$ (as before),
- or an **association** of the form

$$\langle r : \overbrace{\langle role_1 : C_1, \mu_1, P_1, \xi_1, \nu_1, o_1 \rangle, \dots, \langle role_n : C_n, \mu_n, P_n, \xi_n, \nu_n, o_n \rangle}^{\text{association end}} \rangle$$

association name (with arrow pointing to r)

the class where this association end is located (with arrow pointing to C_n)

- $n \geq 2$ (at least two ends),
- $r, role_i$ are just **names**, $C_i \in \mathcal{C}$, $1 \leq i \leq n$,
- the **multiplicity** μ_i is an expression of the form

$$\mu ::= N..M \mid N..* \mid \mu, \mu \quad (N, M \in \mathbb{N})$$

- P_i is a set of **properties** (as before),
- $\xi \in \{+, -, \#, \sim\}$ (as before),
- $\nu_i \in \{\times, -, >\}$ is the **navigability**,
- $o_i \in \mathbb{B}$ is the **ownership**.
- N for $N..N$,
- $*$ for $0..*$ (use with care!)

(Temporarily) Extend Signature: Associations

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$$\langle r : \begin{array}{l} \langle role_1 : C_1, \mu_1, P_1, \xi_1, \nu_1, o_1 \rangle, \\ \vdots \\ \langle role_n : C_n, \mu_n, P_n, \xi_n, \nu_n, o_n \rangle \end{array} \rangle$$

- $n \geq 2$ (at least two ends),
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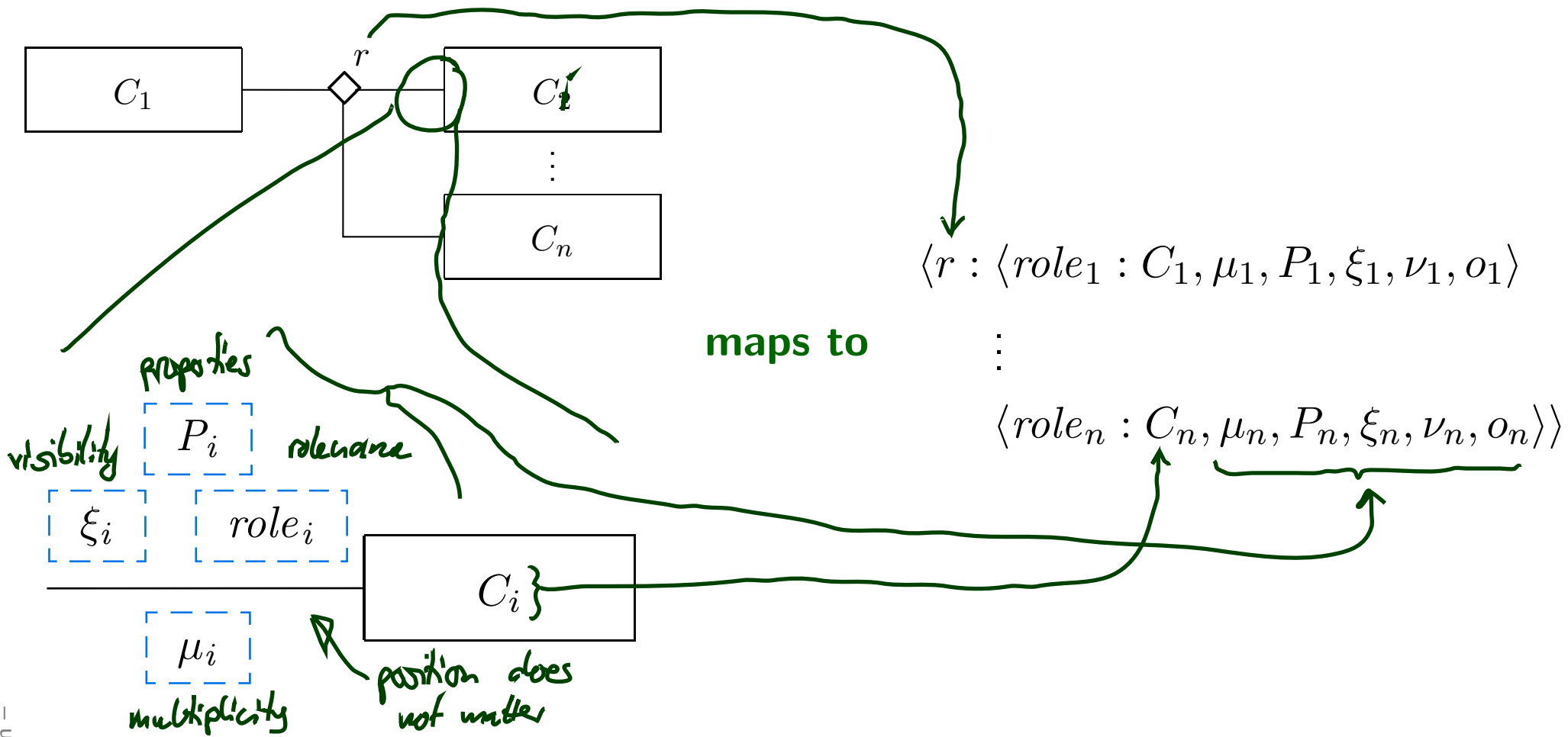
$$\mu ::= N..M \mid N..* \mid \mu, \mu \quad (N, M \in \mathbb{N})$$

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- $\nu_i \in \{\times, -, >\}$ is the **navigability**,
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Multiplicity abbreviations:

- N for $N..N$,
- $*$ for $0..*$ (use with care!)

From Association Lines to Extended Signatures

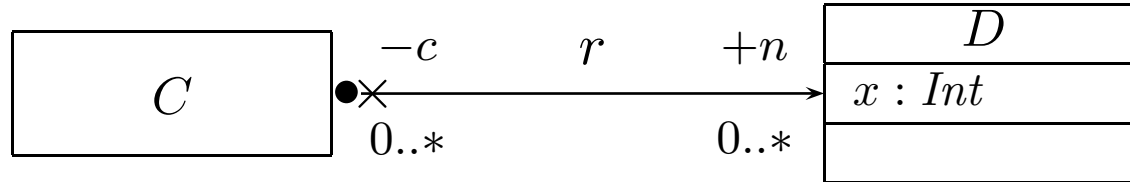


$$o_i = \begin{cases} 1 & , \text{ if } \text{---} \bullet \boxed{C_i} \\ 0 & , \text{ if } \text{---} \boxed{C_i} \end{cases}$$

$$\nu_i = \begin{cases} \times & , \text{ if } \text{---} \times \boxed{C_i} \\ - & , \text{ if } \text{---} \boxed{C_i} \\ > & , \text{ if } \text{---} \rightarrow \boxed{C_i} \end{cases}$$

navigability

Association Example



Signature:

$$\mathcal{I} = \left(\{Int\}, \{C, D\}, \{x: Int, \right.$$

∇ assumption
 \circ
 \downarrow

$$\begin{aligned}
 & \langle r: \langle c: C, 0..*, \{unique\}, -, x, 1 \rangle, \\
 & \langle n: D, 0..*, \{unique\}, +, >, 0 \rangle \rangle \\
 & \left. \{C \mapsto \emptyset, D \mapsto \{x\}\} \right)
 \end{aligned}$$

What If Things Are Missing?

Most components of associations or association end may be omitted.
For instance (OMG, 2011b, 17), Section 6.4.2, proposes the following rules:

- **Name:** Use

$$A_{-}\langle C_1 \rangle_{-}\cdots_{-}\langle C_n \rangle$$

if the name is missing.

Example:



- **Reading Direction:** no default.
- **Role Name:** use the class name at that end in lower-case letters

Example:



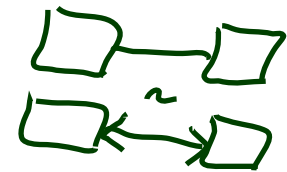
Other convention: (used e.g. by modelling tool Rhapsody)



What If Things Are Missing?

- **Multiplicity:** 1

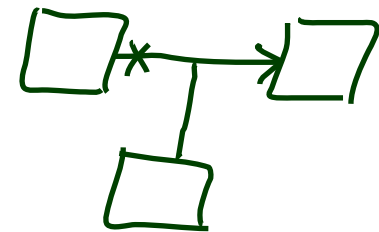
In my opinion, it's safer to assume 0..1 or * (for 0..*) if there are no fixed, written, agreed conventions ("expect the worst").



- **Properties:** \emptyset

- **Visibility:** public

- **Navigability and Ownership:** not so easy. (OMG, 2011b, 43)



“Various options may be chosen for showing navigation arrows on a diagram.

In practice, it is often convenient to suppress some of the arrows and crosses and just show exceptional situations:

- *Show all arrows and x's. Navigation and its absence are made completely explicit.*
- *Suppress all arrows and x's. No inference can be drawn about navigation.*
This is similar to any situation in which information is suppressed from a view.
- *Suppress arrows for associations with navigability in both directions, and show arrows only for associations with one-way navigability.*

In this case, the two-way navigability cannot be distinguished from situations where there is no navigation at all; however, the latter case occurs rarely in practice.”

Wait, If Omitting Things...

- ...**is causing so much trouble** (e.g. leading to misunderstanding), why does the standard say “**In practice, it is often convenient...**”?

Is it a good idea to trade **convenience** for **precision/unambiguity**?

It depends.

- Convenience as such is a **legitimate goal**.
- In UML-As-Sketch mode, precision “**doesn't matter**”, so convenience (for writer) can even be a primary goal.
- In UML-As-Blueprint mode, **precision** is the **primary goal**. And misunderstandings are in most cases annoying.

But: (even in UML-As-Blueprint mode)

If all associations in your model have multiplicity *, then it's probably a good idea not to write all these *'s.

So: tell the reader about your convention and leave out the *'s.

Temporarily (Lecture 8/9) Extended Signature

Definition. An (Extended) Object System **Signature** (with Associations) is a quadruple $\mathcal{S} = (\mathcal{I}, \mathcal{C}, V, atr)$ where

- ...
- each element of V is
 - **either** a **basic type attribute** $\langle v : T, \xi, expr_0, P_v \rangle$ with $T \in \mathcal{I}$
 - **or** an **association** of the form

$$\langle r : \begin{array}{l} \langle role_1 : C_1, \mu_1, P_1, \xi_1, \nu_1, o_1 \rangle, \\ \vdots \\ \langle role_n : C_n, \mu_n, P_n, \xi_n, \nu_n, o_n \rangle \end{array} \rangle$$

- ...
- $atr : \mathcal{C} \rightarrow 2^{\{v \in V \mid v:T, T \in \mathcal{I}\}}$
maps each class to its set of **basic type** (!) attributes.

In other words:

- only **basic type attributes** “**belong**” to a class (may appear in $atr(C)$),
- **associations** are not “owned” by a particular class (do not appear in any $atr(C)$), but “**live on their own**”.

Associations: Semantics

Associations in General

Recall: We consider associations of the following form:

$$\langle r : \langle role_1 : C_1, \mu_1, P_1, \xi_1, \nu_1, o_1 \rangle, \dots, \langle role_n : C_n, \mu_n, P_n, \xi_n, \nu_n, o_n \rangle \rangle$$

Only these parts are relevant for extended system states:

$$\langle r : \langle role_1 : C_1, -, P_1, -, -, - \rangle, \dots, \langle role_n : C_n, -, P_n, -, -, - \rangle \rangle$$

(recall: we assume $P_1 = P_n = \{\text{unique}\}$).

The UML standard thinks of associations as **n-ary relations** which “**live on their own**” in a system state.

That is, **links** (= association instances)

- **do not** belong (in general) to certain objects (in contrast to pointers, e.g.)
- are “first-class citizens” **next to objects**,
- are (in general) **not** directed (in contrast to pointers).

Links in System States

$$\langle r : \langle role_1 : C_1, -, P_1, -, -, - \rangle, \dots, \langle role_n : C_n, -, P_n, -, -, - \rangle \rangle$$

Only for the course of lectures 8 / 9 we change the definition of system states:

Definition. Let \mathcal{D} be a structure of the (extended) signature with associations $\mathcal{S} = (\mathcal{T}, \mathcal{C}, V, atr)$.

A **system state** of \mathcal{S} wrt. \mathcal{D} is a pair (σ, λ) consisting of

- a type-consistent mapping (as before)

only basic type attributes in here

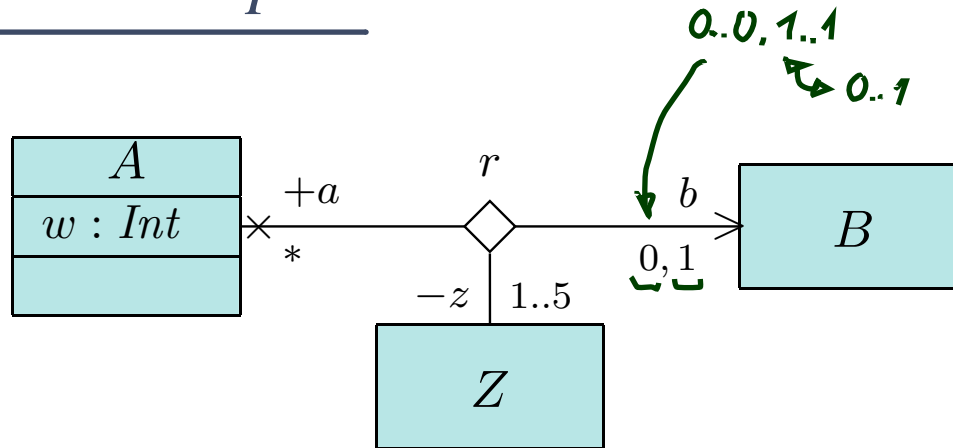
$$\sigma : \mathcal{D}(\mathcal{C}) \rightarrow (atr(\mathcal{C}) \rightarrow \mathcal{D}(\mathcal{T})),$$

- a mapping λ which maps each association $\langle r : \langle role_1 : C_1 \rangle, \dots, \langle role_n : C_n \rangle \rangle \in V$ to a **relation**

$$\lambda(r) \subseteq \mathcal{D}(C_1) \times \dots \times \mathcal{D}(C_n)$$

(i.e. a set of type-consistent n -tuples of identities).

Association / Link Example



Signature:

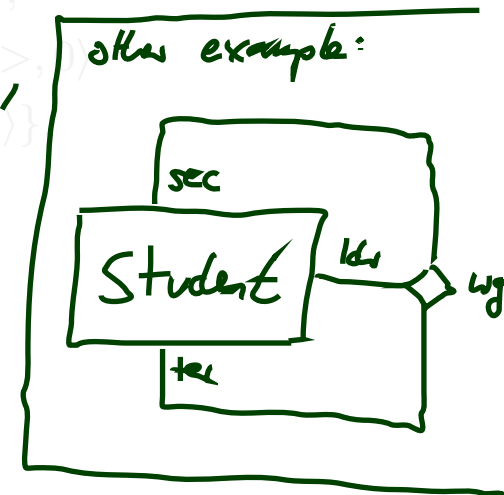
$$\mathcal{Y} = (\{Int\}, \{A, B, Z\}, \{w: Int,$$

$$\mathcal{I} = (\{Int\}, \{A, B, Z\}, \{ \langle r: \langle a: A, 0..*, +, \{unique\}, \times, 0 \rangle, \langle b: B, 0..1, +, \{unique\}, \times, 0 \rangle, \langle z: Z, 1..5, -, \{unique\}, -, 0 \rangle \rangle \}, \{A \mapsto \{w\}, B \mapsto \emptyset, C \mapsto \emptyset\})$$

System state:

$$\sigma: \{ 1_A \mapsto \{w \mapsto 0\}, 2_A \mapsto \{w \mapsto 1\}, 3_A \mapsto \{w \mapsto 2\}, 10_B \mapsto \emptyset, 11_B \mapsto \emptyset, 27_z \mapsto \emptyset, 28_z \mapsto \emptyset \}$$

$$\lambda: \{ r \mapsto \{ (1_A, 10_B, 27_z), (2_A, 10_B, 27_z), (1_A, 11_B, 27_z), (5_A, 13_B, 29_z) \}$$



Associations and OCL

OCL and Associations: Syntax

Recall: OCL syntax as introduced in Lecture 3, interesting part:

$$\begin{array}{l|ll} \text{expr} ::= \dots & r_1(\text{expr}_1) & : \tau_C \rightarrow \tau_D & r_1 : D_{0,1} \in \text{atr}(C) \\ & r_2(\text{expr}_1) & : \tau_C \rightarrow \text{Set}(\tau_D) & r_2 : D_* \in \text{atr}(C) \end{array}$$

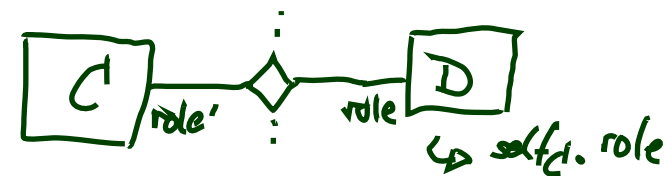
OCL and Associations: Syntax

Recall: OCL syntax as introduced in Lecture 3, interesting part:

$expr ::= \dots$	$ r_1(expr_1) : \tau_C \rightarrow \tau_D$	$r_1 : D_{0,1} \in atr(C)$
	$ r_2(expr_1) : \tau_C \rightarrow Set(\tau_D)$	$r_2 : D_* \in atr(C)$

Now becomes

$expr ::= \dots$	$ role(expr_1) : \tau_C \rightarrow \tau_D$
	$ role(expr_1) : \tau_C \rightarrow Set(\tau_D)$



$\mu = 0..1$ or $\mu = 1..1$
otherwise

if there is

$\langle r : \dots, \langle role : D, \mu, -, -, -, - \rangle, \dots, \langle role' : C, -, -, -, -, - \rangle, \dots \rangle \in V$ or

$\langle r : \dots, \langle role' : C, -, -, -, -, - \rangle, \dots, \langle role : D, \mu, -, -, -, - \rangle, \dots \rangle \in V, \quad \underline{role \neq role'}$.

Note:

- Association name as such **does not occur** in OCL syntax, role names do.
- $expr_1$ has to denote an object of a class which “participates” in the association.

OCL and Associations Syntax: Example

$$\begin{array}{l|l|l} \text{expr} ::= \dots & \text{role}(\text{expr}_1) & : \tau_C \rightarrow \tau_D & \mu = 0..1 \text{ or } \mu = 1..1 \\ & \text{role}(\text{expr}_1) & : \tau_C \rightarrow \text{Set}(\tau_D) & \text{otherwise} \end{array}$$

if there is

$$\langle r : \dots, \langle \text{role} : D, \mu, -, -, -, - \rangle, \dots, \langle \text{role}' : C, -, -, -, -, - \rangle, \dots \rangle \in V \text{ or}$$

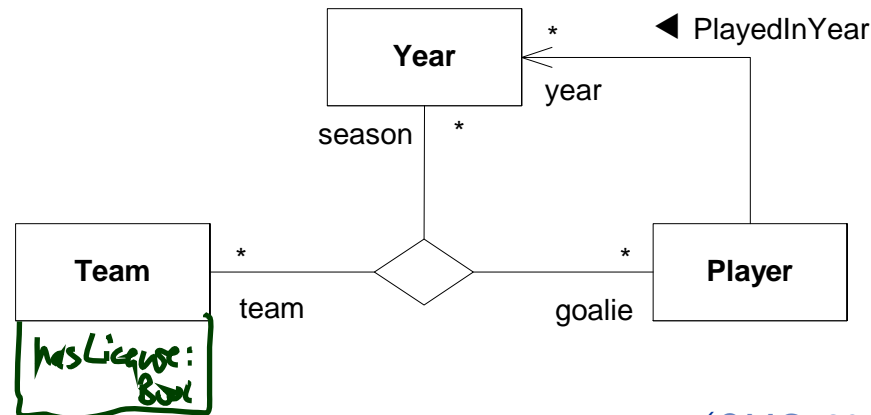
$$\langle r : \dots, \langle \text{role}' : C, -, -, -, -, - \rangle, \dots, \langle \text{role} : D, \mu, -, -, -, - \rangle, \dots \rangle \in V, \text{role} \neq \text{role}'.$$


Figure 7.21 - Binary and ternary associations (OMG, 2011b, 44).

context Player inv: team \rightarrow forall(t | t.hasLicense)

OCL and Associations: Semantics

Recall:

Assume $expr_1 : \tau_C$ for some $C \in \mathcal{C}$. Set $u_1 := I[expr_1](\sigma, \beta) \in \mathcal{D}(T_C)$.

- $I[r_1(expr_1)](\sigma, \beta) := \begin{cases} u & , \text{ if } u_1 \in \text{dom}(\sigma) \text{ and } \sigma(u_1)(r_1) = \{u\} \\ \perp & , \text{ otherwise} \end{cases}$
- $I[r_2(expr_1)](\sigma, \beta) := \begin{cases} \sigma(u_1)(r_2) & , \text{ if } u_1 \in \text{dom}(\sigma) \\ \perp & , \text{ otherwise} \end{cases}$

Now needed:

$$I[role(expr_1)]((\sigma, \lambda), \beta)$$

- We cannot simply write $\sigma(u)(role)$.

Recall: *role* is (**for the moment**) not an attribute of object u (not in $atr(C)$).

- What we have is $\lambda(r)$ (with association name r , not with role name *role*!).

$$\langle r : \dots, \langle role : D, \mu, -, -, -, - \rangle, \dots, \langle role' : C, -, -, -, -, - \rangle, \dots \rangle$$

But it yields a set of n -tuples, of which **some** relate u and some instances of D .

- *role* denotes the position of the D 's in the tuples constituting the value of r .

References

References

Oestereich, B. (2006). *Analyse und Design mit UML 2.1, 8. Auflage*. Oldenbourg, 8. edition.

OMG (2011a). Unified modeling language: Infrastructure, version 2.4.1. Technical Report formal/2011-08-05.

OMG (2011b). Unified modeling language: Superstructure, version 2.4.1. Technical Report formal/2011-08-06.