

Software Design, Modelling and Analysis in UML

Lecture 15: Hierarchical State Machines I

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Contents & Goals

Last Lecture:

- step, RTC-step, divergence
- initial state, UML model semantics (so far)
- create, destroy actions

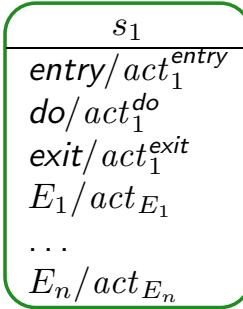
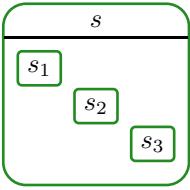
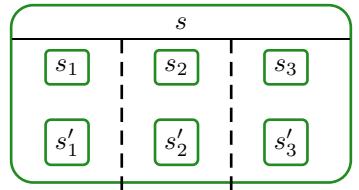
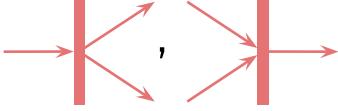
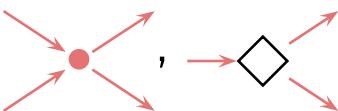
This Lecture:

- **Educational Objectives:** Capabilities for following tasks/questions.
 - What is simple state, OR-state, AND-state?
 - What is a legal state configuration?
 - What is a legal transition?
 - How is enabledness of transitions defined for hierarchical state machines?
- **Content:**
 - Legal state configurations
 - Legal transitions
 - Rules (i) to (v) for hierarchical state machines

Hierarchical State-Machines

The Full Story

UML distinguishes the following **kinds of states**:

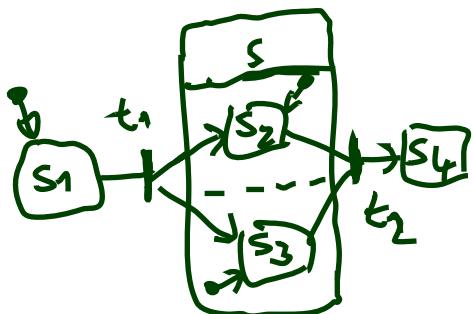
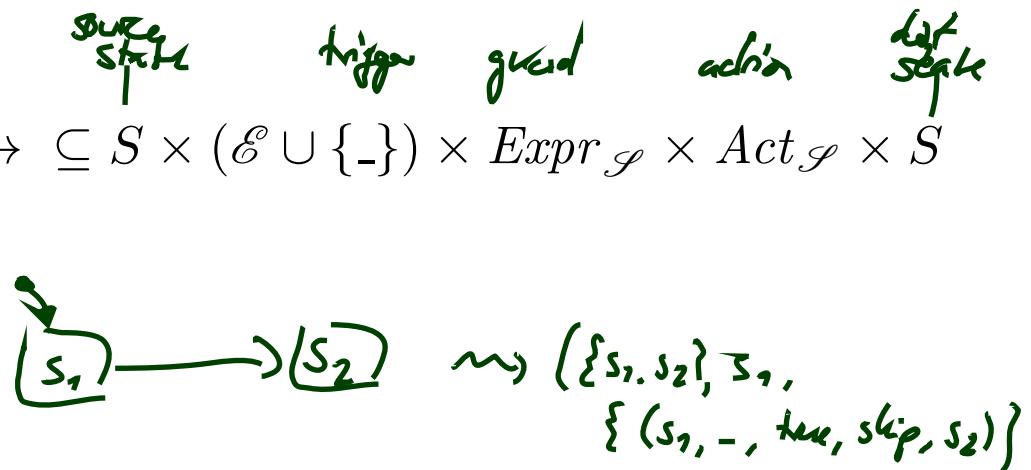
	example	
simple state		
final state		
composite state	 	
OR		pseudo-state
AND		initial (shallow) history deep history fork/join junction, choice entry point exit point terminate
		         

Representing All Kinds of States

- Until now: init state

(S, s_0, \rightarrow) , $s_0 \in S, \rightarrow \subseteq S \times (\mathcal{E} \cup \{-\}) \times Expr_{\mathcal{S}} \times Act_{\mathcal{S}} \times S$

set of states / edges



NEW: $(\{s_1, s_2, s_3, s_4, s, \text{top}\}, \{t_1, t_2\},$
 $\{t_1 \mapsto (\{s_1\}, \{s_2, s_3\}), \dots\}$
 $\{t_2 \mapsto (\{s_1, \text{true}, \text{skip}\}), \dots\})$

Representing All Kinds of States

- Until now:

$$(S, s_0, \rightarrow), \quad s_0 \in S, \rightarrow \subseteq S \times (\mathcal{E} \cup \{-\}) \times Expr_{\mathcal{S}} \times Act_{\mathcal{S}} \times S$$

- From now on: (hierarchical) state machines

$$(S, kind, region, \rightarrow, \psi, annot)$$

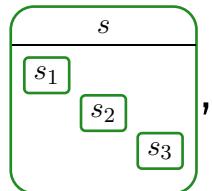
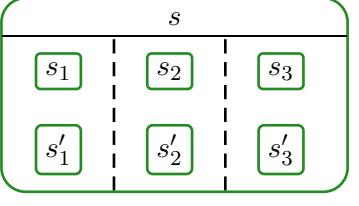
where

- $S \supseteq \{top\}$ is a finite set of states (as before),
- $kind : S \rightarrow \{st, init, fin, shist, dhist, fork, join, junc, choi, ent, exi, term\}$ is a function which labels states with their **kind**, (new)
- $region : S \rightarrow 2^{2^S}$ is a function which characterises the **regions** of a state, (new)
- \rightarrow is a set of transitions, (changed)
- $\psi : (\rightarrow) \rightarrow 2^S \times 2^S$ is an **incidence function**, and (new)
- $annot : (\rightarrow) \rightarrow (\mathcal{E} \cup \{-\}) \times Expr_{\mathcal{S}} \times Act_{\mathcal{S}}$ provides an annotation for each transition. (new)

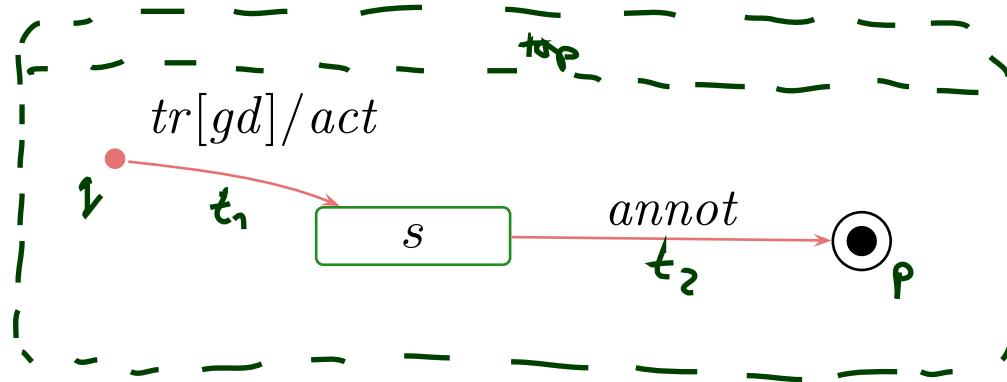
(s_0 is then redundant — replaced by proper state (!) of kind ‘init’.)

From UML to Hierarchical State Machine: By Example

$(S, \text{kind}, \text{region}, \rightarrow, \psi, \text{annot})$

	example	$\in S$	kind	region
simple state		s	st	\emptyset
final state		q	fin	\emptyset
composite state				
OR		s	st	$\{\{s_1, s_2, s_3\}\}$
AND		s	st	$\{\{s_1, s'_1\}, \{s_2, s'_2\}, \{s_3, s'_3\}\}$
submachine state	(later)	—	—	—
pseudo-state	p^* , ...	p	init ...	\emptyset
$(s, \text{kind}(s))$ for short				

From UML to Hierarchical State Machine: By Example



... denotes $(S, kind, region, \rightarrow, \psi, annot) =$

$$\begin{aligned}
 & \underbrace{\left(\left\{ (q, \text{init}), (s, \text{st}), (p, \text{fin}), (\top_p, \text{st}) \right\} \right)}_{S, kind}, \\
 & \underbrace{\left\{ q \mapsto \emptyset, p \mapsto \emptyset, s \mapsto \emptyset, \top_p \mapsto \{\{q, s, p\}\} \right\}}_{region}, \\
 & \underbrace{\left\{ t_1, t_2 \right\}, \quad \underbrace{\left\{ t_1 \mapsto (\{q\}, \{s\}), t_2 \mapsto (\{s\}, \{p\}) \right\}}_{\psi}}_{\rightarrow}, \\
 & \underbrace{\left\{ t_1 \mapsto (tr[gd], act), t_2 \mapsto annot \right\}}_{annot}
 \end{aligned}$$

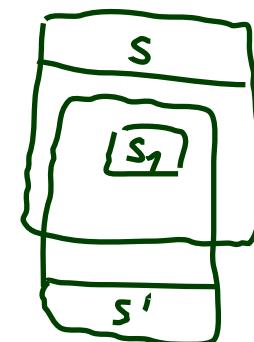
Well-Formedness: Regions

	$\in S$	<i>kind</i>	$region \subseteq 2^S, S_i \subseteq S$	$child \subseteq S$
final state	s	<i>fin</i>	\emptyset	\emptyset
pseudo-state	s	<i>init, ...</i>	\emptyset	\emptyset
simple state	s	<i>st</i>	\emptyset	\emptyset
composite state	s	<i>st</i>	$\{S_1, \dots, S_n\}, n \geq 1$	$S_1 \cup \dots \cup S_n$
implicit top state	<i>top</i>	<i>st</i>	$\{S_1\}$	S_1

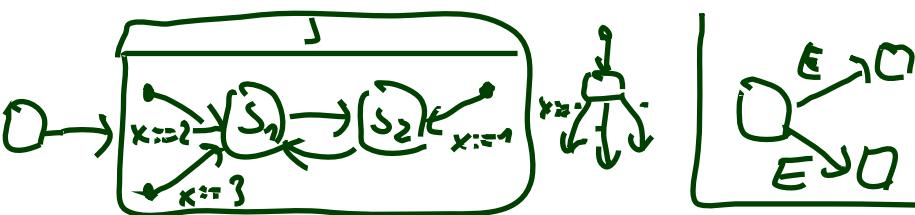
- Final and pseudo states **must not comprise** regions.
- States $s \in S$ with $kind(s) = st$ **may comprise** regions.

Naming conventions can be defined based on regions:

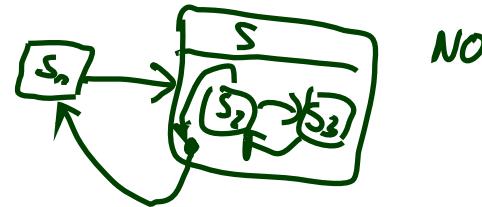
- No region: simple state.
- One region: OR-state.
- Two or more regions: AND-state.
- Each state (except for *top*) **must** lie in exactly one region.
- **Note:** The region function induces a **child** function.
- **Note:** Diagramming tools (like Rhapsody) can ensure well-formedness.



Well-Formedness Continued

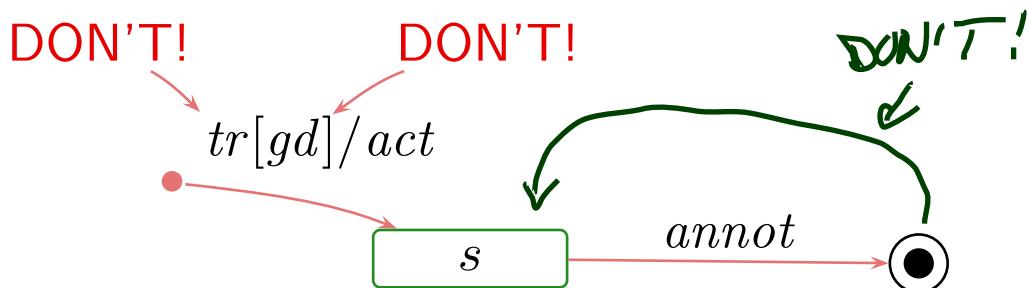


- Each non-empty region has **exactly one** initial pseudo-state and at least one transition from there to a state of the region, i.e.
 - for each $s \in S$ with $\text{region}(s) = \{S_1, \dots, S_n\}$,
 - for each $1 \leq i \leq n$, there exists exactly one initial pseudo-state $(s_1^i, \text{init}) \in S_i$ and at least one transition $t \in \rightarrow$ with s_1^i as source,
- Initial pseudo-states are not targets of transitions.

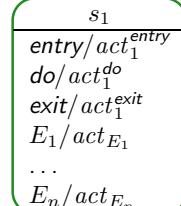
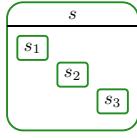
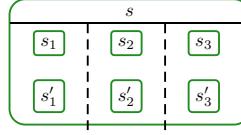


For simplicity:

- The target of a transition with initial pseudo-state source in S_i is (also) in S_i .
- Transitions from initial pseudo-states have no trigger or guard, i.e. $t \in \rightarrow$ from s with $\text{kind}(s) = \text{st}$ implies $\text{annot}(t) = (_, \text{true}, \text{act})$.
- Final states are not sources of transitions.



Plan

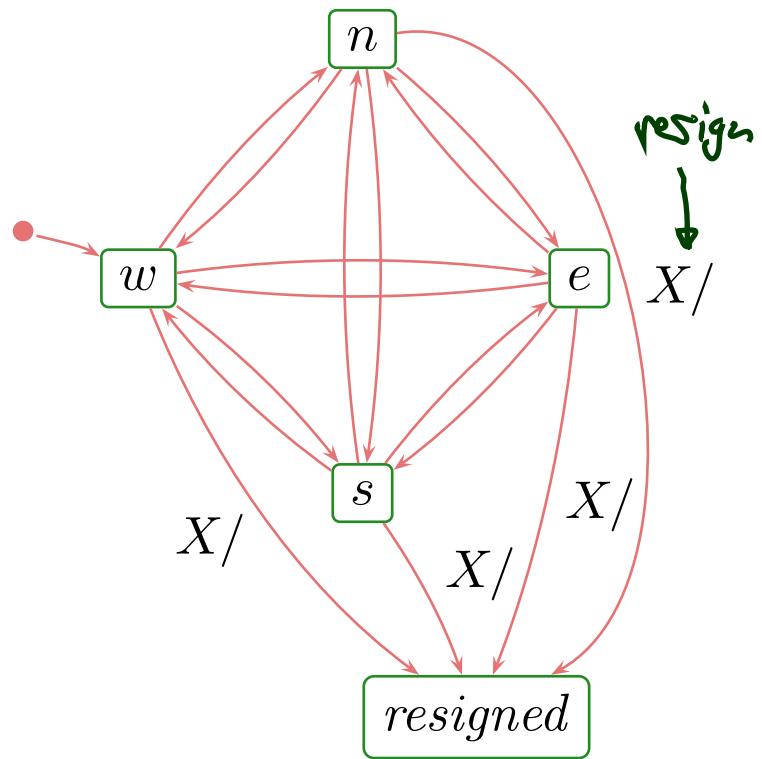
	example	
simple state		pseudo-state
final state		initial (shallow) history
composite state		deep history
OR		fork/join
AND		junction, choice
		entry point
		exit point
		terminate
		submachine state
		

- Composite states.
- Initial pseudostate, final state.
- Entry/do/exit actions, internal transitions.
- History and other pseudostates, the rest.

Composite States

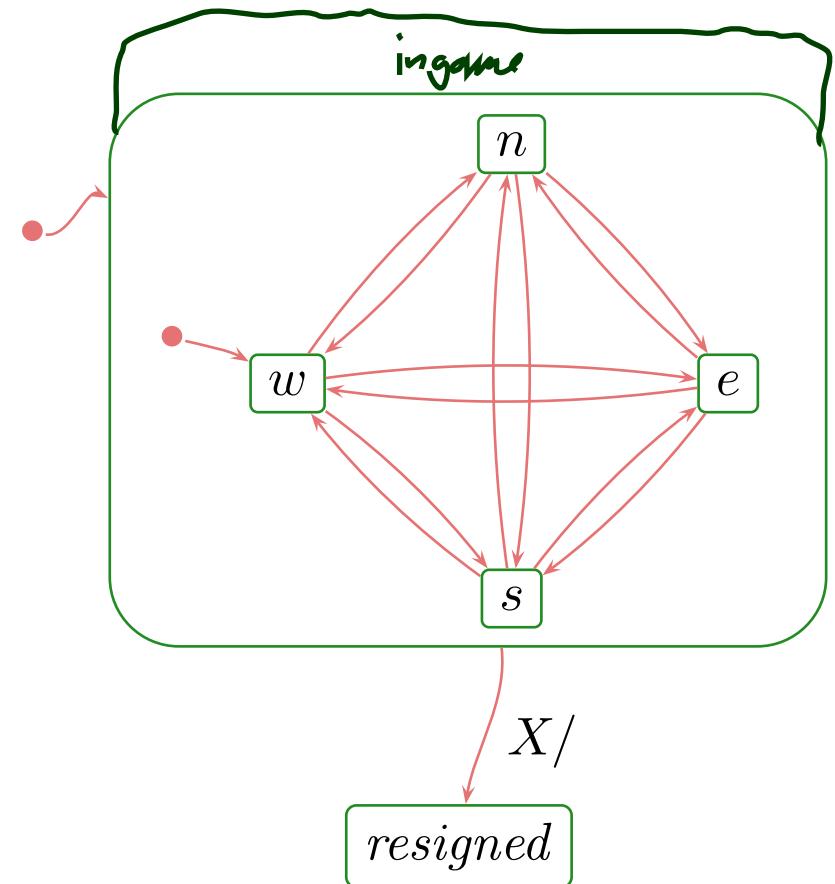
Composite States

- In a sense, composite states are about **abbreviation, structuring**, and **avoiding redundancy**.
- **Idea:** in Tron, for the Player's Statemachine, instead of



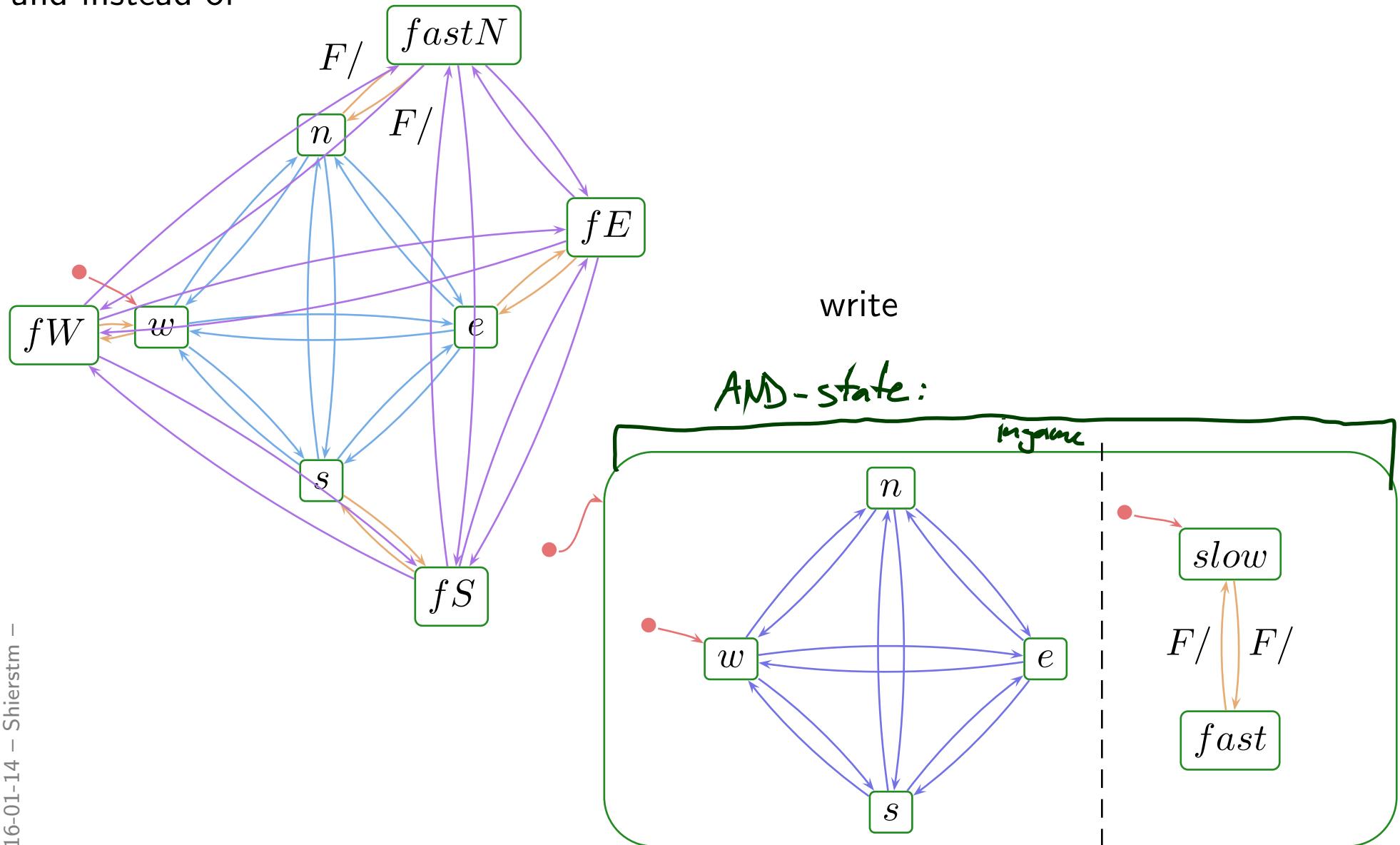
write

OR-state:



Composite States

and instead of



Composite States: Blessing or Curse?

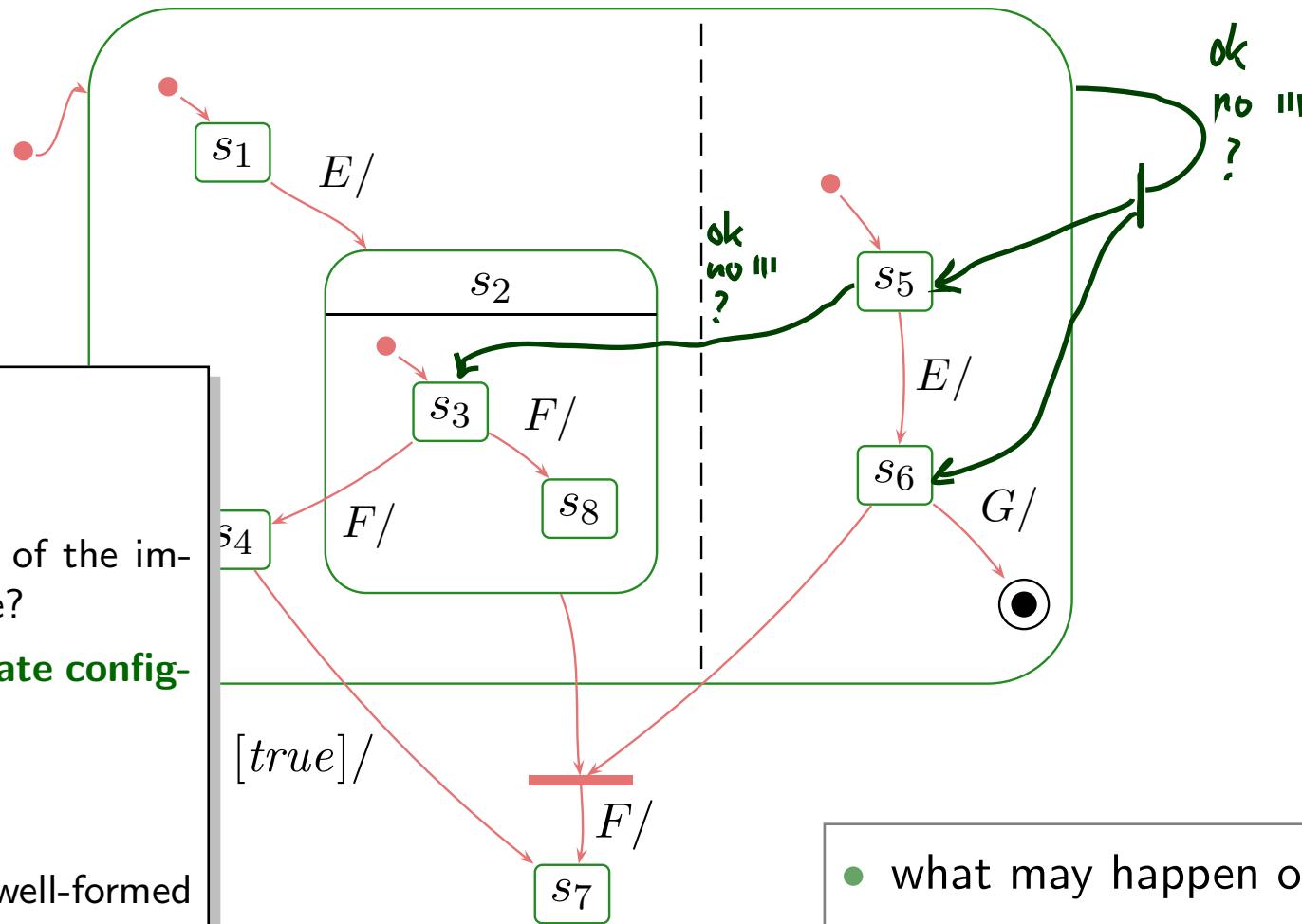
Plan:

States:

- what is the type of the implicit *st* attribute?
- what are **legal state configurations**?

Transitions:

- what are **legal** / well-formed transitions?
- when is a legal transition enabled?
- which effects do transitions have?



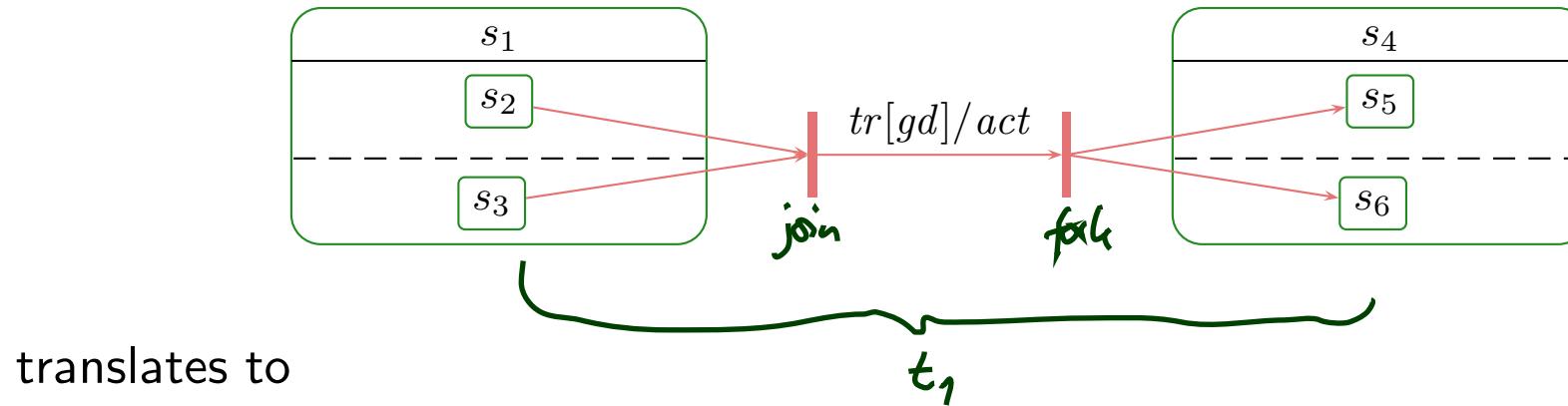
- what may happen on *E*?
- what may happen on *E, F*?
- can *E, G* kill the object?
- ...

Syntax: Fork/Join

- For simplicity, we consider transitions with (possibly) multiple sources and targets, i.e.

$$\psi : (\rightarrow) \rightarrow (2^S \setminus \emptyset) \times (2^S \setminus \emptyset)$$

- For instance,



$$(S, kind, region, \underbrace{\{t_1\}}_{\rightarrow}, \underbrace{\{t_1 \mapsto (\{s_2, s_3\}, \{s_5, s_6\})\}}_{\psi}, \underbrace{\{t_1 \mapsto (tr, gd, act)\}}_{annot})$$

- Naming convention: $\psi(t) = (source(t), target(t))$.

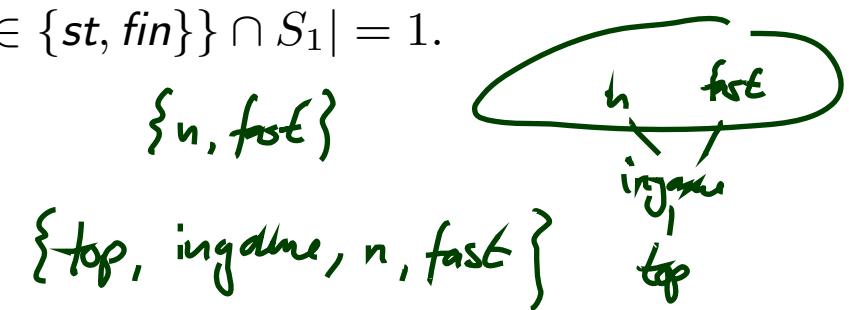
State Configuration

- The type of (implicit attribute) st is from now on **a set of states**, i.e.
 $\mathcal{D}(S_{MC}) = 2^S$
- A set $S_1 \subseteq S$ is called **(legal) state configurations** if and only if
 - $top \in S_1$, and
 - with each state $s \in S_1$ that has a non-empty region $\emptyset \neq R \in region(s)$, exactly one (non pseudo-state) child of s is in S_1 , i.e.

$$|\{s \in R \mid kind(s) \in \{st, fin\}\} \cap S_1| = 1.$$

$\{n, fast, \omega\}$
not OK

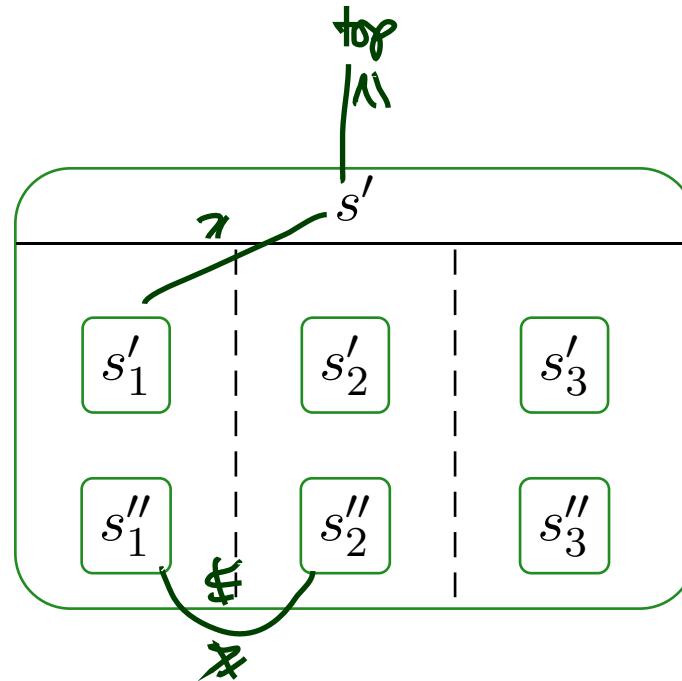
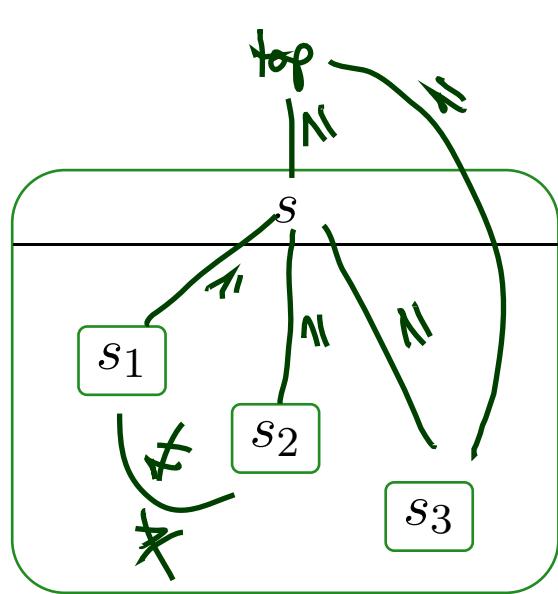
$\{n\}$
not OK



A Partial Order on States

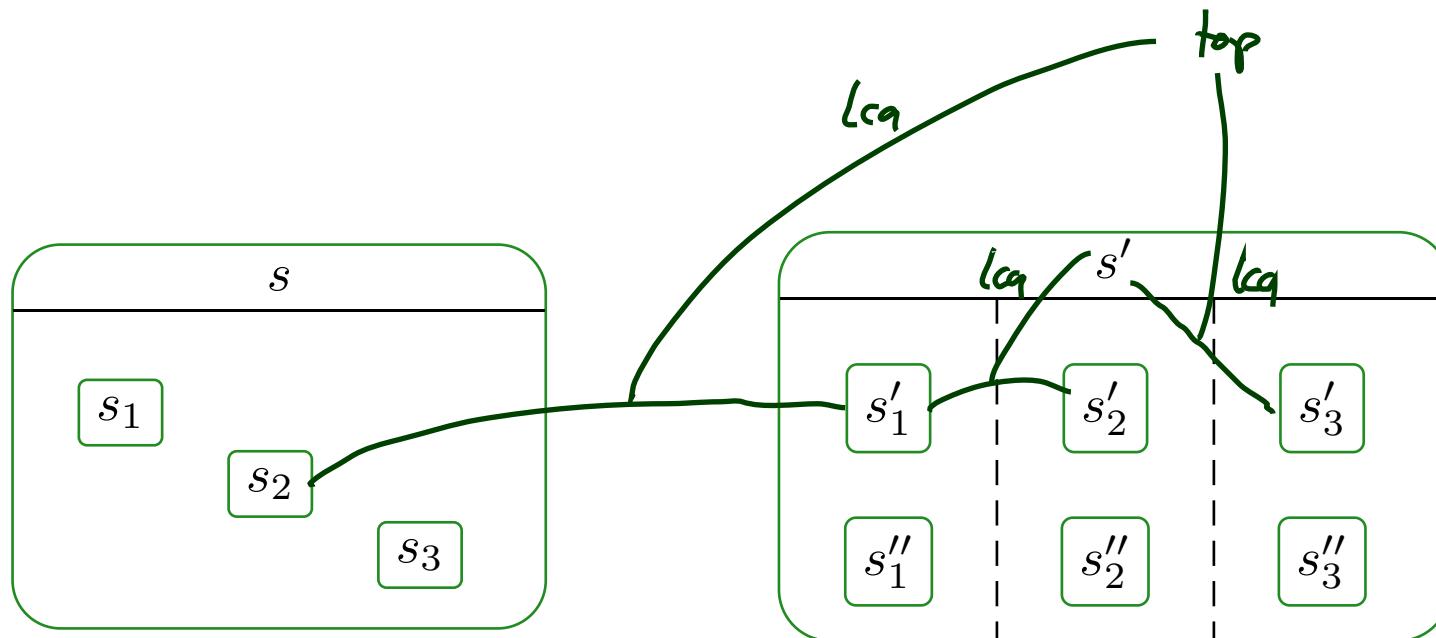
The substate- (or **child-**) relation **induces** a **partial order on states**:

- $\text{top} \leq s$, for all $s \in S$,
- $s \leq s'$, for all $s' \in \text{child}(s)$,
- transitive, reflexive, antisymmetric,
- $s' \leq s$ and $s'' \leq s$ implies $s' \leq s''$ or $s'' \leq s'$.



Least Common Ancestor

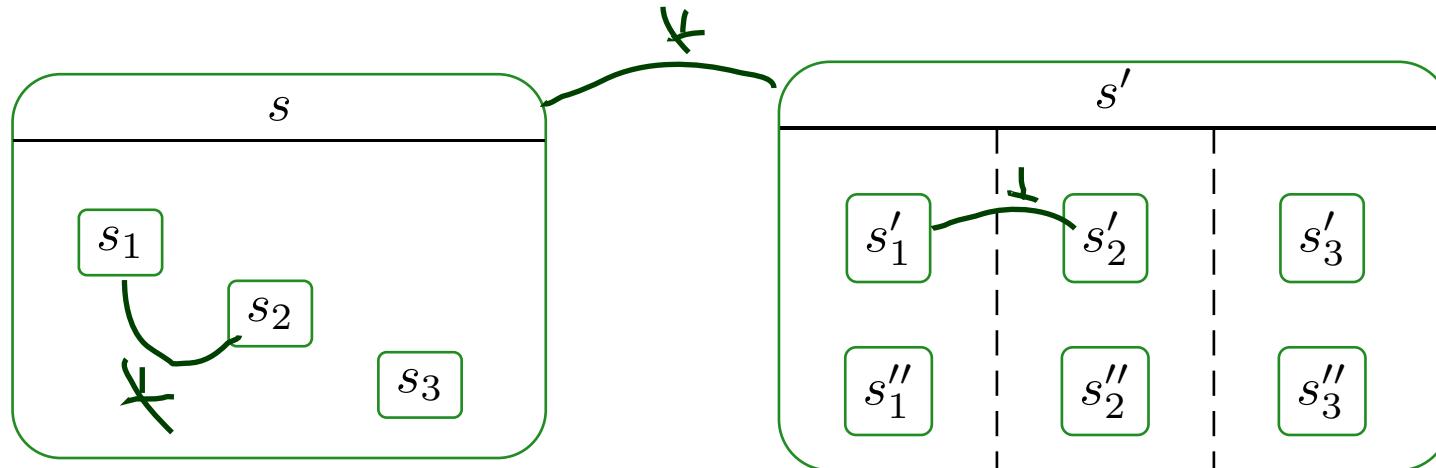
- The **least common ancestor** is the function $lca : 2^S \rightarrow S$ such that
 - The states in S_1 are (transitive) children of $lca(S_1)$, i.e.
$$lca(S_1) \leq s, \text{ for all } s \in S_1 \subseteq S,$$
 - $lca(S_1)$ is **maximal**, i.e. if $\hat{s} \leq s$ for all $s \in S_1$, then $\hat{s} \leq lca(S_1)$
- **Note:** $lca(S_1)$ exists for all $S_1 \subseteq S$ (last candidate: *top*).



Orthogonal States

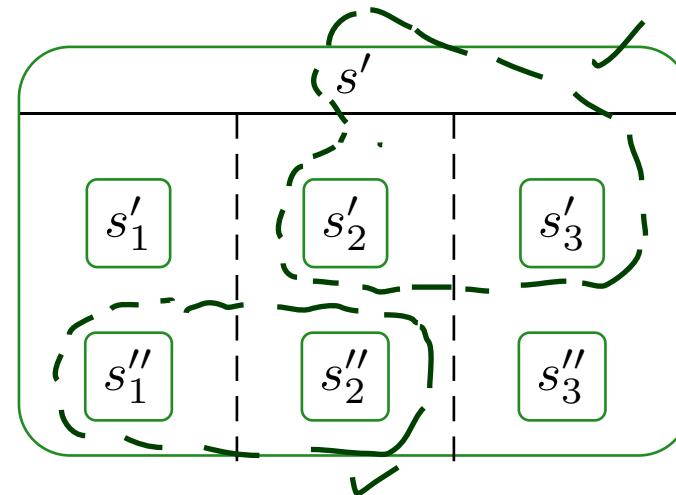
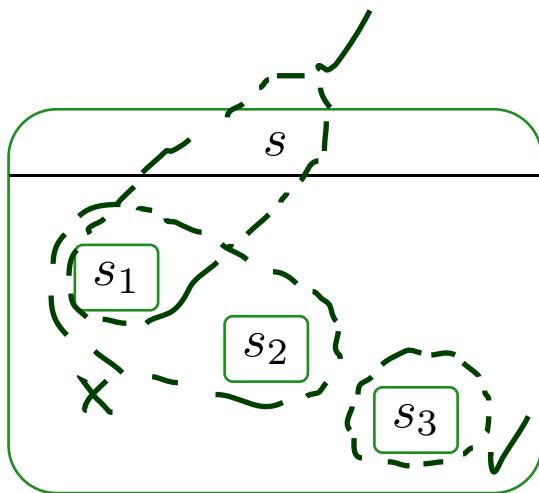
- Two states $s_1, s_2 \in S$ are called **orthogonal**, denoted $s_1 \perp s_2$, if and only if
 - they are unordered, i.e. $s_1 \not\leq s_2$ and $s_2 \not\leq s_1$, and
 - they live in different regions of an AND-state, i.e.

$$\exists s, \text{region}(s) = \{S_1, \dots, S_n\}, 1 \leq i \neq j \leq n : s_1 \in \text{child}(S_i) \wedge s_2 \in \text{child}(S_j),$$



Consistent State Sets

- A set of states $S_1 \subseteq S$ is called **consistent**, denoted by $\downarrow S_1$, if and only if for each $s, s' \in S_1$,
 - $s \leq s'$,
 - $s' \leq s$, or
 - $s \perp s'$.



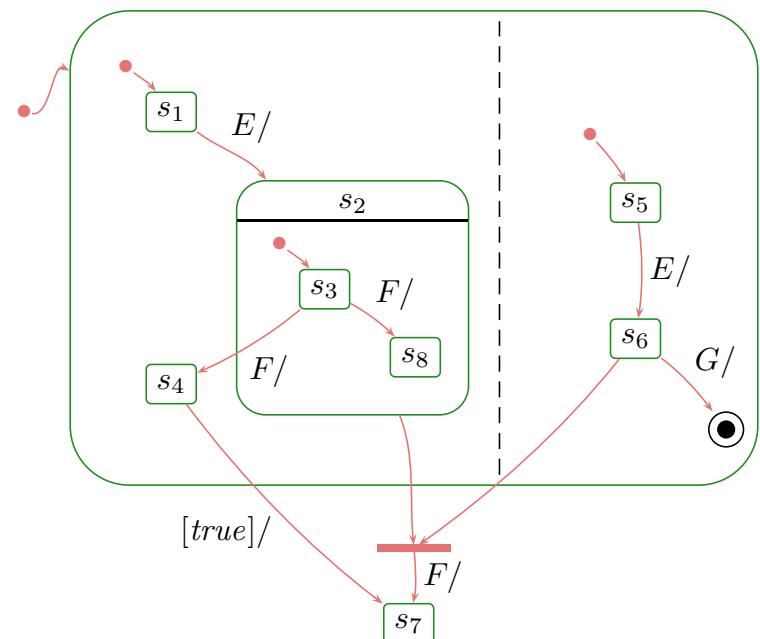
Legal Transitions

A hierarchical state-machine $(S, kind, region, \rightarrow, \psi, annot)$ is called **well-formed** if and only if for all transitions $t \in \rightarrow$,

- source and destination are consistent, i.e. $\downarrow source(t)$ and $\downarrow target(t)$,
- source (and destination) states are pairwise orthogonal, i.e.
 - forall $s, s' \in source(t) (\in target(t))$, $s \perp s'$,
- the top state is neither source nor destination, i.e.
 - $top \notin source(t) \cup target(t)$.

Recall: final states are not sources of transitions.

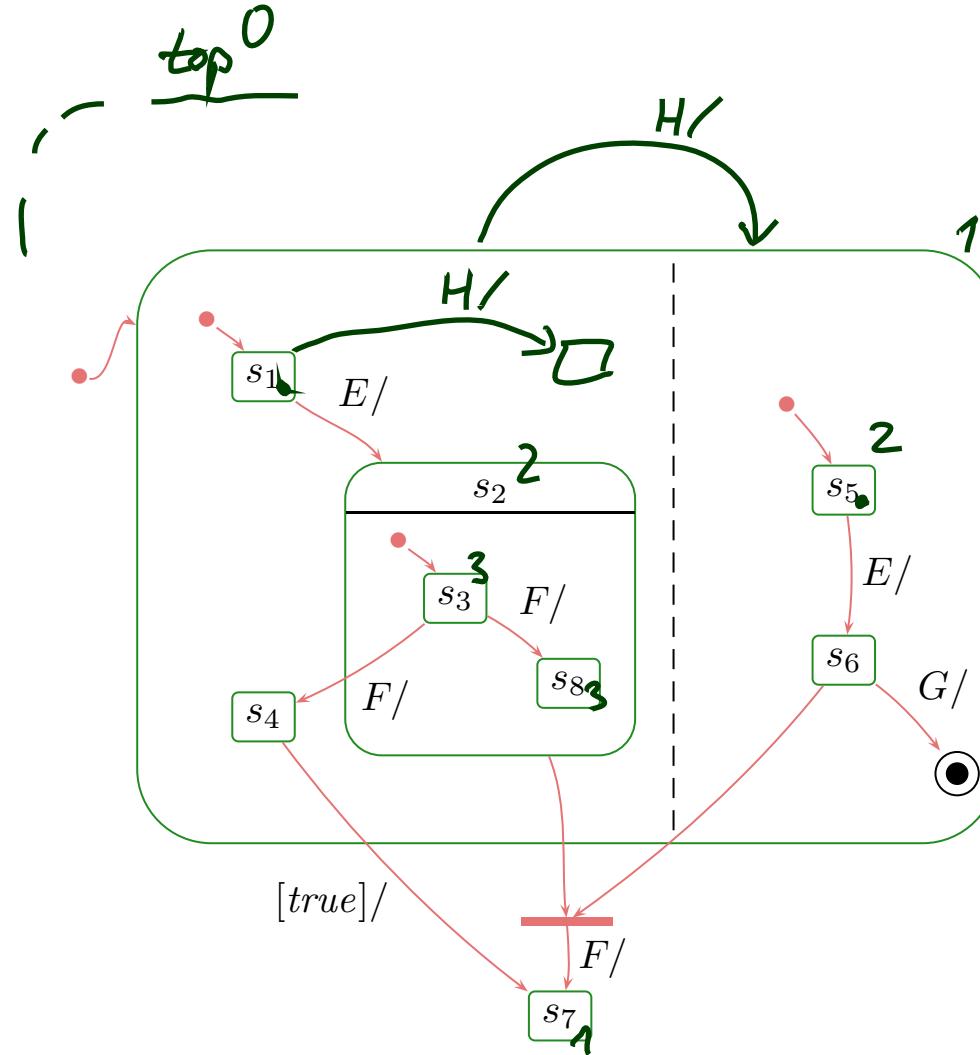
Example:



The Depth of States

- $\text{depth}(\text{top}) = 0$,
- $\text{depth}(s') = \text{depth}(s) + 1$, for all $s' \in \text{child}(s)$

Example:



Enabledness in Hierarchical State-Machines

- The **scope** ("set of possibly affected states") of a transition t is the **least common region** (!) of

$$\text{source}(t) \cup \text{target}(t).$$

- Two transitions t_1, t_2 are called **consistent** if and only if their scopes are orthogonal (i.e. states in scopes pairwise orthogonal).
- The **priority** of transition t is the depth of its innermost source state, i.e.

$$\text{prio}(t) := \max\{\text{depth}(s) \mid s \in \text{source}(t)\}$$

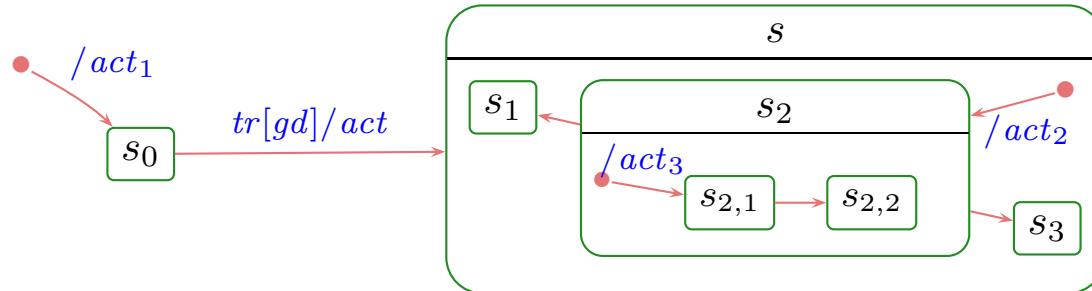
- A set of transitions $T \subseteq \rightarrow$ is **enabled** in an object u if and only if
 - T is consistent,
 - T is maximal wrt. priority (all transitions in T have the ~~same~~ highest priority),
 - all transitions in T share the same trigger,
 - for all $t \in T$, the source states are active, i.e.
$$\text{source}(t) \subseteq \sigma(u)(st) \ (\subseteq S).$$
 - all guards are satisfied by $\sigma(u)$ ~~and~~
- as before*

Transitions in Hierarchical State-Machines

- Let T be a set of transitions enabled in u .
 - Then $(\sigma, \varepsilon) \xrightarrow[u]{(cons, Snd)} (\sigma', \varepsilon')$ if
 - $\sigma'(u)(st)$ consists of the target states of T ,
i.e. for **simple states** the **simple states themselves**,
for **composite states** the **initial states**,
 - $\sigma', \varepsilon', cons$, and Snd are the effect of firing each transition $t \in T$
one by one, in any order, i.e. for each $t \in T$,
 - the exit action transformer (\rightarrow later) of all affected states, highest depth first,
 - the transformer of t ,
 - the entry action transformer (\rightarrow later) of all affected states, lowest depth first.
- ~~> adjust Rules (ii), (iii), (v) accordingly.

Initial and Final States

Initial Pseudostate



Principle:

- when entering a non-simple state,
- then go to the destination state of a transition with initial pseudo-state source,
- execute the action of the chosen initiation transition(s) **between** exit and entry actions (\rightarrow later).

Recall: For simplicity, we assume exactly one initiation transitions — could be more, choose non-deterministically.

Special case: the region of *top*.

- If class C has a state-machine, then “create- C transformer” is the concatenation of
 - the transformer of the “constructor” of C (here not introduced explicitly) and
 - a transformer corresponding to one initiation transition of the top region.

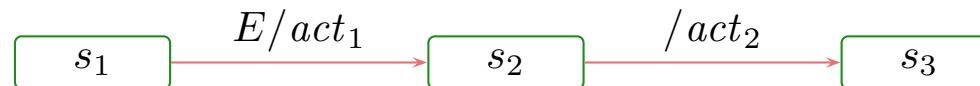
Final States



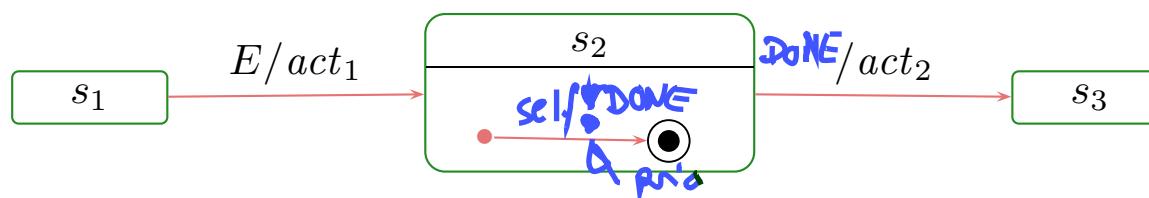
- If $(\sigma, \varepsilon) \xrightarrow[u]{(cons, Snd)} (\sigma', \varepsilon')$ and all **simple states** in $st \in \sigma(u)(st)$ are **final**, i.e. $kind(s) = fin$, then
 - stay **unstable** if there is a common parent of the simple states in $\sigma(u)(st)$ which is source of a transition without trigger and satisfied guard,
 - otherwise **kill** u .
- ~~ adjust Rules (i), (ii), (iii), and (v) accordingly.

Observation: u never “survives” reaching a state (s, fin) with $s \in child(top)$.

Observation:



vs.



References

References

OMG (2011a). Unified modeling language: Infrastructure, version 2.4.1. Technical Report formal/2011-08-05.

OMG (2011b). Unified modeling language: Superstructure, version 2.4.1. Technical Report formal/2011-08-06.