Software Design, Modelling and Analysis in UML

Lecture 6: Class Diagrams I

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UML Mathematics $v_{\pi} = ((\sigma_i, cons_i, Snd_i))_{i \in \mathbb{N}}$

Recall: Signature vs. Class Diagram



UML Class Diagrams: Stocktaking

Content

Course Map

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 (Temporary) Abbreviations

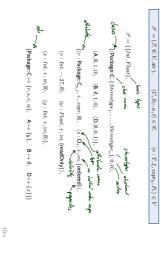
Stereotypes

That'd Be Too Simple

What Do We Want / Have to Cover?

has a set of stereotypes. has name. (belongs to a package.) can be abstract can be active. has a set of attributes. has a set of operations. (~ late.) a visibility, a name, a type, المامان a multiplicity, an order, Wanted: places in the signature to represent the information from the picture. an initial value, and Each attribute has a set of properties, such as readOnly, ordered, etc. (Stereotype, ..., Stereotype,) Package::C :Float {readOnly}

Extended Signature Example



Conventions

- $\bullet~$ We write $\langle C,S_C,a,t\rangle$ if we want to refer to all aspects of class C.
- If the new aspects are irrelevant (for a given context), we simply write C i.e. old definitions (written in terms of C) are still valid.
- $\bullet~$ Similarly, we write $\langle v:T,\xi,expr_0,P_v\rangle$ if we want to refer to all aspects of attribute v
- Write only v:T or v if details are irrelevant

Extended Signature

• '6' is a finite set of classes (C,S_C,a,t) where • S_C is a finite boosibly-empty set of streestypes. • $a \in B$ is a booken flag inducing whether C is absolute. ($a \circ f : \frac{1}{2}! \in d'$ is a foldwed) • $t \in B$ is a booken flag indicating whether C is active.

Definition. An (Extended) Object System Signature is a quadruple $\mathscr{S}=(\mathscr{T},\mathscr{C},V,atr)$ where

F is a set of (basic) types,

• V is a finite set of attributes $\langle v:T,\xi,expr_0,P_v\rangle$ where • T is a type from \mathscr{T} , or $C_{0,1},C_*$ for some $C\in\mathscr{C}$.

* $\xi \in \{\text{public}, \text{private}, \text{protected}, \text{package}\}\$ is the visibility.

Extended Signature

We use $S_\mathscr{C}$ to denote the set $\bigcup_{C\in\mathscr{C}}S_C$ of stereotypes in $\mathscr{S}.$ atr: % → 2^V maps each class to its set of attributes.

an initial value expression capto given as a year from a language for initial value expressions, e.g. OCL or C-in the Rhapsody-root: write for to explicitly not go an initial unite expression.
 a finite (possibly empty) set of properties P_r.

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Structures of Extended Signatures

Recall:

```
Definition. A Basic Object System Structure of a Basic Object System Signature \mathscr{S}=(\mathscr{T},\mathscr{C},V,atr) is a domain function \mathscr{D} which assigns to each type a domain, i.e.

    Sets of object identities for different classes are disjoint, i.e.

                                                                                                                                                                                                                                                                                                                                                                      • C\in\mathscr{C} is mapped to an infinite set \mathscr{D}(C) of (object) identities. Note: Object identities only have the "=" operation
We use \mathscr{D}(\mathscr{C}) to denote \bigcup_{C \in \mathscr{C}} \mathscr{D}(C); analogously \mathscr{D}(\mathscr{C}_*).
                                                                                             • C_* and C_{0,1} for C \in \mathscr{C} are mapped to 2^{\mathscr{D}(C)}.
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                 • \tau \in \mathscr{T} is mapped to \mathscr{D}(\tau).
                                                                                                                                                                                          \forall\, C, D \in \mathscr{C} \colon C \neq D \to \mathscr{D}(C) \cap \mathscr{D}(D) = \emptyset.
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Structures of Extended Signatures

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                                                                                                                                                                                                                                                                                                                                                    • C\in\mathscr C is mapped to an infinite set \mathscr D(C) of (abject) identities. Note. Object identities only have the "=" operation.
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                Definition. An (Object System) Structure of an (Extended Object System) Signature \mathscr{S}=(\mathscr{T}\mathscr{C},V,atr) is a domain function \mathscr{D} which assigns to each type a domain i.e.
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                    • \tau \in \mathscr{T} is mapped to \mathscr{D}(\tau).
                                                                                                                                                                             \forall\, C,D\in\mathcal{C}:C\neq D\rightarrow\mathcal{D}(C)\cap\mathcal{D}(D)=\emptyset.
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Semantical Relevance

- \circ The semantics (or meaning) of an extended object system signature $\mathscr S$ wrt. a structure $\mathscr D$ is the set of system states $\Sigma \mathscr D$.
- \circ The semantics (or meaning) of an extended object system signature $\mathscr S$ is the set of sets of system states wrt. some structure of $\mathscr S$, i.e. the set

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Which of the following aspects is semantically relevant, i.e. does contribute to the constitution of system states?
• has a set of stereotypes, X.
• has name, V
• basing two panings:
• can be abstract, V
• can be advice, X
• has a set of attributes, V
• has a set of operations (later).
                                                                                                                              (2) 9 is structure of タ) ・ 犬ょくこ ガウウン トレスター トレスター トレスター アンスター アンス
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System States of Extended Signatures

System States of Extended Signatures

Definition. Let \mathscr{D} be a structure of extended signature $\mathscr{S}=(\mathscr{T},\mathscr{C},V,atr).$ state of $\mathscr S$ wrt. $\mathscr D$ is a type-consistent mapping $\sigma: \mathcal{D}(\mathcal{C}) \twoheadrightarrow (V \twoheadrightarrow (\mathcal{D}(\mathcal{T}) \cup \mathcal{D}(\mathcal{C}_*))).$

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\bullet \ \ \sigma(u)(v) \in \mathscr{D}(D_*) \ \text{if} \ v : D_{0,1} \ \text{or} \ v : D_* \ \text{with} \ D \in \mathscr{C}
                                                                                                                                                                                                                                                             • \sigma(u)(v) \in \mathscr{D}(\tau) \text{ if } v : \tau, \tau \in \mathscr{T}
                                                                                                                                                                                                                                                                                                                                                                              • dom(\sigma(u)) = atr(C)
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                        A system state of {\mathscr S} wrt. {\mathscr D} is a type-consistent mapping
We call u \in \mathscr{D}(\mathscr{C}) alive in \sigma if and only if u \in \text{dom}(\sigma).
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                  That is, for each u \in \mathscr{D}(C), C \in \mathscr{C}, if u \in \mathrm{dom}(\sigma)
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                      Definition. Let \mathscr D be a basic structure of basic signature \mathscr S=(\mathscr T,\mathscr C,V,dr).
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                        \sigma\colon \mathscr{D}(\mathscr{C}) \nrightarrow (V \nrightarrow (\mathscr{D}(\mathscr{T}) \cup \mathscr{D}(\mathscr{C}_*))).
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We use $\Sigma_{\mathscr{S}}^{\mathscr{D}}$ to denote the set of all system states of \mathscr{S} wrt. \mathscr{D} .

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• $\sigma(u)(v) \in \mathscr{D}(\tau)$ if $v : \tau, \tau \in \mathscr{T}$ • $dom(\sigma(u)) = atr(C)$ That is, for each $u \in \mathscr{D}(C), C \in \mathscr{C}$, if $u \in \text{dom}(\sigma)$

We use $\Sigma_{\mathscr{S}}^{\mathscr{D}}$ to denote the set of all system states of \mathscr{S} wrt. $\mathscr{D}.$

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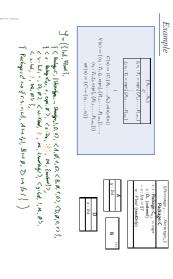
We call $u \in \mathscr{D}(\mathscr{C})$ alive in σ if and only if $u \in \text{dom}(\sigma)$.

Mapping UML Class Diagrams to Extended Signatures

A class box n induces an (extended) signature class as follows: n: $(S_1, \dots, (S_k))$ "abstract" is determined by the font: $V(n) := \{ \langle \widetilde{p_1} : T_1, \xi_1, expr_0^{\ell}, \{P_{1,1}, \dots, P_{1,m_1} \} \rangle, \dots, \langle v_{\ell} : T_{\ell}, \xi_{\ell}, expr_0^{\ell}, \{P_{\ell,1}, \dots, P_{\ell,m_{\ell}} \} \} \}$ $C(n) := \langle C, \{S_1, ...$ $atr(n) := \{C \mapsto \{v_1, \dots, v_\ell\}\}$..., S_k }, a(n), t(n)} "active" is determ

From Class Boxes to Extended Signatures

 $a(n) = \begin{cases} \textit{fue} & \text{, if } n = \boxed{\textit{C}} \text{ or } n = \boxed{\textit{C}_{(A)}} \\ \textit{false} & \text{, otherwise} \end{cases}$ $l(n) = \begin{cases} \textit{fase} & \text{.if } n = \boxed{\texttt{C}} \text{ or } n = \boxed{\texttt{C}} \end{cases}$



From Class Diagrams to Extended Signatures

Is the Mapping a Function?

(1) A class ${\cal C}$ may appear in multiple class diagrams: There are two possible sources for problems: Question: Is $\mathscr{S}(\mathscr{C}\mathscr{D})$ well-defined?

- We view a class diagram CD as a graph with nodes $\{n_1,\dots,n_N\}$ (each "class rectangle" is a node). $\mathcal{C}(CD) := \mathbb{C}(n_k) \mid 1 \leq i \leq N\}$ $V(CD) := \mathbb{U}_{i=1}^N V(n_i)$
- $atr(CD) := \bigcup_{i=1}^{N} atr(n_i)$
- In a UML model, we can have finitely many class diagrams.

which induce the following signature: $\mathscr{C}\mathscr{D} = \{CD_1, \dots, CD_k\},\$

$$\mathcal{S}(\mathcal{C}\mathcal{D}) = \left(\mathcal{F}, \bigcup_{i=1}^{k} \mathcal{C}(\mathcal{CD}_{i}), \bigcup_{i=1}^{k} V(\mathcal{CD}_{i}), \bigcup_{i=1}^{k} atr(\mathcal{CD}_{i})\right).$$

$$\mathcal{S}(\mathcal{C}\mathcal{G}) = \left(\mathcal{F}, \bigcup_{i=1}^k \mathcal{C}(\mathcal{CD}_i), \bigcup_{i=1}^k V\left(\mathcal{CD}_i\right), \bigcup_{i=1}^k atr(\mathcal{CD}_i)\right).$$

(Assuming $\mathcal T$ given. In "reality" (i.e. in full UML), we can introduce types in class diagrams, the class diagram then contributes to $\mathcal T$. Example: enumeration types.)

Simply forbid the case (ii) – easy syntactical check on diagram.

What If Things Are Missing?

- Visibility: There is no "no visibility" an attribute has a visibility in the (extended) signature.
 Some (and we) assume public as default but conventions may vary.

Example Cont'd

$$\begin{split} C(n) &= \{C, \{S_1, ..., S_k\}, a(n), t(n)\} \\ V(n) &= \{(v_1, T_1, \xi_1, copr_{1}^{1}, \{P_{1,1}, ..., P_{sm_k})\}, ..., \\ (v_1, T_k, \xi_r, copr_{1}^{1}, \{P_{1,1}, ..., P_{sm_k})\}\} \\ atr(n) &= \{C \mapsto \{v_1, ..., v_k\}\} \end{split}$$

 $\mathcal{S} = \{\{Int, Float\},\$

 $\{(\mathsf{Package}.\mathsf{C}, \{\mathit{Stereotype}_1, \dots, \mathit{Stereotype}_n\}, 0, 0),$

 $\{\langle r: \mathsf{Package-C}_{0,1}, +, expr, \emptyset \rangle, \quad \langle s: \mathsf{D}_*, +, \bowtie, \{\mathsf{ordered}\} \rangle,$

 $\{ \mathsf{Paclage=C} \mapsto \{r, s, v, w\}, \quad \mathsf{A} \mapsto \{y\}, \quad \mathsf{B} \mapsto \emptyset, \quad \mathsf{D} \mapsto \{x\}\} \})$

 $\langle x: Int, \ ^{\mathbf{L}}, \bowtie, \not {\phi} \rangle, \quad \langle y: Int, +, \bowtie, \not {\phi} \rangle \},$ $\langle v:Int,-,27, \pmb{\mathcal{B}}\rangle, \quad \langle w:Float,+,\bowtie, \{\mathsf{readOnly}\}\rangle,$ $\langle \mathbf{A},\emptyset,1,0\rangle,\quad \langle \mathbf{B},\emptyset,1,0\rangle,\quad \langle \mathbf{D},\emptyset,0,1\rangle\},$

What does the standard say? (OMG, 2011a, 121)

"Presentation Options.

The type, visibility, default, multiplicity, property string may be suppressed from being displayed, even if there are values in the model."

- Initial value: some assume it given by domain (such as "leftmost value".but what is "leftmost" of [Z?].

 Some (and we) understand non-deterministic initialisation if not given.
- Properties: probably safe to assume \(\theta \) if not given at all.

Is the Mapping a Function?

(2) An attribute v may appear in multiple classes with different type: $\frac{C}{v:Bool}$ $\frac{D}{v:Int}$



Two approaches:

• Require unique attribute names. This requirement can easily be established (implicitly, behind the scenes) by viewing v as an abbreviation for $\frac{(\mathcal{L}_{2^{n}})}{\mathcal{L}} \text{ or } \underbrace{\mathcal{D} \in \mathcal{D}_{2^{n}}}_{\mathcal{L}} \mathcal{D} \in \mathbb{R}^{n}) \ .$ depending on the context $(\mathcal{L}_{2^{n}})$: Bol and $D_{2^{n}}$: Int are then unique.)

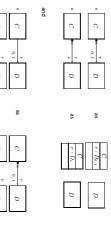
Subtle, formalist's approach: observe that

are different things in V. $\begin{array}{cccc} \langle v:Bool,\ldots\rangle & \text{and} & \langle v:Int,\ldots\rangle \\ & & & & & & & \\ & & & & & & \\ & & & & & & \\ & & & & & & \\ & & & & & & \\ & & & & & & \\ & & & & & & \\ & & & & & & \\ & & & & & \\ & & & & & \\ & & & & & \\ & & & & & \\ & & & & & \\ & & & & & \\ & & & & & \\ & & & & \\ & & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & \\ & & & \\ & \\ & & \\ & & \\ & \\ & & \\ & \\ & \\ & & \\ & \\$

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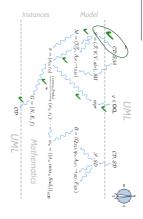
Abbreviations

Since we have not yet discussed associations, for now we read



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Course Map



Stereotypes

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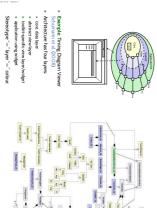
Example: Stereotypes for Documentation

Stereotypes as Labels or Tags

What are Stereotypes?

Not represented in system states.

Not contributing to typing rules / well-formedness.



Useful for documentation and model-driven development, e.g. code-generation:

View stereotypes as (additional) "labelling" ("tags") or as "grouping".

Documentation: e.g. layers of an architecture.

Sometimes, packages (cf. OMG (2011a.b)) are sufficient and "right".

Model Driven Architecture (MDA): later.

Other Examples * Use stereotypes Team; Team; Team; and assign stereotype Team, to class $\mathcal O$ if Team, is responsible for class $\mathcal C$.

Use stereotypes 'Server A', 'Server B' to indicate where objects should be stored.

Use stereotypes to label classes with licensing information (e.g., LGPL vs. proprietary).

Use stereotypes to label classes with states in the development process like "under development," "submitted for testing," "accepted."

Necessary: a common idea of what each stereotype stands for.
(To be defined / agreed on by the team, not the job of the UML consortium.)

- Extended Signatures allow us to represent aspects like
 abstract, active, visibility, initial value expression, ...
- Not all of these aspects are semantically relevant.
- The only change on system states is that abstract classes cannot have instances.

 Class Diagrams map to Extended Signatures, i.e. the meaning of a class diagram is the extended signature which it uniquely denotes.
- Thus a Class Diagram (transitively) denotes a set of system states (given a structure).
- Stereotypes are just labels.

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References

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