Software Design, Modelling and Analysis in UML

Lecture 8: Class Diagrams III

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* Class diagram (such trans y association)

* Class diagram (such

Wanted: places in the signature to represent the information from the picture

According to the second second

Klasse 1

Architect Annual Control

Association

Association

So, What Do We (Have to) Cover?

Recall: Plan & Extended Signature

Content

Recall: Associations
 Overview & Flan
 (*e Overporarily) Extend Signature

From Class Diagrams to Signatures
 Wheat if Things are Missing?
 Association Semantics
 Links in System States
 Associations and OCL

The Rest

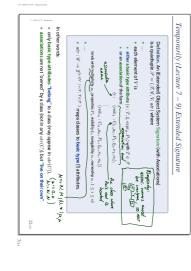
Visibility, Navigability

Multiplicity, Properties,

Ownership, "Diamonds"

Chillip 18

Back to the Main Track



Associations in Class Diagrams

What If Things Are Missing?

Association Example

 $r \blacktriangleright + n$ 0..* x:Int

Most components of associations or association end may be omitted. For instance (OMG, 2011b. 17), Section 6.4.2, proposes the following rules:

Name: Use

 $A = \langle C_1 \rangle = \cdots = \langle C_n \rangle$

Example: if the name is missing.

C ACD D for

 Role Name: use the class name at that end in lower-case letters Reading Direction: no default.

 $\mathcal{S} = \left(\left\{ \left| \mathsf{id} \right\rangle, \left\{ \mathcal{C}, \mathsf{D} \right\rangle, \left\{ \langle x, \mathsf{Ld}_y + \mathsf{M}, \mathcal{B}' \rangle, \left\langle x, \mathsf{C}_y \right\rangle, \left\langle x, \mathsf{C}_y$

C c d D for

Other convention: (used e.g. by modelling tool Rhapsody)

 $C = \frac{itsC - itsD}{D} = D$ for

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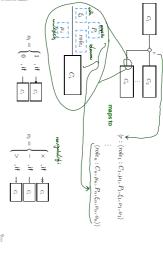
C

 $\label{eq:continuity} Multiplicity: 1 In my opinion, its safet to assume 0..1 or * (for 0..*) if there are no fixed, written, agreed conventions ("expect the worst").$

• Properties: (in course Europe !)

Visibility, public
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From Association Lines to Extended Signatures



What If Things Are Missing?

Some all arrows and × ½. Manigation and its absence are made completely explicit.

See suppress all arrows and × ½ tho elevence can be down about nonligation.

This is similar to any situation in which information is suppressed from a when the second the second to th

Suppress arrows for a sociations with non-goalbilly in both directions, and show arrows only for associations with new even ronaglability in and show arrows only for associations with new even ronaglability countries from the distinguished from situations in this case accurs rarely in practice." where there is no nonigation at all; however, the latter case occurs rarely in practice."

Wait, If Omitting Things...

- ...is causing so much trouble (e.g. leading to misunderstanding).
 why does the standard say "in practice, it is often convenient..."?
- Is it a good idea to trade convenience for precision/unambiguity?

- Convenience as such is a legitimate goal.
- In UML-As-Sketch mode, precision "doesn't matter",
 so convenience (for writer) can even be a primary goal.
- In UML-As-Blueprint mode, precision is the primary goal.
 And misunderstandings are in most cases annoying.
- But: (even in UML-As-Blueprint mode)
- If all associations in your model have multiplicity *, then it's probably a good idea not to write all these *s.
- So: tell the reader about your convention and leave out the \ast 's.

Associations: Semantics

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Links in System States

 $\langle r:\langle role_1:C_1,_,P_1,_,_,_\rangle,\ldots,\langle role_n:C_n,_,P_n,_,_,_\rangle$

Only for the course of lecture 3/8 / 9 we change the definition of system states:

• a type-consistent mapping (as before) $\sigma: \mathcal{Q}(\mathcal{C}) \to (atr(\mathcal{C}) \to \mathcal{Q}(\mathcal{T})),$ A system state of & wrt D is a pair (0, 1) consisting of only banic type articles before the desired attributes here • a mapping λ which maps each association $\langle r:\langle role_1:C_1\rangle,\ldots,\langle role_n:C_n\rangle\rangle\in V$ to a relation Definition. Let $\mathscr D$ be a structure of the (extended) signature with associations $\mathscr S=(\mathscr S,\mathscr C,V,at^{}).$

(i.e. a set of type-consistent n-tuples of identities).

 $\lambda(r) \subseteq \mathscr{D}(C_1) \times \cdots \times \mathscr{D}(C_n)$

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Association / Link Example

(12,35,24) 17/34

Associations in General

Recall: We consider associations of the following form:

Only these parts are relevant for extended system states: $\langle r:\langle role_1:C_1,\mu_1,P_1,\xi_1,\nu_1,o_1\rangle,\ldots,\langle role_n:C_n,\mu_n,P_n,\xi_n,\nu_n,o_n\rangle\rangle$

(recall: we assume $P_1=P_n=\{\mathrm{unique}\}$). $\langle r : \langle role_1 : C_1, _, P_1, _, _, _ \rangle, \dots, \langle role_n : C_n, _, P_n, _, _ \rangle$

The UML standard thinks of associations as n-ary relations which "live on their own" in a system state.

That is, links (= association instances)

 are "first-class citizens" next to objects. do not belong (in general) to certain objects (in contrast to pointers, e.g.)

are (in general) not directed (in contrast to pointers).

Associations and OCL

OCL and Associations: Syntax

Recall: OCL syntax as introduced in Lecture 3, interesting part

 $\begin{array}{lll} expr ::= \dots & |r_1(expr_1) & :\tau_C \to \tau_D & r_1 : D_{0,1} \in atr(C) \\ |r_2(expr_1) & :\tau_C \to Set(\tau_D) & r_2 : D_* \in atr(C) \end{array}$

 $\langle r: \dots, \langle role: D, \mu_{n-1-1}, \dots, \langle role: D, \mu_{n-1-1}, \dots \rangle \in V$ or $\langle r: \dots, \langle role: C, \dots, \dots \rangle \in V$, $role \neq role'$. $\begin{array}{ll} expr ::= \dots & | \ rote (expr_1) &: \tau_C \to \tau_D & \mu = 0..1 \ \text{or} \ \mu = 1..1 \\ | \ rote (expr_1) &: \tau_C \to Set(\tau_D) & \text{otherwise} \end{array}$

- Association name as such does not occur in OCL syntax, role names do.
 expr₁ has to denote an object of a dass which "participates" in the association.

 $\begin{array}{ll} TL \mp 2 ((\alpha_{1}^{2} \lambda_{1}^{2} b) & \underset{(\alpha_{1}^{2} \lambda_{1}^{2} b) = \{ c \\ u_{+} \ge L \ge 3 ((\alpha_{1}^{2} \lambda_{1}^{2} b) \le 5 \} \\ u_{+} \ge L \ge 3 ((\alpha_{1}^{2} \lambda_{1}^{2} b) + \lambda_{1}^{2} + \lambda_{2}^{2} b) + 2 \\ & \underset{(\alpha_{1}^{2} \lambda_{2}^{2} b) = \{ (\alpha_{2}^{2} \lambda_{1}^{2} b) + \lambda_{2}^{2} + \lambda_{2}^{2} b \} \\ & L(b)(u_{+}, \lambda_{1}) = \left\{ (\alpha_{2}^{2} \lambda_{1}^{2} b) + \lambda_{2}^{2} + \lambda_{2}^{2} b \} \right\} \lambda^{2} \\ & L(b)(u_{+}, \lambda_{1}^{2}) = \left\{ (\alpha_{2}^{2} \lambda_{1}^{2} b) + \lambda_{2}^{2} + \lambda_{2}^{2} b \right\} \\ & L(b)(u_{+}, \lambda_{1}^{2}) = \left\{ (\alpha_{2}^{2} \lambda_{1}^{2} b) + \lambda_{2}^{2} + \lambda_{2}^{2} b \right\} \\ & L(b)(u_{+}, \lambda_{1}^{2}) = \left\{ (\alpha_{2}^{2} \lambda_{1}^{2} b) + \lambda_{2}^{2} + \lambda_{2}^{2} b \right\} \\ & L(b)(u_{+}, \lambda_{1}^{2}) = \left\{ (\alpha_{2}^{2} \lambda_{1}^{2} b) + \lambda_{2}^{2} + \lambda_{2}^{2} b \right\} \\ & L(b)(u_{+}, \lambda_{1}^{2}) = \left\{ (\alpha_{2}^{2} \lambda_{1}^{2} b) + \lambda_{2}^{2} + \lambda_{2}^{2} b \right\} \\ & L(b)(u_{+}, \lambda_{1}^{2}) = \left\{ (\alpha_{2}^{2} \lambda_{1}^{2} b) + \lambda_{2}^{2} + \lambda_{2}^{2} b \right\} \\ & L(b)(u_{+}, \lambda_{1}^{2}) = \left\{ (\alpha_{2}^{2} \lambda_{1}^{2} b) + \lambda_{2}^{2} + \lambda_{2}^{2} b \right\} \\ & L(b)(u_{+}, \lambda_{1}^{2}) = \left\{ (\alpha_{2}^{2} \lambda_{1}^{2} b) + \lambda_{2}^{2} + \lambda_{2}^{2} b \right\} \\ & L(b)(u_{+}, \lambda_{1}^{2}) = \left\{ (\alpha_{2}^{2} \lambda_{1}^{2} b) + \lambda_{2}^{2} + \lambda_{2}^{2} b \right\} \\ & L(b)(u_{+}, \lambda_{1}^{2}) = \left\{ (\alpha_{2}^{2} \lambda_{1}^{2} b) + \lambda_{2}^{2} + \lambda_{2}^{2} b \right\} \\ & L(b)(u_{+}, \lambda_{1}^{2}) = \left\{ (\alpha_{2}^{2} \lambda_{1}^{2} b) + \lambda_{2}^{2} + \lambda_{2}^{2} b \right\} \\ & L(b)(u_{+}, \lambda_{1}^{2} b) + \lambda_{2}^{2} + \lambda_{2}^{2} b + \lambda_{2}^{2}$ 0.1 Student 1 Exists($s \mid s.t2 = s.t3$) $m \, b \, (expr_1) \big] \big[(\sigma, \lambda), \beta \big) := \begin{cases} L(role)(u_1, \lambda) & \text{if } u_1 \in \text{dom}(\sigma) \\ & L(role)(u, \lambda) = \big\{ (u_1, \dots, u_n) \\ & \in \lambda(r) \mid u \in \{u_1, \dots, u_n\} \big\} \downarrow i \end{cases}$ $ie(expr_1)]((\sigma,\lambda),\beta):=\begin{cases} u & \text{, if } u_1\in \text{dom}(\sigma) \text{ and } L(role)(u_1,\lambda)=\{u\}\\ \bot & \text{, otherwise} \end{cases}$ ITs@]((6,1), {sh16})- 1

OCL and Associations Semantics: Example

Associations: The Rest

OCL and Associations: Semantics

Recall: $\bullet \ I[r_1(expr_1)](\sigma,\beta) := \begin{cases} u & \text{, if } u_1 \in \text{dom}(\sigma) \text{ and } \sigma(u_1)(r_1) = \{u\} \\ \bot & \text{, otherwise} \end{cases}$ Assume $apr_1: \tau_C$ for some $C\in \mathscr{C}.$ Set $u_1:=I[axpr_1](a,\beta)\in \mathscr{D}(T_C)$ $I[\![r_2(expv_1)]\!](\sigma,\beta) \coloneqq \begin{cases} \sigma(u_1)(r_2) & \text{. if } u_1 \in \text{dom}(\sigma) \\ \bot & \text{. otherwise} \end{cases}$

 $I[\![role(expr_1)]\!]((\sigma,\lambda),\beta)$

- . We cannot simply write $\sigma(u)(rote)$. Recall: rote is (for the moment) not an attribute of object u (not in atr(C)).
- What we have is $\lambda(r)$ (with association name r, not with role name role!).
- $\begin{array}{ll} \langle r:\ldots,\langle role:D,\mu,\ldots,\dots,\langle role:G,\dots,\langle role:G,\dots,\dots,\rangle \rangle \\ \text{But tyields a set of } n\text{--tuples, of which some relate } \text{u} \text{ and some instances of } D. \\ * \textit{role denotes the position of the D in the tuples constituting the value of r.} \end{array}$

OCL and Associations: Semantics Cont'd

 $\mathbf{Assume}\ expr_1: \tau_C \ \mathsf{for} \ \mathsf{some}\ C \in \mathscr{C}. \ \mathsf{Set}\ u_1:=I[\![expr_1]\!]((\sigma,\lambda),\beta) \in \mathscr{D}(T_C).$



 $\bullet \ I[\![role(expr_1)]\!]((\sigma,\lambda),\beta) := \begin{cases} L(role)(u_1,\lambda) & \text{, if } u_1 \in \text{dom}(\sigma) \\ \bot & \text{, otherwise} \end{cases}$

 $L(\underline{vdc})(u'_{k}\lambda) = \{(u_{1}, \dots, u_{n}) \in \lambda(p) \mid u \in \{u_{1}, \dots, u_{n}\}\} \downarrow \{(x'_{1}, \dots, x_{n}) \in \lambda(p) \mid u \in \{u_{1}, \dots, u_{n}\}\} \downarrow \{(x'_{1}, \dots, x_{n}) \in \lambda(p) \mid u \in \{u_{1}, \dots, u_{n}\}\} \downarrow \{(x'_{1}, \dots, x_{n}) \in \lambda(p) \mid u \in \{u_{1}, \dots, u_{n}\}\} \downarrow \{(x'_{1}, \dots, x_{n}) \in \lambda(p) \mid u \in \{u_{1}, \dots, u_{n}\}\} \downarrow \{(x'_{1}, \dots, x_{n}) \in \lambda(p) \mid u \in \{u_{1}, \dots, u_{n}\}\} \downarrow \{(x'_{1}, \dots, x_{n}) \in \lambda(p) \mid u \in \{u_{1}, \dots, u_{n}\}\} \downarrow \{(x'_{1}, \dots, x_{n}) \in \lambda(p) \mid u \in \{u_{1}, \dots, u_{n}\}\} \downarrow \{(x'_{1}, \dots, x_{n}) \in \lambda(p) \mid u \in \{u_{1}, \dots, u_{n}\}\} \downarrow \{(x'_{1}, \dots, x_{n}) \in \lambda(p) \mid u \in \{u_{1}, \dots, u_{n}\}\} \downarrow \{(x'_{1}, \dots, x_{n}) \in \lambda(p) \mid u \in \{u_{1}, \dots, u_{n}\}\} \downarrow \{(x'_{1}, \dots, x_{n}) \in \lambda(p) \mid u \in \{u_{1}, \dots, u_{n}\}\} \downarrow \{(x'_{1}, \dots, x_{n}) \in \lambda(p) \mid u \in \{u_{1}, \dots, u_{n}\}\} \downarrow \{(x'_{1}, \dots, x_{n}) \in \lambda(p) \mid u \in \{u_{1}, \dots, u_{n}\}\} \downarrow \{(x'_{1}, \dots, x_{n}) \in \lambda(p) \mid u \in \{u_{1}, \dots, u_{n}\}\} \downarrow \{(x'_{1}, \dots, x_{n}) \in \lambda(p) \mid u \in \{u_{1}, \dots, u_{n}\}\} \downarrow \{(x'_{1}, \dots, x_{n}) \in \lambda(p) \mid u \in \{u_{1}, \dots, u_{n}\}\} \downarrow \{(x'_{1}, \dots, x_{n}) \in \lambda(p) \mid u \in \{u_{1}, \dots, u_{n}\}\} \downarrow \{(x'_{1}, \dots, x_{n}) \in \lambda(p) \mid u \in \{u_{1}, \dots, u_{n}\}\} \}$ $(-), \dots (role_n : -, -, -, -, -),), \quad role = role_i$

Given a set of n-tuples A, $A \downarrow i$ denotes the element-wise projection onto the i-th component.

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The Rest

Recapitulation: Consider the following association:

 $\langle r : \langle role_1 : C_1, \mu_1, P_1, \xi_1, \nu_1, o_1 \rangle, \dots, \langle role_n : C_n, \mu_n, P_n, \xi_n, \nu_n, o_n \rangle \rangle$

- Association name r and role names f types $role_{\ell}/C_{\ell}$ induce extended system states (σ,λ) .
- Multiplicity µ is considered in OCL syntax.
- Visibility ξ / Navigability ν : well-typedness (in a minute).

Now the rest:

- Multiplicity \(\mu\): we propose to view them as constraints.
- Properties P_i: even more typing.
- ullet Ownership o: getting closer to pointers/references.
- Diamonds: exercise.

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Navigability

Navigability is treated similar to visibility: Using names of non-navigable association ends ($\nu=\times$) are forbidden.

Example: Given



the following OCL expression is not well-typed wrt. navigability.

context D inv : role.x > 0

The standard says: navigation is... 'x':...not possible • '>':..efficient

'--':...possible

So, in general, UML associations are different from pointers / references in general!

But: Pointers / references can faithfully be modelled by UML associations.

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Multiplicities as Constraints

Recall: Multiplicity is a term of the form $N_1..N_2, \ldots, N_{2k-1}..N_{2k}$ where $N_i \leq N_{i+1}$ for $1 \leq i \leq 2k, \quad N_1, \ldots, N_{2k-1} \in \mathbb{N}. \quad N_{2k} \in \mathbb{N} \cup \{*\}.$

for each $\langle r:\dots,\langle role:D,\mu,_,_,_,_\rangle,\dots,\langle role':C,_,_,_,_,_\rangle,\dots\rangle\in V$ or $\operatorname{context} C \operatorname{inv}: (N_1 \leq \operatorname{role} \Rightarrow \operatorname{size}() \leq N_2) \text{ or } \dots \text{ or } (N_{2k-1} \leq \operatorname{role} \Rightarrow \operatorname{size}() \leq N_{2k})$ $\operatorname{contit}(N_{2k} = *$

 $\langle r:\dots,\langle role':C,_,_,_,_,\rangle,\dots,\langle role:D,\mu,_,_,_,_\rangle,\dots\rangle\in V,$ with $role\neq role'$, if $\mu\neq 0..1,\mu\neq 1..1$, and

 $\mu_{\mathrm{OCL}}^{C}(nole) := \mathsf{context}\ C\ \mathsf{inv} : \mathsf{not}(\mathsf{ocllsUndefined}(nole))$

Note: in n-ary associations with n>2, there is redundancy.

Ownership

Properties

We don't want to cover association properties in detail, only some observations (assume binary associations):



Intuitively it says:

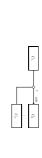
| Property | Insultion | Property | Pro

Association r is not a "thing on its own" (i.e. provided by λ), but association end 'rvie' is owned by C (!). (That is, it's stored inside C object and provided by σ).

So: if multiplicity of role is 0.1 or 1..1, then the picture above is $\underline{\text{yety.dose}}$ to concepts of pointers/references.

Actually, ownership is seldom seen in UML diagrams. Again: if target platform is clear, one may well live without (cf. (OMG, 201b, 42) for more details).

Not clear to me:



For subsets, redefines, union, etc. see (?, 127).

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 $\begin{array}{c|c} & \text{OCL Typing of expression } nle(expr) \\ \hline T_D \to Set(T_C) \\ \hline T_D \to Bag(T_C) \\ \hline T_D \to Bag(T_C) \\ \hline \text{ince} & T_D \to Seq(T_C) \\ \end{array}$

Multiplicities as Constraints Example

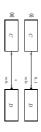
 $\mu^{G}_{\mathrm{CL}}(mk) = \mathrm{context}\ C\ \mathrm{inv}:$ $(N_1 \leq mke \to \mathrm{size}() \leq N_2)\ \mathrm{or}\ \ldots\ \mathrm{or}\ (N_{2k-1} \leq mke \to \mathrm{size}() \leq N_{2k})$



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Back to the Main Track

from now on, we only use associations of the form



- \bullet Form (i) introduces $role: C_{0,1}.$ and form (ii) introduces $role: C_{\bullet}$ in V.
- In both cases, role ∈ atr(C).
- We drop λ and go back to our nice σ with $\sigma(u)(rok) \subseteq \mathscr{D}(D)$.

Tell Them What You've Told Them. . .

- From class diagrams with (general) associations, we obtain extended signatures.
- \bullet Links (instances of associations) "live on their own" in the λ in extended system states (σ,λ) .

- OCL considers role names,
 the semantics is (more or less) straightforward. The Rest
- navigability is treated like visibility.
 view multiplicities as shorthand for constraints.
 properties ownership: "diamonds": exist/

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References

References

OMG (201ta). Unified modeling language. Infrastructure, version 2.4.1. Technical Report formal/2011-08-05.
OMG (201tb). Unified modeling language. Superstructure, version 2.4.1. Technical Report formal/2011-08-06.

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