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**Real-Time Systems**

<http://swt.informatik.uni-freiburg.de/teaching/WS2017-18/rtsys>

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Exercise Sheet 7

Early submission: Monday, 2018-01-29, 14:00

Regular submission: Tuesday, 2018-01-30, 14:00

**Exercise 1 — Queries**

**(2/20 Points)**

Task 3.(ii).f of Exercise Sheet 6 asked for a query to check that “a sensor failure is detected 10s after the failure the latest”. Consider the corresponding model with query which is available with the ILIAS exercise.

- (i) First, explain the query and outcome of the check in your own words (you may re-use your own submission to Task 6:3.(ii).f if you like).

Then disable (or delete) the two edges to fail which originate at send and at the unnamed, rightmost location in the sensor and re-check the query.

What do you observe? Discuss!

(2)

**Exercise 2 — Disjunctive Constraints**

**(1/20 Points)**

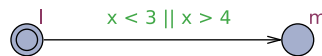


Figure 1: Disjunctive clock constraint.

Assume we need to model timed behaviour where a transition from a location  $l$  to a location  $m$  may happen at any time except for the interval  $[3, 4]$  after system startup. A naïve solution as shown in Figure 1 neither satisfies the definitions from the lectures nor is it accepted by Uppaal.

- (i) Is it principally impossible to model such behaviour, or is there a way to obtain the desired behaviour with a well-formed Uppaal model? If no, argue why not, if yes, explain your approach. (1)

**Exercise 3 — TA vs. DC Implementables**

**(5/20 Points)**

In Exercise 3.(ii) from Exercise Sheet 4, the task was to describe a design for a jamming device using implementables. One implementable from the tutorial read

$$F := [o] \xrightarrow{1} [\neg o]$$

where ‘ $o$ ’ modelled “all four channels free”.

- (i) Check (using Uppaal) whether your timed automata model of the jamming device correctly realises implementable  $F$ . (3)
- (ii) To cross-check your solution for Task (i), modify your timed automata model of the jamming device such that it *does not* realise  $F$  any more and re-check the resulting model. (1)
- (iii) Check the query ‘ $\mathbf{A}\langle\rangle \varphi$ ’ (where  $\varphi$  characterises reaching the “bad location” of the constructed observer) for the changed model from Task (ii) and interpret the outcome. (1)

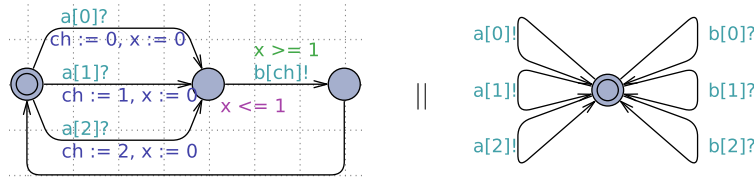


Figure 2: Network of timed automata with a timed response behaviour.

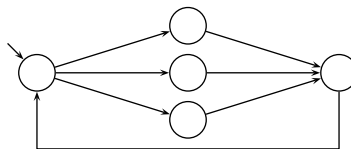
```

1 % verifyta -u -s AB.xml
2 Options for the verification:
3   Generating no trace
4   Search order is breadth first
5   Using conservative space optimisation
6   Seed is 1515780592
7   State space representation uses minimal constraint systems
8
9 Verifying formula 1: A[] true
10
11 -- Formula is satisfied.
12 -- States stored : 9 states
13 -- States explored : 9 states
14 -- CPU user time used : 0 ms
15 -- Virtual memory used : 24168 KiB
16 -- Resident memory used : 4744 KiB
    
```

Figure 3: Statistics from model-checking the network from Figure 2.

Consider the network of timed automata shown in Figure 2. The automaton on the left accepts inputs on the channels  $a[0]$ ,  $a[1]$ , and  $a[3]$ , and responds to each input on  $a[i]$  with an output on the matching channel  $b[i]$  one time unit later. To this end, the number of the input channel is stored in the variable  $ch$  to be used in the subsequent output.

At first sight, one would expect the reachable zone graph of the network to look like this:



That is, one initial configuration, one configuration for between input on  $a[i]$  and reply  $b[i]$ , and one before going back to the initial configuration.

Model-checking the network against query  $A[] \text{ true}$ <sup>1</sup> with the Uppaal command-line model-checker yields the statistics shown in Figure 3.

- (i) Why do we see a number of 9 states (i.e. configurations) explored, instead of the 5 we expected? (2)
- (ii) Can we change the model such that the behaviour on  $a$  and  $b$  (including timing) is preserved but only the expected 5 configurations are explored? (1)

*By-the-way: the model from Figure 2 can be written more concisely using select statements:*



Here,  $c : chID$  works like a free local variable with the edge as scope; the edge is enabled if one can choose (or select) a value for  $c$  such that the guard is satisfied etc.

<sup>1</sup>We check an invariant which is satisfied by the model to get an idea of the number of reachable configurations, since for a satisfied invariant, the whole set of reachable configurations needs to be examined. Otherwise, the search could stop as soon as a counter-example is found. Yet, the particular trivial query ' $A[] \text{ true}$ ' has to be used with care: some tools may recognise its triviality and immediately report "yes, satisfied" without doing any examination. Uppaal seems not to be a tool of this kind.

## Exercise 5 — Deadlock

(4/20 Points)

The query language actually supports one more atom in addition to locations and constraints, namely

`deadlock`

so that, e.g.,  $A[] \text{ deadlock}$  is a valid query.

- (i) What is the formal semantics of “`deadlock`”? Explain the intuition in your own words. (2)

*Hint: there are online documentation and offline tutorials and manuals available for Uppaal.*

- (ii) Why would anybody be interested in using this atom?

Provide a timed automata model which is supposed not to have a deadlock in this sense, check that it doesn't have a deadlock.

Then introduce an error into your model to obtain a new model which has a deadlock and check that you introduced the error correctly. (1)

- (iii) If the timed automaton



is added to your defective model from the previous task, does the resulting network still have a deadlock?

Answer the question without actually using the tool. (1)

## Exercise 6 — Model-Checking Cost

(5/20 Points)

Recall the WFAS model from Exercise Sheet 6. The model of the deterministic jamming device is parameterised in upper and lower bound on reaction time (`eps1` and `eps1`).

- (i) Collect statistics on the explored/stored states when checking your queries from Exercise Sheet 6 using the command line model-checking tool `verifyta(1)`. (3)

Consider at least the value pairs  $(1, 5)$ ,  $(0, 0)$ , and a third one of your choice. (3)

*Hint: research results should be reproducible, so please provide a small shell script (or makefile (or ...)) so that your tutor can easily check your claims from the submission.*

- (ii) Which queries are the most and second-most expensive ones wrt. explored/stored states? (1)

- (iii) What is the reason for the different costs of checking the model with different value pairs? (1)