



Tutorial for Cyber-Physical Systems - Discrete Models

Exercise Sheet 12

Exercise 1: Safety-Liveness Decomposition I 4 Points + 1 Bonus Points
The following theorem was discussed in the lecture. A proof is given below.

Theorem 1 (Decomposition). *Let AP be a set of atomic propositions, and let $\Sigma = 2^{AP}$. For every LT-property $E \subseteq \Sigma^\omega$, there exists a safety property P_{safe} and a liveness property P_{live} such that $E = P_{safe} \cap P_{live}$.*

Proof. Recall that the (prefix) closure of a property E is defined as

$$cl(E) := \{ \sigma \in \Sigma^\omega \mid pref(\sigma) \subseteq pref(E) \}$$

i.e., the set of all traces σ such that for every finite prefix w of σ , there exists some $\sigma' \in E$ such that w is also a prefix of σ' .

We set $P_{safe} := cl(E)$, and $P_{live} := (\Sigma^\omega \setminus cl(E)) \cup E$. It is easy to show that $P_{safe} \cap P_{live} = cl(E) \cap ((\Sigma^\omega \setminus cl(E)) \cup E) \stackrel{(\diamond)}{=} E$. It remains to show that P_{safe} is a safety property, and P_{live} is a liveness property. To show the former, we must prove that $cl(P_{safe}) = cl(cl(E)) \stackrel{(\clubsuit)}{=} cl(E)$. This holds due to idempotence of cl . To show the latter, we must show that $cl(P_{live}) = \Sigma^\omega$. It holds that

$$\begin{aligned} cl(P_{live}) &= cl((\Sigma^\omega \setminus cl(E)) \cup E) \\ &\stackrel{(\spadesuit)}{=} cl(\Sigma^\omega \setminus cl(E)) \cup cl(E) \\ &\stackrel{(\heartsuit)}{\supseteq} \Sigma^\omega \setminus cl(E) \cup cl(E) \\ &= \Sigma^\omega \end{aligned}$$

As it is trivially the case that $cl(P_{live}) \subseteq \Sigma^\omega$, we conclude that indeed $P_{live} = \Sigma^\omega$ and P_{live} is a liveness property. □

Consider the set of atomic propositions $AP = \{a, b\}$. Apply Theorem 1 to the property

$$E = \{ A_0 A_1 A_2 \dots \mid \exists i. b \in A_i \wedge (\forall j. j < i \rightarrow a \in A_j) \}$$

or, informally stated, “*a holds until b holds*” (and b does eventually hold).

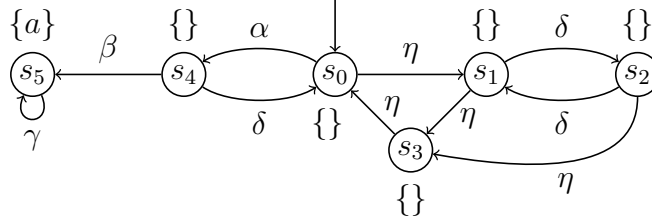
- (a) Give the decomposition of this property into a safety property P_{safe} and a liveness property P_{live} , following the construction in the proof. Prove that $E = P_{safe} \cap P_{live}$ (equation (\diamond)), that P_{safe} is indeed a safety property (equation (\clubsuit)), and that P_{live} is indeed a liveness property (equations (\spadesuit) , (\heartsuit)).
- (b) Consider the property $P'_{live} = \{ A_0 A_1 A_2 \mid \exists i. a \notin A_i \}$. Why can this property not be used in place of P_{live} in part (a)?

Exercise 2: Satisfaction under Fairness Assumptions

12 Points

The goal of this task is to train your ability to identify fair and unfair traces of a given transition system, in order to reason about properties of a system under given fairness assumptions.

Consider the following transition system:



For the fairness assumptions given in (a)–(h), perform the following tasks.

- (i) For each of the fairness assumptions below, give an execution that fulfills the fairness assumption (a fair execution) and an execution that violates the fairness assumption (an unfair execution).
- (ii) A system satisfies a property P under a given fairness assumption, if all fair traces (i.e., traces corresponding to fair executions) satisfy property P . Under which of the following fairness assumptions does the system satisfy the property “eventually a ”? Justify your answer.
 - (a) unconditional fairness for $A = \{\gamma\}$
 - (b) unconditional fairness for $A_1 = \{\alpha\}$ and for $A_2 = \{\gamma\}$
 - (c) unconditional fairness for $A = \{\alpha, \gamma\}$
 - (d) strong fairness for $A = \{\beta\}$
 - (e) strong fairness for $A_1 = \{\alpha\}$ and for $A_2 = \{\beta\}$
 - (f) strong fairness for $A_1 = \{\alpha\}$ and for $A_2 = \{\beta\}$ and for $A_3 = \{\eta\}$
 - (g) weak fairness for $A = \{\eta\}$
 - (h) weak fairness for $A_1 = \{\alpha\}$ and for $A_2 = \{\beta\}$ and for $A_3 = \{\eta\}$