

Formal Methods for Java

Lecture 15: Loops in Key

Jochen Hoenicke



Software Engineering
Albert-Ludwigs-University Freiburg

December 14, 2011

Proving loops with invariant and variant

To prove a loop in key, one needs

- a loop invariant; it must be
 - initially valid,
 - inductive, i.e. hold after executing the body if it held before,
 - strong enough to prove the post-condition (use case).
- a modifies set; this must contain all variables changed by the loop body.
- a loop variant (or ranking function); it must be
 - non-negative,
 - strictly decreasing for every execution of the loop body.

The loop variant guarantees that the loop terminates.

The rule while_invariant_with_variant_dec

The rule `while_invariant_with_variant_dec` takes an invariant inv , a modifies set $\{m_1, \dots, m_k\}$ and a variant v . The following cases must be proven.

- Initially Valid: $\implies inv \wedge v \geq 0$
- Body Preserves Invariant:

$$\implies \{m_1 := x_1 \| \dots \| m_k := x_k\} (inv \wedge [\{b = \text{COND};\}] b = \text{true}) \\ \rightarrow \langle BODY \rangle inv$$

- Use Case:

$$\implies \{m_1 := x_1 \| \dots \| m_k := x_k\} (inv \wedge [\{b = \text{COND};\}] b = \text{false}) \\ \rightarrow \langle \dots \rangle \phi$$

- Termination:

$$\implies \{m_1 := x_1 \| \dots \| m_k := x_k\} (inv \wedge v \geq 0 \wedge [\{b = \text{COND};\}] b = \text{true}) \\ \rightarrow \{old := v\} \langle BODY \rangle v \leq old \wedge v \geq 0$$

Example: Multiplication

```
/*@
 * requires a >= 0 && b >= 0;
 * ensures \result == a*b;
 */
public static int mul(int a, int b) {
    int sum = 0;
    while (b > 0) {
        sum = sum + a;
        b--;
    }
    return sum;
}
```

Loop invariant for Multiplication

One possible loop invariant is $sum + a*b = a*\text{\old}(b)$:

```
/*@ requires a >= 0 && b >= 0;
 @ ensures \result == a*b;
 @*/
public static int mul(int a, int b) {
    int sum = 0;
    /*@ loop_invariant sum + a*b == a*\text{\old}(b);
     @ modifies sum, b;
     @ decreases b;
     @*/
    while (b > 0) {
        sum = sum + a;
        b--;
    }
    return sum;
}
```

This is enough to prove it in KeY ([Demo](#))

Linear Search

Algorithm to check if an array contains an element.

```
/*@
 * @ requires arr != null;
 * @ ensures \result == (\exists int k; 0 <= k && k < arr.length;
 * @                           arr[k] == elem);
 */
public static boolean find(int[] arr, int elem) {
    for (int i = 0; i < arr.length; i++) {
        if (arr[i] == elem)
            return true;
    }
    return false;
}
```

Loop-Invariant

What is the loop invariant?

```
/*@
 * loop_invariant !(\exists k; 0 <= k && k < i; arr[k] == elem);
 * loop_invariant 0 <= i && i <= arr.length;
 * modifies i;
 * decreases arr.length - i;
 */
for (int i = 0; i < arr.length; i++) {
    if (arr[i] == elem)
        return true;
}
return false;
}
```

Demo: Binary Search

```
/*@ requires arr != null;
 @ requires (\forall int j,k; 0 <= j && j <= k && k < arr.length;
 @           arr[j] <= arr[k]); // array is sorted
 @ ensures \result == (\exists int k; 0 <= k && k < arr.length;
 @           arr[k] == elem);
 */
boolean binary(int[] arr, int elem) {
    int lower = 0, upper = arr.length - 1;
    while (lower <= upper) {
        int mid = (lower + upper) / 2;
        assert lower <= mid && mid <= upper;
        if (arr[mid] == elem) {
            return true;
        } else if (arr[mid] > elem) {
            upper = mid-1;
        } else {
            lower = mid+1;
        }
    }
    return false;
}
```

Bubble Sort

```
/*@ requires arr != null && arr.length > 0;
 @ ensures (\forall int j,k; 0 <= j && j <= k && k < arr.length;
 @           arr[j] <= arr[k]); // array is sorted
 @*/
public static boolean bubblesort(int[] arr) {
    for (int i = arr.length-1; i > 0; i--) {
        for (int j = 0; j < i; j++) {
            if (a[j] > a[j+1]) {
                int t = a[j];
                a[j] = a[j+1];
                a[j+1] = t;
            }
        }
    }
}
```

BubbleSort

Function `BubbleSort` sorts integer array `arr`



by “bubbling” the largest element of the left unsorted region of `arr` toward the sorted region on the right.

Each iteration of the outer loop expands the sorted region by one cell.

Sample execution of BubbleSort

