Formal Methods for Java Lecture 16: Abnormal Termination in Key

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Abnormal termination in Java is caused by

- a break statement,
- a continue statement,
- a return statement,
- a throw statement, or
- a statement that throws a exception.

Abnormal Termination in Dynamic Logic

The formula $\langle \alpha \rangle \phi$ holds,

• iff α terminates normally and ϕ holds afterwards.

The formula $[\alpha]\phi$ holds,

- if α terminates normally and ϕ holds afterwards.
- if α terminates abnormally.
- if α does not terminate at all.

Reasoning about exceptions.

How can we express that statement α throws an exception?

- $\langle\{\alpha\}\rangle\phi$ is equivalent to false if α throws an exception or does not terminate
- $[\{\alpha\}]\phi$ is equivalent to **true** if α throws an exception or does not terminate
- The trick is to put an exception handler into the code:

 $\langle \{ \text{Throwable thrown} = null; \\ try {\alpha; } \\ catch (Throwable ex) \{ thrown = ex; \} \} \rangle thrown \neq null$

Reasoning with try-catch blocks

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Many DL-rules in KeY just skip opening of try blocks,e.g.
\find( \<{ .. #loc = #se ... }\> post )
\replacewith( { #loc := #se } \< { .... } \> post )
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Here .. stands for an arbitrary number of opening try-blocks, labelled blocks and normal blocks.

Example:

$$\langle \{ try \{ label : \{ try \{ x = 5 \dots \} \} \} \rangle \phi$$

is replaced with

$$\{x := 5\} \langle \{ \mathsf{try} \ \{ \mathsf{label} : \{ \mathsf{try} \ \{ \ldots \} \} \}
angle \phi$$

Reasoning with try-catch blocks (2)

When an exception is thrown, the surrounding try blocks become important: \find(\<{ .. try { throw #se; #slist1 } catch (#t #v0) { #slist2 } ... }\> post) throwing a handled exception: #se instanceof #t 2 throwing an unhandled exception: ! (#se instanceof #t) \replacewith(\<{ .. throw #se; ... }\> post) Ithrowing a null pointer: #se = null \replacewith(\< { .. try { throw new NullPointerExc(); #slist1</pre> catch (#t #v0) { #slist2 } ... }\> post) The KeY system defines a single rule: \replacewith(\< { .. if (#se = null) then</pre> try { throw new NullPointerExc(); #slist1 catch (#t #v0) { #slist2 } else if (#se instanceof #t) then #t v0 = #se; #slist2 else throw #se: $\ldots \} > post)$

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