

# *Software Design, Modelling and Analysis in UML*

## *Lecture 18: Inheritance I*

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# *Contents & Goals*

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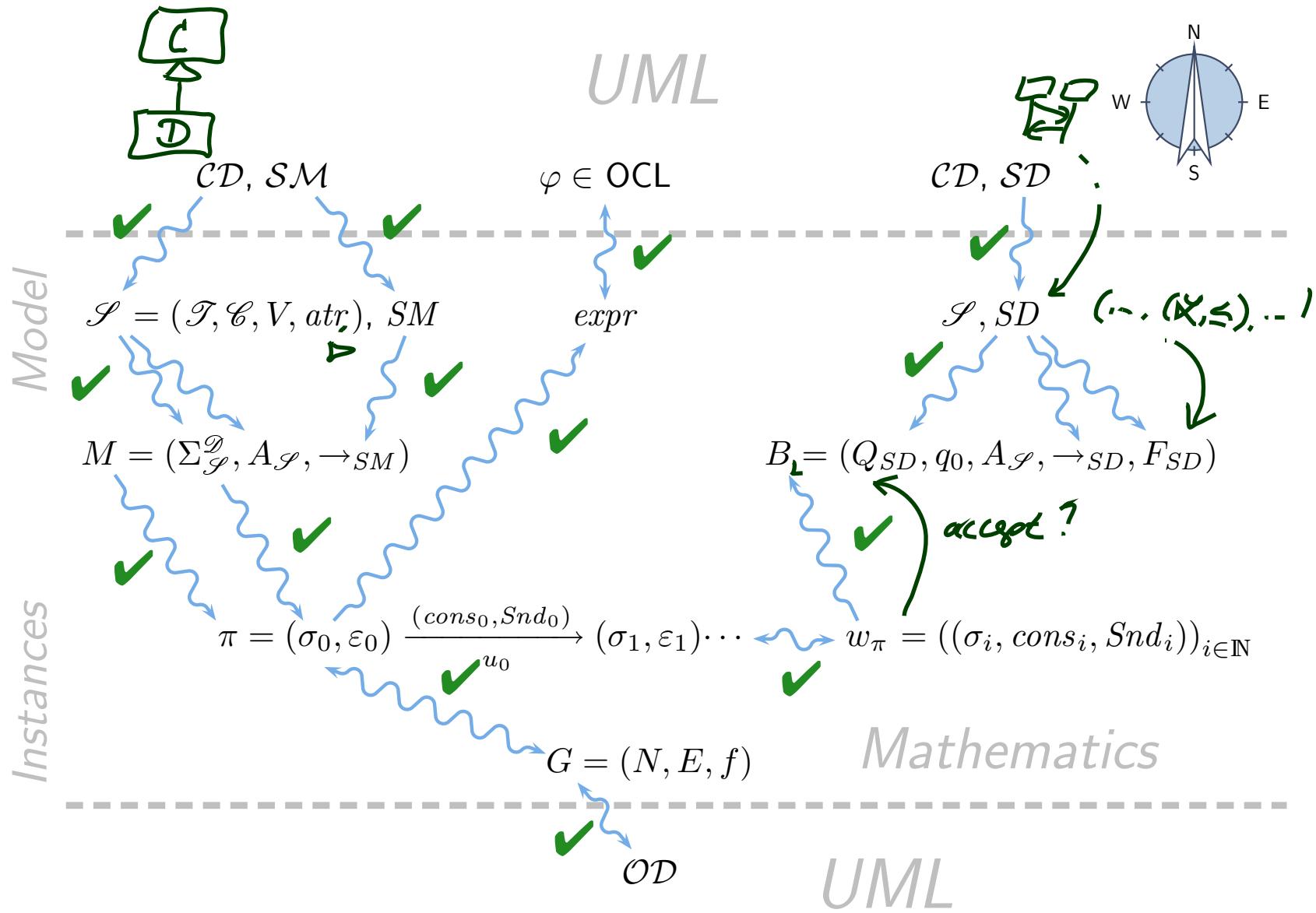
## Last Lecture:

- Live Sequence Charts Semantics

## This Lecture:

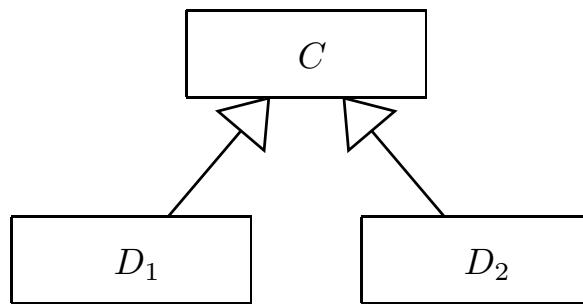
- **Educational Objectives:** Capabilities for following tasks/questions.
  - What's the Liskov Substitution Principle?
  - What is late/early binding?
  - What is the subset, what the uplink semantics of inheritance?
  - What's the effect of inheritance on LSCs, State Machines, System States?
  - What's the idea of Meta-Modelling?
- **Content:**
  - Inheritance in UML: concrete syntax
  - Liskov Substitution Principle — desired semantics
  - Two approaches to obtain desired semantics

# Course Map

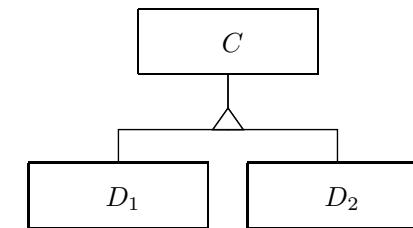
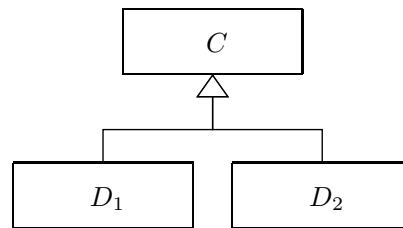
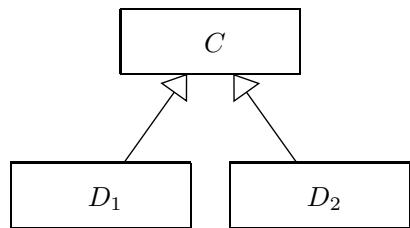


## *Inheritance: Syntax*

# Inheritance: Generalisation Relation



- Alternative renderings:



- Read:

- $C$  **generalises**  $D_1$  and  $D_2$ ;  $C$  is a **generalisation** of  $D_1$  and  $D_2$ ,
- $D_1$  and  $D_2$  **specialise**  $C$ ;  $D_1$  **is a** (specialisation of)  $C$ ,
- $D_1$  **is a**  $C$ ;  $D_2$  **is a**  $C$ .



- **Well-formedness rule:** No **cycles** in the generalisation relation.

# Abstract Syntax

**Recall:** a signature (with signals) is a tuple  $\mathcal{S} = (\mathcal{T}, \mathcal{C}, V, atr)$ .

**Now (finally):** extend to

$$\mathcal{S} = (\mathcal{T}, \mathcal{C}, V, atr, F, mth, \triangleleft)$$

*behav. fct.*  
 $\downarrow$        $\vdash : \mathcal{C} \rightarrow \mathcal{F}$

where  $F/mth$  are methods, analogously to attributes and

$$\triangleleft \subseteq (\mathcal{C} \times \mathcal{C}) \cup (\mathcal{E}(B) \times \mathcal{E}(Y))$$

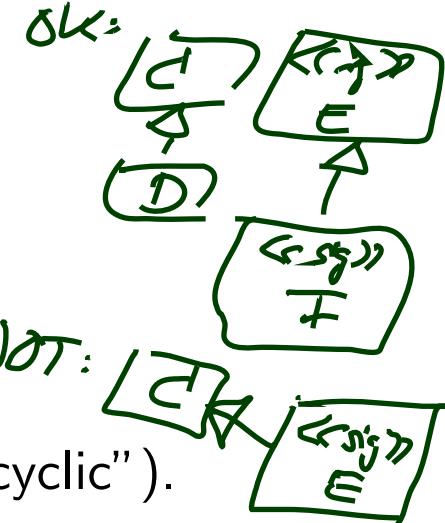
$\mathcal{E}(B)$      $\mathcal{E}(Y)$

is a **generalisation** relation such that  $C \triangleleft^+ C$  for **no**  $C \in \mathcal{C}$  ("acyclic").

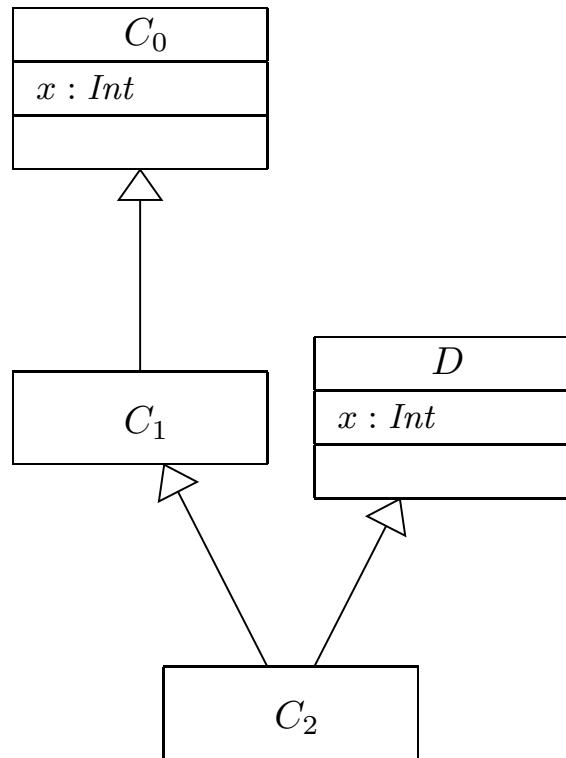
*transitiv closure*

$C \triangleleft D$  reads as

- $C$  is a generalisation of  $D$ ,
- $D$  is a specialisation of  $C$ ,
- $D$  inherits from  $C$ ,
- $D$  is a sub-class of  $C$ ,
- $C$  is a super-class of  $D$ ,
- ...



# Mapping Concrete to Abstract Syntax by Example



$\mathcal{F} = \left( \{\text{let}\}, \right.$

$\left\{ C_0, C_1, D, C_2 \right\},$

$\left\{ C_0 :: x : \text{Int}, \right.$

$D :: x : \text{Int} \left. \right\},$

$\left\{ C_0 \mapsto \{C_0 :: x\}, \right.$

$D \mapsto \{D :: x\}, C_1 \mapsto \emptyset,$

$C_2 \mapsto \emptyset ?, \right.$

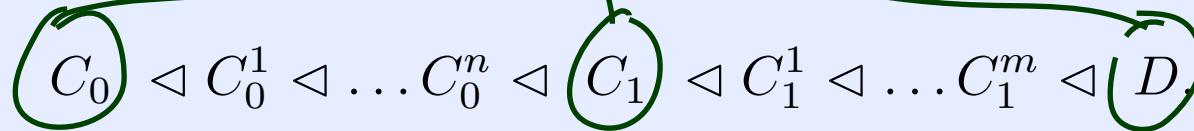
$\left. \{C_0 \Delta C_1, C_1 \Delta C_2, D \Delta C_2\} \right)$

NOT:  $\text{addr}(C_2) = \{C_0 :: x, D :: x\}$

**Note:** we can have **multiple inheritance**.

# Reflexive, Transitive Closure of Generalisation

**Definition.** Given classes  $C_0, C_1, D \in \mathcal{C}$ , we say  $D$  inherits from  $C_0$  via  $C_1$  if and only if there are  $C_0^1, \dots, C_0^n, C_1^1, \dots, C_1^m \in \mathcal{C}$  such that



We use ' $\preceq$ ' to denote the reflexive, transitive closure of ' $\lhd$ '.

In the following, we assume

- that all attribute (method) names are of the form

$$C::v, \quad C \in \mathcal{C} \cup \mathcal{E} \quad (C::f, \quad C \in \mathcal{C}),$$

- that we have  $C::v \in \text{atr}(C)$  resp.  $C::f \in \text{mth}(C)$  **if and only if**  $v$  ( $f$ ) appears in an attribute (method) compartment of  $C$  in a class diagram.

We still want to accept “context  $C$   $\text{inv} : v < 0$ ”, which  $v$  is meant? Later!

## *References*

# References

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